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Space engineering — Space data links — Telecommand protocols, synchronization and channel coding

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Space engineering - Space data links - Telecommand protocols, synchronization and channel coding

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Protocoles, synchronisation et codage canal des liaisons
télécommande

Raumfahrttechnik - Telekommando-Protokolle, -
synchronisation und -kanalkodierung

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Foreword

This document (EN 16603-50-04:2014) has been prepared by Technical Committee CEN/CLC/TC 5 "Space", the secretariat of which is held by DIN.

This standard (EN 16603-50-04:2014) originates from ECSS-E-ST-50-04C.

This European Standard shall be given the status of a national standard, either by publication of an identical text or by endorsement, at the latest by March 2015, and conflicting national standards shall be withdrawn at the latest by March 2015.

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This document has been developed to cover specifically space systems and has therefore precedence over any EN covering the same scope but with a wider domain of applicability (e.g. : aerospace).

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1

Scope

This Standard specifies the data structures and protocols for a telecommand space data link and the procedure for physical layer operation.

Usually, the source of data on a telecommand space data link is located on the ground and the receiver is located in space. However, the Standard may also be used for space-to-space telecommand data links.

Further provisions and guidance on the application of this standard can be found, respectively, in the following documents:

- The higher level standard ECSS-E-ST-50 “Communications”, which defines the principle characteristics of communication protocols and related services for all communication layers relevant for space communication (physical- to application-layer), and their basic relationship to each other.
- The handbook ECSS-E-HB-50 “Communications guidelines”, which provides information about specific implementation characteristics of these protocols in order to support the choice of a certain communications profile for the specific requirements of a space mission.

Users of this present standard are invited to consult these documents before taking decisions on the implementation of the present one.

This standard may be tailored for the specific characteristics and constraints of a space project in conformance with ECSS-S-ST-00.

2

Normative references

The following normative documents contain provisions which, through reference in this text, constitute provisions of this ECSS Standard. For dated references, subsequent amendments to, or revisions of any of these publications, do not apply. However, parties to agreements based on this ECSS Standard are encouraged to investigate the possibility of applying the most recent editions of the normative documents indicated below. For undated references the latest edition of the publication referred to applies.

EN reference	Reference in text	Title
EN 16601-00-01	ECSS-S-ST-00-01	ECSS system - Glossary of terms
	CCSDS 135.0-B-3 Issue 3, October 2006	Space Link Identifiers – Blue Book

3

Terms, definitions and abbreviated terms

3.1 Terms from other standards

For the purpose of this Standard, the terms and definitions from ECSS-S-ST-00-01 apply.

3.2 Terms specific to the present standard

3.2.1 mission phase

period of a mission during which specified telecommand characteristics are fixed

NOTE The transition between two consecutive mission phases can cause an interruption of the telecommand data link services.

3.2.2 octet

group of eight bits

NOTE 1 The numbering for octets within a data structure starts with 0.

NOTE 2 Refer to 3.4 for the convention for the numbering of bits.

3.2.3 packet

variable-length data structure consisting of higher layer user data encapsulated within standard header information

3.2.4 static

unchanged within a specific virtual channel

NOTE 1 This Standard contains requirements on the invariability, throughout one or all mission phases, of certain characteristics of the data structures specified in it.

NOTE 2 For virtual channel refer to 4.7.

3.2.5 physical layer operation procedure

procedure in the physical layer to activate and deactivate the physical communications channel by invoking RF carrier and modulation techniques

3.3 Abbreviated terms

For the purpose of this Standard, the abbreviated terms from ECSS-ST-00-01 and the following apply:

Abbreviation	Meaning
AD	acceptance check and data
ARQ	automatic repeat request
BC	bypass of acceptance check and control
BCH	Bose-Chaudhuri-Hocquenghem
BD	bypass of acceptance check and data
BER	bit error rate
BTG	bit transition generator
CCSDS	Consultative Committee for Space Data Systems
CLCW	communications link control word
CLTU	communications link transmission unit
CMM	carrier modulation mode
COP	communications operation procedure
FARM	frame acceptance and reporting mechanism
FDU	telecommand frame data unit
FECF	Frame Error Control Field
FOP	frame operation procedure
ID	identifier, or, identification
LLIF	lower layer interface
MAP	multiplexer access point
MSB	most significant bit
NRZ-M	non-return-to-zero-mark
NW	negative window
PAC	packet assembly controller
PLOP	physical layer operation procedure
PW	positive window
RF	radio frequency
SEC	single error correction
SS	Suspend_State
TC	telecommand
TED	triple error detection
TT	Timeout_Type

NOTE AC, CB and BD frames are defined in Table 6-3.

3.4 Conventions

3.4.1 Bit 0, bit 1, bit N-1

To identify each bit in an N-bit field, the first bit in the field to be transferred (i.e. the most left justified in a graphical representation) is defined as bit 0; the following bit is defined as bit 1 and so on up to bit N-1.

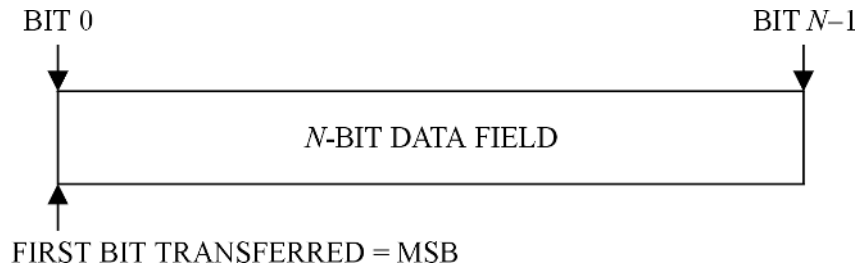


Figure 3-1: numbering convention

3.4.2 Most significant bit

When an N-bit field is used to express a binary value (such as a counter), the most significant bit is the first bit of the field, i.e. bit 0 (see Figure 3-1).

3.4.3 Use of capitals for the names of data structures and fields

In this Standard initial capitals are used for the names of data structures and fields.

This enables field names to be easily identified in the surrounding text.

For example, the field Frame Length is easier to see than frame length in text containing words such as frame and length. It also prevents ambiguity over where the name begins and ends.

4

Overview

4.1 Presentation

This Standard describes the data structures and protocols used in space missions to transmit data on telecommand space data links. It is designed to meet the requirements of space missions for the reliable transfer of telecommands over ground-to-space or space-to-space communications links.

Annex E describes the mission configuration parameters within the scope of this Standard.

The functions and protocols of the telecommand space link are grouped into layers and sublayers. The functions and protocols addressed in this Standard are grouped as follow:

- the segmentation sublayer,
- the transfer sublayer,
- the synchronization and channel coding sublayer, and
- the procedures in the physical layer at the sending end.

The segmentation sublayer, the transfer sublayer and the synchronization and channel coding sublayer are sublayers of the data link layer.

The layers and sublayers defined in this Standard provide a means for transferring data units on behalf of higher layers.

In general, the procedures, protocols and functions defined in this Standard are not concerned with the format, contents or applications of the higher layer data units. The only exception is the blocking function of the segmentation sublayer: the higher layer data units which use this function are restricted to certain packet types.

4.2 Protocol profiles

Figure 4-1 illustrates the following:

- The relationship of the layers and sublayers specified in this Standard.
- The data structures passed between the peer entities. The data structures are introduced in the overview clauses 4.3 to 4.5.

The protocol profile for a project can differ from the profile represented by Figure 4-1.

When defining the layers and sublayers in this Standard, reference is sometimes made to one of the other layers or sublayers, but such references are not intended to exclude differences from the profile represented by Figure 4-1.

The protocol profile in Figure 4-1 does not show any security sublayers. As part of the tailoring process described in Clause 1 a project with security requirements can insert security sublayers into the protocol profile of the telecommand space data link.

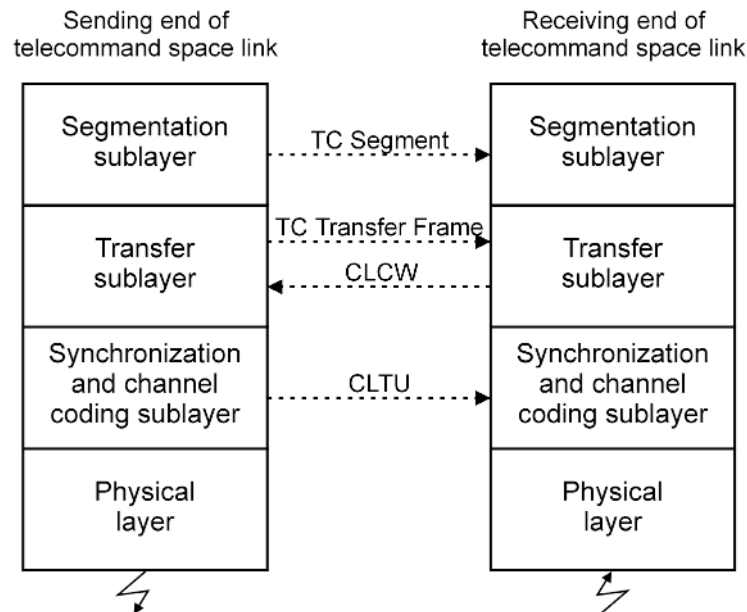


Figure 4-1: Layers and sublayers specified in this Standard

4.3 Segmentation sublayer

The TC segment contains fields to support two of the functions of the segmentation sublayer:

- the segmentation function, and
- a multiplexing function.

The segmentation function breaks long data units into smaller portions so that they are consistent with the limited length of the frames in the transfer sublayer. At the receiving end, the segmentation function reassembles the portions into the original data units.

The segmentation sublayer also includes a blocking function for packets so that multiple packets can be carried in a frame.

Clause 5 provides the specification of the segmentation sublayer and of the TC segment.

4.4 Transfer sublayer

The TC Transfer Frame is a variable-length data structure, designed for the efficient transfer of variable-length data units.

The transfer sublayer contains procedures for generating and processing TC Transfer Frames. The procedures include a multiplexing function.

The procedures are defined in Clause 6, which also includes the definition of the TC Transfer Frame.

The transfer sublayer includes the communications operation procedure (COP-1), which operates the transmission protocol of the transfer sublayer. The detailed specification of COP-1 is given in Clause 7.

COP-1 supports two service types:

- the sequence-controlled service, and
- the expedited service.

COP-1 uses protocol status information which is passed from the transfer sublayer at the receiving end to the transfer sublayer at the sending end. The information is contained in a communications link control word (CLCW), which can be transmitted in a telemetry transfer frame.

4.5 Synchronization and channel coding sublayer

The synchronization and channel coding sublayer establishes an error-controlled data channel through which data can be transferred. The data is encoded, using a block code with error-correcting capabilities, to reduce the effects of noise on the data in the space channel.

The synchronization and channel coding sublayer includes procedures for information bit randomizing, to ensure frequent bit transitions in the communications channel. Synchronization for the codeblock and delimiting of the beginning of user data are provided by the communications link transmission unit (CLTU) data structure.

Clause 8 defines the synchronization and channel coding sublayer and also contains the definition of the CLTU.

In Annex D the performance of the coding and data structures in the synchronization and channel coding sublayer is discussed and the related frame error rates in the transfer sublayer.

4.6 Physical layer

The physical layer is the lowest layer of the telecommand space link. In this Standard the physical layer is designed to work together with the synchronization and channel coding sublayer. It accepts CLTUs from the synchronization and channel coding sublayer at the sending end. At the receiving end it supplies the synchronization and channel coding sublayer with a symbol stream containing the CLTUs.

Clause 9 defines the physical layer procedures, including the associated data structures.

Other aspects of the physical layer, such as radio frequency characteristics, are outside the scope of this Standard.

4.7 Virtual channels

The TC Transfer Frame supports the division of the physical channel into multiple logically-separate virtual channels by means of identifier fields in the header of a TC Transfer Frame. The identifier fields are the Virtual Channel Identifier and the Spacecraft Identifier. The virtual channel multiplexing is a procedure of the transfer sublayer.

This Standard does not specify the manner in which the virtual channels are multiplexed into a physical channel: the Spacecraft Identifier can be included as part of the multiplexing process in a given implementation. Clause 6.9 defines virtual channel multiplexing.

Segmentation sublayer

5.1 Overview

The segmentation sublayer accepts variable-length octet-aligned data units from the next higher layer or sublayer at the sending end and delivers them to the next higher layer or sublayer at the receiving end. Within the segmentation sublayer these data units are referred to as user data units.

The segmentation sublayer provides three functions: blocking, segmentation and multiplexing.

- The blocking function combines multiple user data units so that they are carried in a single frame of the transfer sublayer. The user data units for the blocking function are limited to packets.
- The segmentation function breaks long user data units into multiple shorter portions at the sending end and reassembles them at the receiving end.
- User data units allocated to different multiplexer access points (MAPs) are multiplexed together so that they share one virtual channel of the transfer sublayer.

The principal data structure in the segmentation sublayer is the TC Segment. This is a variable-length data structure containing an integral number of octets. The TC Segment has fields in its header to support segmentation and MAP multiplexing.

If segmentation or MAP multiplexing is used for a virtual channel, then the virtual channel uses TC Segments.

If neither segmentation nor MAP multiplexing is used for a virtual channel, then the use of TC Segments on the virtual channel is optional. A physical channel can have some virtual channels that use TC Segments and some virtual channels that do not.

At the sending end, the data units generated by the segmentation sublayer are passed to the next lower sublayer:

- It is assumed that the next lower sublayer accepts variable-length data units from the segmentation sublayer.
- The next lower sublayer can set maximum length limitations for the data units it accepts from the segmentation sublayer.
- When the segmentation function is in use, it is also assumed that the next lower sublayer can provide a service which ensures that the data units are delivered to the receiving end in sequence and without gaps or duplication.

5.2 TC Segment

5.2.1 General

- a. The length of the TC Segment shall not exceed the maximum length accepted by the next lower sublayer below the segmentation sublayer.

NOTE 1 The TC Segment is a variable-length data structure.

NOTE 2 The maximum length accepted by the transfer sublayer can vary across different virtual channels.

- b. The length of a TC Segment shall be an integer number of octets.
- c. A TC Segment shall consist of two fields, positioned contiguously, in the following sequence:

1. Segment Header 8 bits
2. Segment Data Field variable length

NOTE The format of the TC Segment is shown in Figure 5-1.

- d. If the next lower sublayer below the segmentation sublayer is the transfer sublayer, and the segmentation sublayer at the sending end is generating TC Segments for a particular virtual channel, then all data units delivered by the transfer sublayer at the receiving end for that virtual channel shall be TC Segments.

NOTE TC Segments do not share a virtual channel with other data structures.

- e. If the next lower sublayer below the segmentation sublayer is the transfer sublayer, then each TC Segment shall be placed in the Transfer Frame Data Field of a TC Transfer Frame such that the Transfer Frame Data Field contains exactly one complete TC Segment.

NOTE In this case, the length of the Transfer Frame Data Field is equal to the length of the TC Segment.

SEGMENT HEADER		SEGMENT DATA FIELD
SEQUENCE FLAGS	MAP IDENTIFIER	
2 bits	6 bits	variable

Figure 5-1: TC Segment

5.2.2 Segment Header

5.2.2.1 General

- a. The Segment Header shall always be present in a TC Segment.
- b. The Segment Header shall be contained in bits 0-7 of the TC Segment.
- c. The Segment Header shall consist of two fields, positioned contiguously, in the following sequence:
 1. Sequence Flags 2 bits
 2. MAP Identifier 6 bits

NOTE Both fields are always present in a Segment Header.

5.2.2.2 Sequence Flags

- a. The Sequence Flags shall always be present in a Segment Header.
- b. The Sequence Flags shall be contained in bits 0-1 of the Segment Header.
- c. The Sequence Flags shall be set as follows:
 1. '01' for the first segment of a user data unit;
 2. '00' for a continuing segment of a user data unit;
 3. '10' for the last segment of a user data unit;
 4. '11' no segmentation.

NOTE The Sequence Flags provide information about the sequential position of the segment to enable the receiving end to reassemble the user data units from a stream of segments. However, the flags do not provide a sequence count, so the reassembly process depends on the stream of segments being delivered in sequence and without omissions. The sequence-controlled service of the transfer sublayer can provide delivery in sequence and without omissions.

5.2.2.3 MAP Identifier

- a. The MAP Identifier shall always be present in a Segment Header.
- b. The MAP Identifier shall be contained in bits 2-7 of the Segment Header.

NOTE The six-bit MAP Identifier enables up to 64 MAPs, from MAP 0 to MAP 63, to be associated with each virtual channel of the transfer sublayer.

- c. The MAP Identifier shall provide the identification of the MAP to which the TC Segment belongs.

NOTE 1 There are no restrictions on the selection of MAP Identifiers. The MAPs need not be

numbered consecutively. However, if the packet assembly controller defined in clause 5.5.4 is in use, then MAPs are handled in pairs and the additional requirements specified in clause 5.5.4.2 apply.

NOTE 2 When a user data unit is placed in a sequence of TC Segments, then all the TC Segments for the user data unit have the same value for the MAP Identifier.

- d. If MAP multiplexing is not used on a particular virtual channel of the transfer sublayer, the MAP Identifier shall have a constant value for all TC Segments on that virtual channel.

NOTE In this case the virtual channel has a single MAP.

5.2.3 Segment Data Field

- a. The Segment Data Field shall always be present in a TC Segment, except as specified in 5.2.3e.
- b. The Segment Data Field shall start at bit 8 of the TC Segment.
- c. The maximum length of the Segment Data Field is one octet less than the maximum length of a TC Segment.

NOTE The length of the Segment Data Field is variable.

- d. The length of the Segment Data Field shall be an integer number of octets.
- e. If the packet assembly controller defined in clause 5.5.4 is in use, then the Segment Data Field may be absent from some TC Segments.

NOTE When the packet assembly controller is in use, some MAPs can carry control segments. A control segment is a special case of a TC Segment that has no Segment Data Field.

5.3 Transfer notification

5.3.1 Overview

The sequence-controlled service of the COP-1 protocol of the transfer sublayer uses acknowledgements to obtain positive confirmation that data units have arrived at the transfer layer at the receiving end. The data units are the frame data units carried in type-A TC Transfer Frames.

In the transfer notification (AD service) signal, COP-1 uses the Positive Confirm and Negative Confirm responses to provide information about the delivery status of a data unit. The transfer notification signal is defined in clause 7.4.3.3, along with explanations of the meaning of the confirm responses.

The segmentation sublayer at the sending end takes the information obtained from the confirm responses and makes it available to the next higher layer or sublayer. If a user data unit is carried in multiple segments, the segmentation sublayer combines the confirm responses for the individual segments.

The mechanism by which the segmentation sublayer makes the information available to the next higher layer of sublayer is outside the scope of this Standard.

The handling by the segmentation sublayer at the sending end of other responses from the transfer notification (AD service) signal, or of information provided by other COP-1 signals, is outside the scope of this Standard.

5.3.2 Requirements

- a. If a user data unit of the segmentation sublayer is transferred using the sequence-controlled service of the transfer sublayer, then the segmentation sublayer at the sending end shall provide, to the next higher layer or sublayer, the information about the delivery of the user data unit specified in clauses 5.3.2b to 5.3.2d.
- b. If a user data unit is not divided into multiple segments, then the segmentation sublayer shall provide a Positive Confirm or Negative Confirm to the next higher layer or sublayer, according to the response contained in the COP-1 signal relating to the transfer of the user data unit.
- c. If a user data unit is divided into multiple segments, then the segmentation sublayer shall provide a Positive Confirm to the next higher layer or sublayer if COP-1 has signalled Positive Confirm responses for all the segments.
- d. If a user data unit is divided into multiple segments, then the segmentation sublayer shall provide a Negative Confirm to the next higher layer or sublayer if COP-1 has signalled a Negative Confirm response for one or more of the segments.

NOTE If some segments of a user data unit have been signalled by a Positive Confirm but at least one segment has been signalled by a Negative Confirm, then the segmentation sublayer provides a Negative Confirm for the user data unit.

5.4 Blocking of packets

5.4.1 Overview

Blocking of packets is performed by a blocking function at the sending end and a deblocking function at the receiving end.

The blocking function blocks multiple packets into a single data unit. The function is constrained by the limitations on the length of the data unit. Figure 5-2 illustrates an example.

Additional constraints can apply to the way the segmentation sublayer uses the blocking function.

For example, an implementation can set a limit for the delay imposed on a packet while waiting for more packets to be blocked with it. The additional constraints and the mechanisms for ensuring transmission of all packets are implementation dependent and outside the scope of this Standard.

The deblocking function uses fields in the standard packet headers in order to extract the packets. The function therefore depends on knowledge of the position, size and meaning of certain fields in the headers.

Blocking can be used with or without TC Segments.

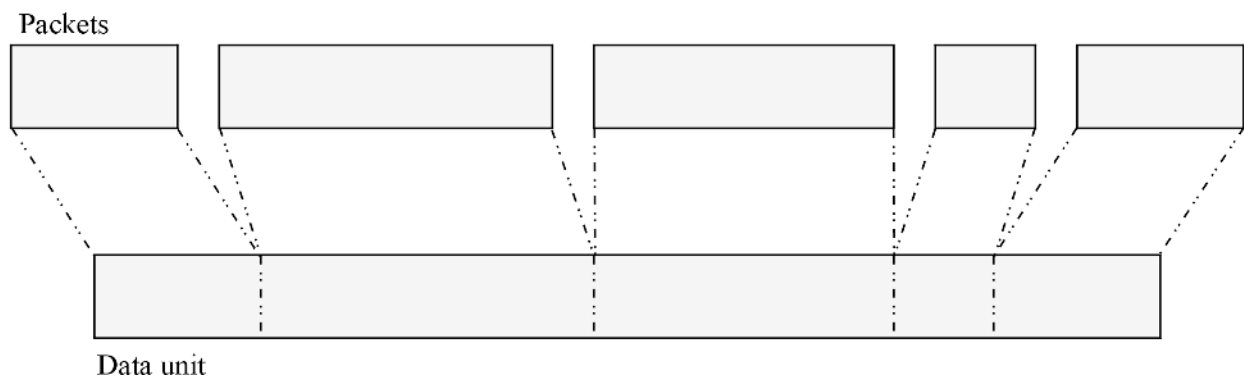


Figure 5-2: Example of blocking of packets

5.4.2 Virtual channels where TC Segments are used

- The blocking and deblocking functions shall be conducted independently for each MAP.
- The length of the data unit resulting from blocking shall not exceed the maximum length of the Segment Data Field of a TC Segment.
- A virtual channel may have some MAPs where the blocking of packets is enabled and some MAPs where the blocking of packets is not enabled.
- If the blocking of packets is enabled for a MAP, then all data units for that MAP, that are passed to the segmentation sublayer from the next higher layer or sublayer at the sending end, shall be packets with the properties defined in clause 5.4.4.

- e. If the blocking function blocks multiple packets into a data unit, then the Sequence Flags of the TC Segment which carries the data unit shall be set to indicate no segmentation.

NOTE A data unit containing multiple packets is not divided into multiple segments by the segmentation function.

5.4.3 Virtual channels where TC Segments are not used

- a. The blocking and deblocking functions shall be conducted independently for each virtual channel.
- b. The length of the data unit resulting from blocking shall not exceed the maximum length of the data units accepted by the next lower sublayer below the segmentation sublayer.
- c. A physical channel may have some virtual channels where the blocking of packets is enabled and some virtual channels where the blocking of packets is not enabled.
- d. If the blocking of packets is enabled for a virtual channel, then all data units for that virtual channel, which are passed to the segmentation sublayer from the next higher layer or sublayer at the sending end, shall be packets with the properties defined in clause 5.4.4.

5.4.4 Packet properties

- a. A packet handled by the blocking function shall have a defined Packet Version Number in conformance with CCSDS 135.0-B-3 clause 7.6 and with the definition of the related data structure specified in the same clause.

NOTE 1 The Packet Version Number occupies the first three bits of the packet header.

NOTE 2 CCSDS 135.0-B-3 clause 7.6 contains a list of the defined Packet Version Numbers, which are also referred to as authorized Packet Version Numbers. It provides also references to the documents that specify the related packet data structures.

NOTE 3 For missions that use the blocking function, it is the responsibility of the mission designer to verify the availability of support by the telecommand transfer services for each defined Packet Version Number to be used by the mission.

5.4.5 Blocking function

- a. Multiple complete packets may be placed contiguously into a data unit.

NOTE 1 The blocking function is only applied to complete packets.

NOTE 2 Due to other constraints, such as implementation-dependent time limits, the segmentation sublayer can choose not to block some packets.

- b. The order of the packets shall not be changed to improve the blocking of the packets into the data units.

NOTE The order of arrival of the packets is preserved.

- c. Packets with different Packet Version Numbers may be blocked into the same data unit.

- d. If the blocking function processes packets for the sequence-controlled service and packets for the expedited service, then all packets that are blocked into the same data unit shall be for the same service.

NOTE The transfer sublayer supports the sequence-controlled service and the expedited service.

5.4.6 Deblocking function

- a. The deblocking function shall extract the packets from the data units it receives.

- b. To perform the extraction, the deblocking function shall use the packet length fields, together with the length of the received data unit.

NOTE The deblocking function inspects the Packet Version Number of each packet in order to locate and interpret the packet length field.

5.5 Segmentation

5.5.1 Overview

A segmenting function at the sending end breaks long user data units into multiple portions which are placed in a sequence of TC Segments. At the receiving end, a reassembly function reassembles the user data units. The functions are conducted independently for each MAP.

All TC Segments belonging to a specific user data unit have the same value for the MAP Identifier and use the same service (sequence-controlled service or expedited service) in the transfer sublayer.

Successful operation of the segmentation function relies on the stream of TC Segments being delivered in sequence and without omissions. The sequence-controlled service can provide delivery in sequence and without omissions.

Figure 5-3 illustrates an example of segmentation of a user data unit.

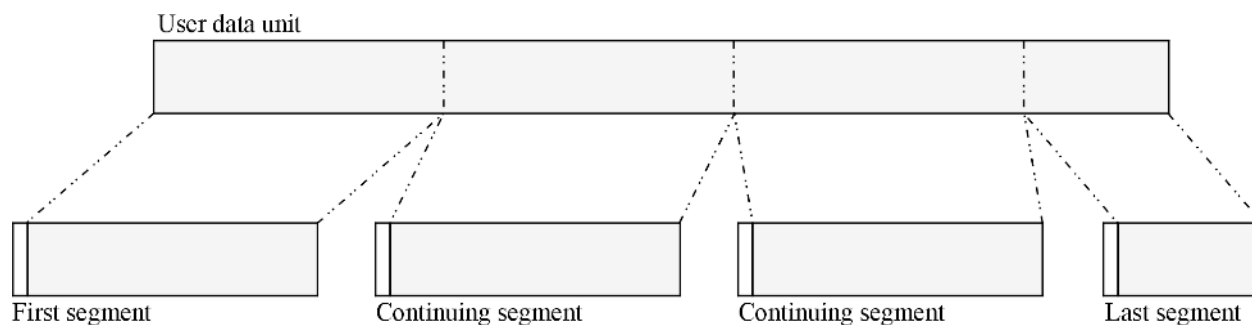


Figure 5-3: Example of segmentation of a user data unit

5.5.2 Segmenting function

- a. The segmenting function shall be conducted independently for each MAP.
- b. If a user data unit exceeds the maximum length of a Segment Data Field, the segmenting function shall split it as follows:
 1. For each portion, select the length of the portion such that it is an integer number of octets that does not exceed the maximum length of a Segment Data Field.
 2. Place the portions in order into a sequence of TC Segments.
 3. Place the first octet of the user data unit in the first octet of the Segment Data Field of the first TC Segment.
 4. Set the Sequence Flags of each TC Segment in the sequence to indicate whether it is a first, continuing or last segment.

NOTE 1 The Segment Data Field in the first and continuing TC Segments can have a length equal to the maximum length of the Segment Data Field. The Segment Data Field in the last TC Segment has a length equal to the residue of the user data unit.

NOTE 2 If a user data unit does not exceed the maximum length of a Segment Data Field, the complete user data unit is placed in a single TC Segment. In this case the Sequence Flags are set to indicate no segmentation.

5.5.3 Reassembly function

- a. The reassembly function shall be conducted independently for each MAP.
- b. The reassembly function shall reassemble the Segment Data Fields from a sequence of TC Segments to recreate the original user data unit.
- c. To perform the reassembly, the reassembly function shall use the Sequence Flags together with the lengths of the TC Segments.

NOTE 1 If a failure occurs during reassembly, the actions are undefined. The choice of whether

to deliver the incomplete or faulty user data unit to the next higher layer or sublayer is implementation dependent.

NOTE 2 The optional packet assembly controller defined in clause 5.5.4 specifies additional actions for the reassembly function, including the behaviour in the event of a reassembly failure.

5.5.4 Packet assembly controller

5.5.4.1 Overview

The packet assembly controller (PAC) can form part of the reassembly function at the receiving end of the segmentation function.

The reassembly function reassembles the segments for each MAP connection, in order to recreate the original user data units. The packet assembly controller carries out the reassembly, including the handling of exceptions.

When the PAC detects an exception it enters a lockout state. In the lockout state, the PAC does not reassemble user data units nor pass user data units to the next higher layer or sublayer. When it receives a MAP reset command, the PAC exits lockout state.

Despite the word “packet” in its name, the actions of the PAC apply to all user data units carried in TC Segments. The actions are not limited to any specific data structures. The name “packet assembly controller” is inherited from earlier standards.

If the segmentation function is not in operation, then the PAC is not used.

5.5.4.2 PAC for a pair of MAPs

- a. There shall be one PAC for each pair of MAP connections.
- b. A pair of MAP connections shall consist of one MAP for data and one MAP for control.
- c. The MAP Identifier for the data MAP shall have the most significant bit set to ‘0’.
- d. The MAP Identifier for the control MAP shall have the same value as the data MAP except that the most significant bit shall be set to ‘1’.
- e. A TC Segment carrying user data shall have the MAP Identifier of the data MAP.
- f. A TC Segment which is a PAC control segment shall have the MAP Identifier of the control MAP.

NOTE 1 The MAP Identifier is a 6-bit value, carried in the Segment Header. Therefore, for a pair of MAPs:

- the data MAP has an identifier in the range 0 to 31;

- the control MAP has an identifier in the range 32 to 63;
- control MAP Identifier = data MAP Identifier + 32;
- the least significant 5 bits of the two MAP Identifiers are the same.

NOTE 2 The PAC requirements on the use of the MAP Identifier apply in addition to the MAP Identifier specification in clause 5.2.2.3.

5.5.4.3 Reassembly of user data units

- a. The PAC reassembly function shall use the Sequence Flags in the Segment Header of each TC Segment to reassemble the user data units.

NOTE 1 The reassembly function implies that the number of octets in each segment is known by the PAC.

NOTE 2 The PAC reassembly function does not use any knowledge of the structure of the user data units.

For example, it does not make use of any length fields contained in the user data units. If the user data units have length fields, then verification that the length of the reassembled data unit is consistent with the value in the length field is outside the scope of this Standard.

5.5.4.4 PAC state

- a. The PAC shall enter a lockout state when it detects one of the following errors:
1. an incorrect sequence of data segments, as indicated by the Sequence Flags;
 2. a control segment with an invalid format.
- b. When the PAC is in a lockout state, it shall not reassemble user data units nor pass user data units to the next higher layer or sublayer.
- c. When the PAC is in a lockout state, it shall remain in that state until it receives a valid control segment as specified in clause 5.5.4.6.

NOTE 1 The following is a list of the incorrect sequences of segments that causes the PAC to enter lockout state. The values for the Sequence Flags are shown in parentheses.

- A first segment ('01') followed by a first segment ('01').
- A first segment ('01') followed by a no segmentation ('11').

- A continuing segment ('00') followed by a first segment ('01').
- A continuing segment ('00') followed by a no segmentation ('11').
- A last segment ('10') followed by a continuing segment ('00').
- A last segment ('10') followed by a last segment ('10').
- A no segmentation ('11') followed by a continuing segment ('00').
- A no segmentation ('11') followed by a last segment ('10').

NOTE 2 The lockout state of the PAC is unrelated to the FARM-1 Lockout state defined in clause 7.2.3.2.

5.5.4.5 PAC status report

- a. The PAC shall have a reassembly status flag set as follows:
 1. '0' when the PAC has completed reassembly of a user data unit;
 2. '1' when reassembly of a user data unit is in progress in the PAC.
- b. The PAC shall have a lockout status flag set as follows:
 1. '1' when the PAC is in a lockout state;
 2. '0' when the PAC is not in a lockout state.
- c. The PAC shall report its status to the sending end, including the following:
 1. the MAP Identifier of the data MAP, and
 2. the reassembly status flag, and
 3. the lockout status flag.

NOTE The correct operation of the PAC relies on the status of the PAC being known by the sending end. This Standard does not specify the format of the status information nor the mechanism to be used to transport it from the PAC to the appropriate entity at the sending end. It also does not specify any resulting behaviour at the sending end, such as the decision to send a control segment.

5.5.4.6 Control segment

- a. A control segment shall have a length of one octet.

NOTE A control segment is a special case of a TC Segment. It has no Segment Data Field and therefore consists of a Segment Header only.

- b. The Sequence Flags in the Segment Header of a control segment shall be set to '11'.
- c. The MAP Identifier in the Segment Header of a control segment shall contain the MAP Identifier of a control MAP.
- d. A valid control segment shall be considered to be a MAP reset command.

NOTE 1 If the PAC receives a segment that has the MAP Identifier of a control MAP but the segment does not conform to the format rules the PAC enters lockout state.

NOTE 2 This Standard does not specify any type of service (sequence-controlled service or expedited service) to be used for the transfer of a control segment.

5.5.4.7 MAP reset actions

- a. When the PAC receives a MAP reset command and the PAC reassembly status flag is '1', the PAC shall:
 - 1. discard the partially reassembled data unit, and
 - 2. set the reassembly status flag to '0'.
- b. When the PAC receives a MAP reset command and the PAC is in lockout state, the PAC shall exit lockout state.

NOTE The MAP reset command is used, for example, to recover from breaks in the sequence of TC Segments due to link difficulties or unplanned termination of transfer services.

5.6 MAP multiplexing

User data units allocated to different MAPs are multiplexed together so that they can share a virtual channel of the transfer sublayer. Each virtual channel can have up to 64 distinct MAPs (from MAP 0 to MAP 63) associated with it.

The MAP multiplexer selects a TC Segment depending on the priorities of the MAPs and the algorithm in use. This Standard does not specify any multiplexing algorithms. The type of priority scheme employed and the allocation of priorities to the individual MAPs is implementation dependent.

Although the multiplexing algorithms are implementation dependent, cross-support missions can specify the use of algorithms described in CCSDS 912.3-B-1.

If MAP multiplexing is in use, the order of arrival of user data units is not preserved across MAPs.

For example, a user data unit on a high-priority MAP can be processed before an earlier data unit on a low-priority MAP. Also, segments from user data units on different MAPs can be interleaved.

6

Transfer sublayer

6.1 Overview

6.1.1 Data structures in the transfer sublayer

The principal data structure in the transfer sublayer is the TC Transfer Frame, specified in clause 6.2. At the sending end, the transfer sublayer delivers TC Transfer Frames to the next lower sublayer.

The transfer sublayer accepts variable-length data units from the next higher sublayer at the sending end and delivers them to the next higher sublayer at the receiving end. Within the transfer sublayer, these data units are referred to as frame data units (FDUs) because each one is carried in the data field of a TC Transfer Frame. The transfer sublayer imposes constraints on the lengths of the data units.

The transfer sublayer at the receiving end accepts variable-length blocks containing an integral number of octets from the next lower sublayer. The transfer sublayer processes each received block as a candidate frame, to verify that it contains a valid TC Transfer Frame.

The operation of the transfer sublayer protocol uses status information which is passed from the receiving end to the sending end of the sublayer. The information is contained in a communications link control word (CLCW), which can be transmitted in a telemetry transfer frame. The format of the CLCW is defined in clause 6.3, and the processing for CLCWs in Clause 7.

6.1.2 Procedures in the transfer sublayer

6.1.2.1 Communications operation procedure

The transfer sublayer includes the communications operation procedure (COP), which operates the transmission protocol of the sublayer. A COP consists of a pair of synchronized procedures which execute within one virtual channel at the sending and receiving ends of the sublayer. This Standard defines one COP, called COP-1, whose detailed specification is contained in Clause 7. COP-1 supports two service types for data transfer: the sequence-controlled service and the expedited service.

6.1.2.2 Sending end

Table 6-1 shows the procedures in the transfer sublayer at the sending end.

At the sending end, the first procedure in the transfer sublayer is the frame operation procedure (FOP-1), which is the sending end of COP-1.

The COP-1 procedures operate within a virtual channel so the transfer sublayer has an instance of FOP-1 for each virtual channel.

There are two procedures to complete the fields in each Transfer Frame:

- the frame header procedure places values in the fields of the Transfer Frame Primary Header, and
- the frame error control procedure then generates the frame check sequence and places it in the Frame Error Control Field.

The transfer sublayer includes virtual channel multiplexing. Because FOP-1 has an instance for each virtual channel, the virtual channel multiplexing is after FOP-1, but its position relative to the other two procedures is implementation dependent.

Table 6-1: Sending-end procedures in the transfer sublayer

Procedure	Clause	Position in the sequence of execution	(Peer procedure at the receiving end)
FOP-1	7	First procedure in the transfer sublayer	(FARM-1)
Frame header procedure	6.4	- After FOP-1 - Before frame error control procedure	(Frame header validation procedure)
Frame error control procedure	6.5	After frame header procedure	(Frame error control procedure)
Virtual channel multiplexing	6.9	After FOP-1	(Virtual channel demultiplexing)

6.1.2.3 Receiving end

Table 6-2 shows the procedures in the transfer sublayer at the receiving end.

At the receiving end, the first procedure is the frame delimiting and fill removal procedure. This validates the length and removes any fill data from the end of each candidate frame.

The frame error control procedure verifies the frame check sequence in the Frame Error Control Field. The frame header validation procedure checks the remaining fields of the Transfer Frame Primary Header.

The last procedure is the frame acceptance and reporting mechanism (FARM-1), which is the receiving end of COP-1. The COP-1 procedures operate within a

virtual channel so the transfer sublayer has an instance of FARM-1 for each virtual channel.

The transfer sublayer includes virtual channel demultiplexing, which is after the frame delimiting and fill removal procedure. Because FARM-1 has an instance for each virtual channel, the virtual channel demultiplexing is before FARM-1, but its position relative to the other two procedures is implementation dependent.

Table 6-2: Receiving-end procedures in the transfer sublayer

Procedure	Clause	Position in the sequence of execution	(Peer procedure at the sending end)
Frame delimiting and fill removal procedure	6.6	First procedure in the transfer sublayer	(none)
Frame error control procedure	6.7	- After frame delimiting and fill removal procedure - Before frame header validation procedure	(Frame error control procedure)
Frame header validation procedure	6.8	- After frame error control procedure - Before FARM-1	(Frame header procedure)
Virtual channel demultiplexing	6.9	- After frame delimiting and fill removal procedure - Before FARM-1	(Virtual channel multiplexing)
FARM-1	7	Last procedure in the transfer sublayer	(FOP-1)

6.2 TC Transfer Frame definition

6.2.1 General

- a. The TC Transfer Frame shall consist of the major fields, positioned contiguously, in the following sequence:

1. Transfer Frame Primary Header 5 octets
2. Transfer Frame Data Field variable length
3. Frame Error Control Field 2 octets

NOTE 1 All three fields are always present in a TC Transfer Frame.

NOTE 2 Figure 6-1 illustrates the detailed format of the TC Transfer Frame.

- b. The length of a TC Transfer Frame shall be an integer number of octets.

NOTE 1 The TC Transfer Frame is a variable-length data structure.

NOTE 2 The length of a TC Transfer Frame is indicated by the Frame Length field defined in clause 6.2.2.8. Because of constraints imposed by the Frame Length field, the length of a TC Transfer Frame does not exceed 1024 octets. For a mission, a value of less than 1024 octets can be selected for the length of the longest TC Transfer Frame that can be carried on a particular virtual channel.

NOTE 3 The length of the Transfer Frame Data Field defined in clause 6.2.3 is not less than 1 octet. As a result, a TC Transfer Frame has a length which is not less than 8 octets.

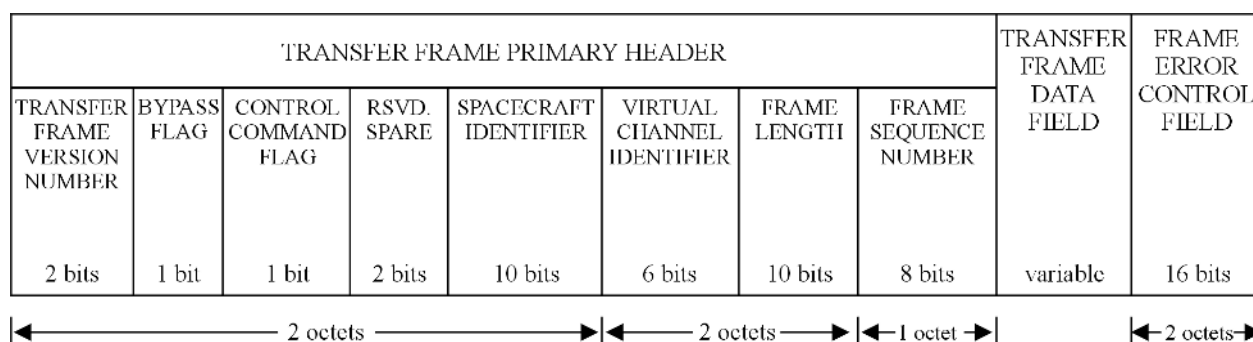


Figure 6-1: TC Transfer Frame format

6.2.2 Transfer Frame Primary Header

6.2.2.1 General

- a. The Transfer Frame Primary Header shall always be present in a TC Transfer Frame.
- b. The Transfer Frame Primary Header shall consist of eight fields, positioned contiguously, in the following sequence:
- Transfer Frame Version Number 2 bits
 - Bypass Flag 1 bit
 - Control Command Flag 1 bit
 - Reserved Spare 2 bits
 - Spacecraft Identifier 10 bits
 - Virtual Channel Identifier 6 bits

- | | | |
|----|-----------------------|---------|
| 7. | Frame Length | 10 bits |
| 8. | Frame Sequence Number | 8 bits |

NOTE All eight fields are always present in a Transfer Frame Primary Header.

6.2.2.2 Transfer Frame Version Number

- a. The Transfer Frame Version Number shall always be present in a Transfer Frame Primary Header.
- b. The Transfer Frame Version Number shall be contained in bits 0-1 of the Transfer Frame Primary Header.
- c. The Transfer Frame Version Number shall be set to the binary value '00'.

NOTE The value '00' defines version 1 of the TC Transfer Frame. Other values of this field are reserved.

6.2.2.3 Bypass Flag

- a. The Bypass Flag shall always be present in a Transfer Frame Primary Header.
- b. The Bypass Flag shall be contained in bit 2 of the Transfer Frame Primary Header.
- c. The Bypass Flag shall be set as follows:
 - 1. '0' for a type-A Transfer Frame;
 - 2. '1' for a type-B Transfer Frame.
- d. The Bypass Flag shall be interpreted in combination with the Control Command Flag as defined in Table 6-3.

NOTE 1 When the Bypass Flag is set to '1', the TC Transfer Frame bypasses the sequence control checks at the receiving end of the COP-1 procedures.

NOTE 2 To operate the sequence-controlled service, a virtual channel carries type-A and type-B Transfer Frames.

6.2.2.4 Control Command Flag

- a. The Control Command Flag shall always be present in a Transfer Frame Primary Header.
- b. The Control Command Flag shall be contained in bit 3 of the Transfer Frame Primary Header.
- c. The Control Command Flag shall be set as follows:
 - 1. '0' for a type-D Transfer Frame, carrying data;
 - 2. '1' for a type-C Transfer Frame, carrying a control command.

- d. The Control Command Flag shall be interpreted in combination with the Bypass Flag as defined in Table 6-3.

NOTE A control command is a protocol control instruction used in COP-1 to configure the procedures at the receiving end.

Table 6-3: Combined Bypass Flag and Control Command Flag

Bypass Flag	Control Command Flag	Interpretation
0	0	AD frame, carrying data in the sequence-controlled service
0	1	Reserved for future application
1	0	BD frame, carrying data in the expedited service
1	1	BC frame, carrying a control command

6.2.2.5 Reserved Spare

- a. The spare field contained in bits 4-5 of the Transfer Frame Primary Header shall be set to '00'.

NOTE The use of this spare field is reserved for future applications.

6.2.2.6 Spacecraft Identifier

- a. The Spacecraft Identifier shall always be present in a Transfer Frame Primary Header.
- b. The Spacecraft Identifier shall be contained in bits 6-15 of the Transfer Frame Primary Header.
- c. The Spacecraft Identifier shall provide the identification of the destination spacecraft for the TC Transfer Frame.
- d. The Spacecraft Identifier shall be static throughout all mission phases.

NOTE 1 The Secretariat of the CCSDS assigns Spacecraft Identifiers according to the procedures in CCSDS 320.0-B-4.

NOTE 2 Different Spacecraft Identifiers can be assigned for normal operations and for development vehicles using the ground networks during pre-launch test operations, and for simulated data streams.

NOTE 3 A single physical channel can carry virtual channels with different Spacecraft Identifiers. Therefore the Spacecraft Identifier forms part of the full identification of a virtual channel within the physical channel.

6.2.2.7 Virtual Channel Identifier

- a. The Virtual Channel Identifier shall always be present in a Transfer Frame Primary Header.
- b. The Virtual Channel Identifier shall be contained within bits 16-21 of the Transfer Frame Primary Header.
- c. The Virtual Channel Identifier, together with the Spacecraft Identifier, shall identify the virtual channel to which the Transfer Frame belongs.

NOTE 1 The Virtual Channel Identifier is also used in the CLCW defined in clause 6.3. As described in clause 6.3.6, the choice of values for the identifiers of the virtual channels can affect the risk of misidentifying the virtual channel for a CLCW.

NOTE 2 The Spacecraft Identifier forms part of the identification of a virtual channel within a physical channel.

6.2.2.8 Frame Length

- a. The Frame Length shall always be present in a Transfer Frame Primary Header.
- b. The Frame Length shall be contained in bits 22-31 of the Transfer Frame Primary Header.
- c. The Frame Length shall contain a binary number equal to the number of octets in the TC Transfer Frame minus 1.

NOTE The value contained in the Frame Length is in the range 7 to 1023, corresponding to TC Transfer Frame lengths of 8 to 1024 octets.

- d. The value contained in the Frame Length may be variable.

NOTE Earlier definitions limited the TC Transfer Frame length to a maximum of 256 octets. In the earlier definitions, the Frame Length is contained in bits 24-31 of the Transfer Frame Primary Header, and is preceded by a 2-bit reserved field containing zeroes. Therefore, frames that conform to the earlier definition are valid under the current definition.

6.2.2.9 Frame Sequence Number

- a. The Frame Sequence Number shall always be present in a Transfer Frame Primary Header.
- b. The Frame Sequence Number shall be contained in bits 32-39 of the Transfer Frame Primary Header.
- c. In type-A Transfer Frames the Frame Sequence Number shall contain the COP-1 value N(S).

NOTE 1 Clause 7.9.2.6 specifies the setting of the COP-1 value N(S).

NOTE 2 The value is used in the COP-1 procedures at the receiving end to check that incoming type-A Transfer Frames are in sequence.

- d. In type-B Transfer Frames the Frame Sequence Number shall be set to zero.

NOTE The Bypass Flag defined in clause 6.2.2.3 indicates if a TC Transfer Frame is type-A or type-B.

6.2.3 Transfer Frame Data Field

- a. The Transfer Frame Data Field shall always be present in a TC Transfer Frame.
- b. The Transfer Frame Data Field shall follow, without gap, the Transfer Frame Primary Header.
- c. The length of the Transfer Frame Data Field shall be an integral number of octets.

NOTE 1 The length of the Transfer Frame Data Field is constrained by the length of the total TC Transfer Frame. The constraints on the length of a TC Transfer Frame are described in clause 6.2.1.

NOTE 2 The length of the Transfer Frame Data Field is not less than 1 octet and does not exceed 1017 octets.

- d. In type-D Transfer Frames the Transfer Frame Data Field shall contain a frame data unit.

NOTE The Transfer Frame Data Field of a type-D Transfer Frame contains exactly one frame data unit. The length of a frame data unit is therefore constrained to be a valid length for the field.

- e. In type-C Transfer Frames the Transfer Frame Data Field shall contain a COP-1 control command.

NOTE The Control Command Flag defined in clause 6.2.2.3 indicates if a TC Transfer Frame is type-D or type-C. Clause 6.1.1 describes the frame data unit. The COP-1 control commands are specified in clause 7.8.

6.2.4 Frame Error Control Field

6.2.4.1 General

- a. The Frame Error Control Field shall always be present in a TC Transfer Frame.
- b. The Frame Error Control Field shall follow, without gap, the Transfer Frame Data Field.
- c. The Frame Error Control Field shall have a length of two octets.
- d. The error detection encoding and decoding procedures specified in clauses 6.2.4.2 and 6.2.4.3 shall be used.

NOTE 1 This field provides a capability for detecting errors introduced into the frame during the transmission and data handling process. One reason for including this field in all TC Transfer Frames is to protect against residual errors after codeblocks are decoded in error-correcting mode as specified in clause 8.8.2.

NOTE 2 When applied to an encoded block of less than 32 768 bits (4096 octets), the code has the following capabilities:

- All error sequences composed of an odd number of bit errors are detected.
- All error sequences containing two-bit errors anywhere in the encoded block are detected.
- If a random error sequence containing an even number of bit errors (≥ 4) occurs within the block, the probability that the error is undetected is approximately 2^{-15} (or 3×10^{-5}).
- All single error bursts spanning 16 bits or less are detected provided no other errors occur within the block.

6.2.4.2 Encoding Procedure

- a. The encoding procedure shall be as follows:
 1. The encoding procedure accepts an $(n-16)$ -bit TC Transfer Frame, excluding the Frame Error Control Field, and generates a systematic binary $(n, n-16)$ block code by appending a 16-bit Frame Error Control Field as the final 16 bits of the codeblock.
 2. The equation for the contents of the Frame Error Control Field is:

$$\text{FECF} = [(X^{16} \cdot M(X)) + (X^{(n-16)} \cdot L(X))] \text{ modulo } G(X)$$

where

- all arithmetic is modulo 2;
- FECF is the 16-bit Frame Error Control Field and the first bit transferred is the most significant bit, taken as the coefficient of the highest power of X ;
- n is the number of bits in the encoded message;
- $M(X)$ is the $(n-16)$ -bit information message to be encoded expressed as a polynomial with binary coefficients, and the first bit transferred is the most significant bit M_0 taken as the coefficient of the highest power of X ;
- $L(X)$ is the presetting polynomial given by:
- $L(X) = \text{Error!}$
- $G(X)$ is the generating polynomial given by:
- $G(X) = X^{16} + X^{12} + X^5 + 1$

NOTE 1 The $X^{(n-16)} \cdot L(X)$ term has the effect of presetting the shift register to all '1' state prior to encoding.

NOTE 2 The encoding has the characteristics of a cyclic redundancy code.

NOTE 3 A sample implementation of an encoder is described in Annex A.

6.2.4.3 Decoding Procedure

- a. The decoding procedure shall use an error detection syndrome, $S(X)$, given by:

$$S(X) = [(X^{16} \cdot C^*(X)) + (X^n \cdot L(X))] \text{ modulo } G(X)$$

where

- $C^*(X)$ is the received block, including the Frame Error Control Field, in polynomial form, with the first bit transferred being the most significant bit C_0 taken as the coefficient of the highest power of X , and
- $S(X)$ is the syndrome polynomial, which is zero if no error is detected and non-zero if an error is detected, with the most significant bit S_0 taken as the coefficient of the highest power of X .

NOTE A sample implementation of a decoder is described in Annex A.

- b. The Frame Error Control Field shall not be used for error correction.

NOTE The code is intended for error detection purposes only.

6.3 CLCW definition

6.3.1 General

- a. The length of a CLCW shall be 4 octets.
- b. A CLCW shall consist of the fields shown in Table 6-4.

NOTE The table defines the position and length of each field. All the fields are always present in a CLCW. The format of a CLCW is shown in Figure 6-2.

Table 6-4: Fields in a CLCW

Field	Location (bits)	Length in bits
Control Word Type	0	1
CLCW Version Number	1-2	2
Status Field	3-5	3
COP in Effect	6-7	2
Virtual Channel Identification	8-13	6
Reserved Spare	14-15	2
No RF Available Flag	16	1
No Bit Lock Flag	17	1
Lockout Flag	18	1
Wait Flag	19	1
Retransmit Flag	20	1
FARM B Counter	21-22	2
Reserved Spare	23	1
Report Value	24-31	8

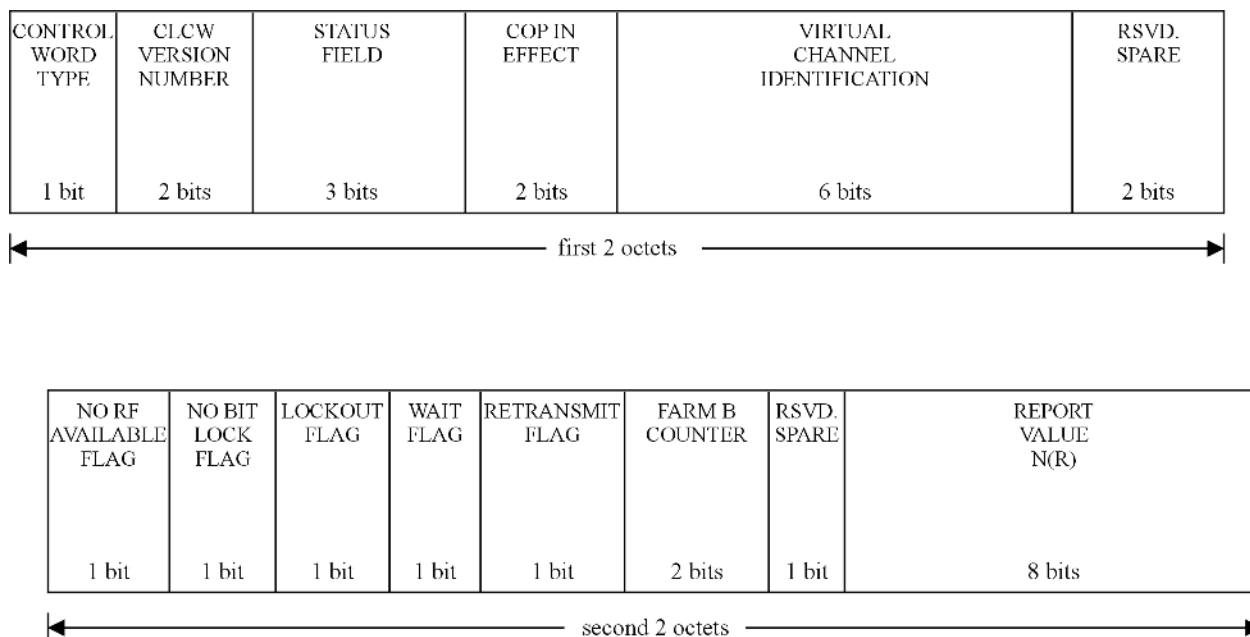


Figure 6-2: Format of a CLCW

- c. A CLCW shall be conveyed in the Operational Control Field of a telemetry transfer frame.

NOTE The telemetry transfer frame can be a version 1 telemetry transfer frame as defined in ECSS-E-ST-50-03 or a version 2 telemetry transfer frame as defined in CCSDS 732.0-B-2.

- d. A CLCW shall only be delivered to the FOP-1 if it has been extracted from an error-free telemetry transfer frame.

NOTE If any errors in the frame are successfully corrected, the frame is considered error-free.

- e. A CLCW shall only be delivered to the FOP-1 of a virtual channel if it has been extracted from a telemetry transfer frame from the spacecraft to which that telecommand virtual channel belongs.

NOTE 1 The purpose of a CLCW is to carry COP-1 protocol status information from the FARM-1 procedure at the receiving end to the FOP-1 procedure at the sending end.

NOTE 2 The COP-1 procedures operate within a virtual channel. CLCWs are generated for each active virtual channel.

NOTE 3 Although no reporting rate is specified, efficient operation of COP-1 can only be achieved if some minimum CLCW sampling rate is provided. The sampling rate is one of the factors used in selecting the FOP-1 timer value described in clause 7.2.2.9.

6.3.2 Control Word Type

- a. The Control Word Type shall be set to the binary value '0'.

NOTE The Control Word Type is used to distinguish between a CLCW and other uses of the Operational Control Field. For a CLCW, this field is always zero.

6.3.3 CLCW Version Number

- a. The CLCW Version Number shall be set to the binary value '00'.

NOTE The value zero indicates that the CLCW is a Version-1 CLCW. Other values of this field are reserved for future use.

6.3.4 Status Field

- a. The use of the Status Field may be defined for mission-specific enhancements to telecommand operations.
- b. If the use of the Status Field is not defined for a given mission, the Status Field shall be set to the binary value '000'.

6.3.5 COP in Effect

- a. The COP in Effect shall be set to the binary value '01'.

NOTE The value 1 indicates that COP-1 is in effect. Other values of this field are reserved for future use.

6.3.6 Virtual Channel Identification

- a. The Virtual Channel Identification shall identify the virtual channel to which the CLCW belongs.

NOTE 1 The spacecraft identifier of the spacecraft that generated the CLCW also forms part of the identification of the virtual channel. The spacecraft identifier is not carried in the CLCW but is established by other means. The spacecraft identifier and the Virtual Channel Identification together identify the instance of the FOP-1 procedure which is the destination of the CLCW.

NOTE 2 Telecommand virtual channel identifier values are assigned to maximize digital distance (i.e. hamming distance) between the values for the various virtual channels used by a spacecraft. This reduces the risk that an error introduced in a CLCW during on-board

data handling changes the virtual channel identifier value into the value of another virtual channel used by the spacecraft.

NOTE 3 A virtual channel identifier value consisting of all zeroes or all ones can be generated when a decoder fails or is powered off. Therefore virtual channel identifier values consisting of all zeroes or all ones are normally avoided.

NOTE 4 The CLCW can be conveyed in a telemetry transfer frame, that contains a virtual channel identifier in its header. The field in the header identifies the telemetry virtual channel which carried the CLCW and is not related to the telecommand virtual channel described here.

6.3.7 Reserved Spare

- a. The spare field contained within bits 14-15 of the CLCW shall be set to the binary value '00'.

NOTE The use of this spare field is reserved for future applications.

6.3.8 No RF Available Flag

- a. The No RF Available Flag shall be set as follows:
1. '0' when the radio frequency physical connection is available through at least one of the spacecraft transponders;
 2. '1' when the radio frequency physical connection is not available through any of the spacecraft transponders.

NOTE 1 The value of this flag provides a logical indication of the ready status of the radio frequency elements within the physical data channel provided by the physical layer at the receiving end.

NOTE 2 The presence of this flag in the CLCW fulfils an operational requirement. This does not preclude the transmission of telemetry data with more complete status information for monitoring purposes.

NOTE 3 When the flag is '1', TC Transfer Frames cannot be transferred without corrective action.

NOTE 4 This flag is used by the physical layer procedures as defined in clause 9.3.

- b. The No RF Available Flag shall always carry a true report of the status of the physical channel, even when other fields in the CLCW report the status of an inactive virtual channel.