

MINI-PROCESSOR
PRODUCTURES

**IEEE Standard for an
8-Bit Backplane
Interface: STEbus**

Includes ANSI/IEEE Std 1101-1987,
IEEE Standard for Mechanical Core
Specifications for Microcomputers



Published by
The Institute of Electrical and
Electronics Engineers, Inc

SH12112

An American National Standard

**IEEE Standard for an
8-Bit Backplane Interface: STEbus**



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**ANSI/IEEE
Std 1000-1987**

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**IEEE Standard for an
8-Bit Backplane Interface: STEbus**

Sponsor
**Technical Committee on Microprocessors and Microcomputers
of the
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Foreword

(This Foreword is not a part of ANSI/IEEE Std 1000-1987, IEEE Standard for an 8-Bit Backplane Interface: STEbus.)

The initial concept for STEbus was to produce a European version of the STD Bus using the Eurocard form factor with the DIN41612 connector. From that concept STE became known as STD-European.

When IEEE formed Working Group P1000 the brief specified a Standard 8-Bit Backplane Interface. At the inaugural meeting of Working Group P1000 it quickly became apparent that the opportunity was there to create a completely new, modern, high-performance 8-Bit bus, and all ideas of merely repinning the old STD Bus were rapidly forgotten.

At the initial meeting of P1000 it was decided that the bus should be a part of the same family as VMEbus and Futurebus and as such should be an asynchronous bus with multimaster capability. Today it is often referred to as the baby brother of VMEbus. Unlike VMEbus though it was to be processor and manufacturer independent. This has proven to be an excellent decision as today there are many varied types of processor available on STEbus, from microcontrollers such as 8031, through Intel's 8085, 8088, and 80188; National Semiconductor's 32008 and 32016; Motorola's 6809, 68000, and 68008; Zilog's Z80 and Z280; Hitachi's 64180, and the Inmos Transputer with the promise of more to come.

A presentation was made to a packed audience at the IEE in London, England in early 1983. It met with critical acclaim. The first article about STEbus was also published about this time in an international magazine (EDN May 26, 1983).

Work continued internationally and in late 1984 Draft D3.1 was produced. This draft eradicated the daisy-chain bus request mechanism of D2.0 in favor of a simple solution that allowed position independence of cards in the rack.

This was the first firm specification and encouraged more manufacturers to look at the bus seriously. Among them were BICC-Vero, a major manufacturer of Eurocard enclosures and backplanes, and British Telecom, the UK's Telephone Utility. Market ground zero was early 1985 and since this time the number of manufacturers has continued to grow from 18 companies in Spring 1986 to more than 30 in mid-1987, with over 700 products available.

Much credit and praise is due Tim Elsmore who first conceived the idea for STEbus during his employment with GMT Electronic Systems Ltd. Paul Borrill was instrumental in negotiating with IEEE the formation of Working Group P1000 and Bill Shields was appointed Chairman.

This standard was prepared by Working Group P1000 of the Microprocessor

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1. General

1.1 Scope. The overall level of performance that may be achieved by any computer system is determined, in large part, by the system bus that is used to effect communication between the various system elements. System performance characteristics, measured in terms of speed, reliability, suitability to a variety of purposes, and adaptability to changing technology are ultimately dependent on the particular bus structure that is used and its associated protocols.

This standard defines the IEEE Std 1000 Bus, which may be used to implement general purpose, high-performance 8-bit microcomputer systems. Such a system may be used in a stand-alone configuration, or in larger multiple-bus architectures, as a private (or secondary) bus or a high-speed I/O channel. This standard is applicable to those systems and system elements with the common commercial designation STE Bus. It is intended for those users who plan to evaluate, implement, or design various system elements that are compatible with the IEEE 1000 Std Bus system structure.

The physical attributes and method of interconnect utilized by boards and modules conforming to this standard are derived from several International Electrotechnical Commission (IEC) standards. These standards, when implemented jointly in a systems environment, result in a mechanical configuration commonly referred to as *Eurocard*. Appendix B lists such applicable standards which, where referenced, are considered as if incorporated with this standard. In particular the connector used by IEEE Std 1000 Bus boards is a 64-pin male connector utilizing the outside two rows (designated *a* and *c* rows), specified in IEC Publication 603-2,¹ and the mating female connector is used on IEEE Std 1000 Bus backplanes. The recommended size for IEEE Std 1000 Bus boards is 100 mm × 160 mm (3.937 in × 6.299 in), commonly referred to as a *single height standard depth* Eurocard.

The IEEE Std 1000 Bus structure is based on the master-slave concept in which a master, having gained control of the bus, may address and command slaves. Masters and slaves communicate with each other by use of an asynchronous interlocked handshake protocol. This technique allows for the construction of computer systems that incorporate devices of widely varying speeds. Multiple masters may be implemented within a single system.

¹ IEC publications are available from the Sales Department, American National Standards Institute, 1430 Broadway, New York, NY 10018.

Two independent address spaces are supported: memory and I/O. Memory transactions reference a 1 megabyte physical address space, while I/O transactions reference a 4 kilobyte physical address space. System integrity during all such transactions is enhanced by provision of a transfer error signal.

Provision is made for interboard condition alerts such as interrupt requests, DMA requests, system-specific error conditions, or other specialized status conditions. Within this scheme eight prioritized attention request levels, each with vectored response capability, are available for user assignment.

This standard deals only with those characteristics that must be specified so as to ensure the successful design and implementation of compatible boards and systems. Issues relating to individual design specifications, and performance or safety requirements are not addressed.

1.2 Features. The fundamental features offered by IEEE Std 1000 Bus are as follows:

- 8-bit Data Field Width
- 1 Megabyte Memory Address Range
- 4 Kilobyte I/O Address Range
- Asynchronous Data Transfer
- Transfer Error Signal
- Multiple Masters
- Eight Attention Request Lines
- IEC 603-2 Connector
- Single or Double Eurocard Boards and Modules
- 5 V, ± 12 V and Standby Power Supply Distribution
- Total Position Independence of Boards and Modules in Backplane
- Total Inter-Board Compatibility
- Total Central Processing Unit (CPU) Generic Device Family Independence
- Potential 5 Megabyte per Second Data Transfers

1.3 Objects. This standard is intended to

- (1) Define a general purpose microcomputer board interface.
- (2) Specify those device-independent electrical, mechanical, and functional interface parameters that must be met so as to effect unambiguous communication between system elements and to effect physical compatibility.
- (3) Specify the terminology and definitions related to the specification.
- (4) Enable the interconnection of a wide variety of independently manufactured boards within a single functional system.
- (5) Define a standard that places the minimum number of restrictions on the performance characteristics of boards within a conforming system.

(6) Allow microcomputer system users of relatively modest experience to assemble modularly expandable computer systems with a high probability of success.

1.4 Definitions. The following general definitions apply throughout this standard. Additional detailed definitions are given where appropriate.

1.4.1 General System Terms

compatibility. The degree to which boards may be interconnected and used without modification when designed according to the specifications contained within this standard.

interface. A shared boundary between two or more systems, or between two or more elements within a system, through which information is conveyed.

interface system. The device-independent electrical, mechanical, and functional interface elements required for unambiguous communication between two or more devices. Typical elements include

- driver and receiver circuitry
- signal line descriptions
- timing and control conventions
- communication protocols
- functional logic circuits.

system. A set of interconnected boards that achieve a specified objective by the performance of designated functions.

1.4.2 Signals and Paths

address. The reference to a unit of data or the value represented by the address lines while ADRSTB* is active.

addressed board. A board that recognizes its address while ADRSTB* is active.

arbitration. The means whereby masters compete for control of the bus and the process by which a master is granted control of the bus.

backplane. A printed circuit board (pcb) on which connectors are mounted, into which boards or plug-in units are inserted.

block transfer. A sequence of data transfers, in the same direction, that occur during a single bus transaction.

board. A printed circuit board (pcb) that complies with this standard.

bus. A signal line or set of lines used by an interface system to connect a number of devices, and over which information is conveyed.

byte. A set of eight signals, individually referred to as bits, which are operated on as a unit.

handshake. An interlocked sequence of signals between interconnected boards in which each board waits for an acknowledgement of its previous signal before proceeding.

high state. The more positive voltage level used to represent one of two logical binary states.

low state. The more negative voltage level used to represent one of two logical binary states.

module. A plug-in unit consisting of one or more boards that contains at least one bus interface conforming to this standard, which plugs into the backplane.

protocol. The signaling rules used to convey information or commands between boards connected to the bus.

release. The action of a transmitter in ceasing to hold a signal line in the asserted state.

sequence. An indivisible bus transaction comprising one or more transfers.

settling time. The time taken for a signal line to settle unambiguously to a logical state when making a transition from one state to another.

signal. The physical representation of data.

signal level. The relative magnitude of a signal when considered in relation to an arbitrary reference. The unit of representation used within this standard is the volt.

signal line. One of a set of signal conductors in an interface system used to transfer data among interconnected boards.

signal parameter. That element of an electrical quantity whose values or sequence of values convey information.

tenure. The time during which a master has control of the bus.

transaction. The combination of data transfer sequences controlled by a master during a single bus tenure.

transfer. The movement of a single byte of data from the current master to the addressed slave(s) or from the addressed slave to the master.

1.4.3 Generic Signal Names. Throughout this standard bus request and acknowledge signals and attention request signals are sometimes referred to as $BUSRQ_n^*$, $BUSAK_n^*$, and $ATNRQ_n^*$ respectively. Such general references are equivalent to specific references $RUSRQ0^*$ or $BUSRQ1^*$ etc.

1.4.4 Notation for Bus Signals. Throughout this standard signals on a particular bus are referred to collectively using the form $A<19..0>$. This notation should be taken as an abbreviation of all of the address bus signals from A19 through to A0 inclusive.

In addition to the address bus signals, the notation is also used for the data lines (for example, $D<7..0>$), the common lines (for example, $CM<2..0>$), the attention request lines (for example, $ATNRQ<7..0>^*$), the bus request lines (for example, $BUSRQ<1..0>^*$) and the bus acknowledge lines (for example, $BUSAK<1..0>^*$).

1.5 Logical and Electrical State Relationships. Throughout this standard the term *asserted* is used to indicate the logical true state of the particular signal referenced. The corresponding term *negated*, however, is not used because it comprises a potentially ambiguous representation when describing signals, which may be low or high true.

All signals that are low in their asserted state are designated by a nathan (asterisk), which follows the signal name (for example, $ADRSTB^*$). The correlation between the terms true:false, high:low, and asserted:released is demonstrated in the following table, utilizing the signals $ADRSTB^*$ and $CM<2..0>$ as an example.

Function	Electrical	Logical	State
CM<2..0>	High	1 True	Active, asserted
	Low	0 False	Active, released
	High Z	—	Inactive
ADRSTB*	Low	1 True	Active, asserted
	High	0 False	Active, released
	High Z	—	Inactive

2. Functional Description

This section describes the functional elements of IEEE Std 1000 Bus interface. They are

- (1) System Controller
- (2) Arbiter
- (3) Masters
- (4) Slaves

An individual board attached to IEEE Std 1000 Bus backplane may consist of one or more of these elements.

2.1 System Controller. Within any IEEE Std 1000 Bus system there shall be one, and only one, system controller. The system controller provides essential facilities for the proper operation of IEEE Std 1000 Bus systems. The system controller may be combined on a board with a master.

The system controller shall provide as a minimum the following requirements:

- (1) SYSCLK. A general-purpose clock signal in accordance with the specifications detailed in Section 3 of this standard.
- (2) SYSRST*. An initial power-on system reset signal in accordance with the specifications detailed in Section 3 of this standard.
- (3) TFRERR*. A transfer error signal in accordance with the specifications detailed in Section 3 of this standard.

2.2 Arbiter. All bus allocation grants shall be provided by the arbiter in accordance with the protocol described in Section 4 of this standard. There shall be one, and only one, arbiter within any IEEE Std 1000 Bus system. The arbiter may be combined on a board with a master.

2.3 Masters. A master is a board that is capable of controlling the transfer of data on the bus, by means of the protocols defined in Sections 4 and 5 of this standard. A master may contain a central processing unit (cpu) or logic necessary to transfer data over the bus (for example, DMA controller).

2.3.1 Master Types. All masters must request allocation of IEEE Std 1000 Bus from the arbiter before they can control data transfers except in the special case of a default master.

default master. A master that is allocated control of the bus by the arbiter whenever the bus is not in use by another master. A default master is necessarily combined with the arbiter on the same board and has the lowest priority for bus allocation. There can be only one default master within a IEEE Std 1000 Bus system, though a IEEE Std 1000 Bus system need not include a default master.

All other masters (termed potential masters) request allocation of bus control from the arbiter. This request shall be made by asserting one of the two bus request (BUSRQ n^*) lines, or in the case of a master that is on the same board as the arbiter but is not configured as a default master, by asserting a third bus request line that is private to that board.

Multiple masters may exist within a system. The master's priority for bus control allocation and the means by which such allocation is accomplished are as described in Section 4.

2.3.2 Master Modes. Masters may retain control for a period of time constrained only by the specific system requirements. Two modes of operation may be used by masters.

(1) *Release-When-Done.* The master retains control of the bus until all desired transfers have been accomplished.

(2) *Release-On-Request.* The master retains control of the bus indefinitely, relinquishing control when it determines that another master requires allocation of the bus. This determination may be made by several methods including receipt of an attention request signal, hardware polling, or by detection of a BUSRQ n^* signal becoming active.

NOTE: By definition a default master always operates in release-on-request mode.

2.4 Slaves. Boards that are capable of being controlled over the IEEE Std 1000 Bus are designated slaves. Slaves decode the address lines and act upon the command provided by the current master. A slave may be combined with other functional elements on a board (for example, a board that contains both a master and memory that is accessible by other masters within the system).

Figure 1 shows one possible system configuration that utilizes a variety of boards in combination.

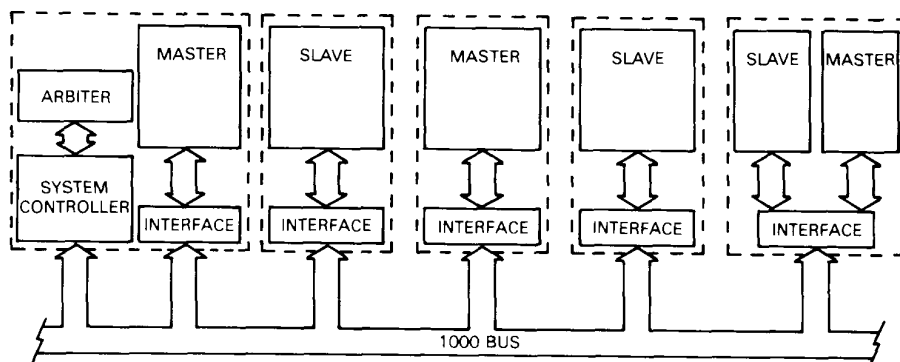


Fig 1
Example of System Configuration

3. Signal Lines

This section provides specific definitions for all signal lines that are part of the IEEE Std 1000 Bus. Each of these signals has been assigned to one of five functional groups. These groups are

- (1) Information Lines
 - (a) Address Lines
 - (b) Data Lines
 - (c) Command Lines
- (2) Synchronization Lines
- (3) Attention Request Lines
- (4) Bus Allocation Lines
- (5) Utility Lines

3.1 Information Lines

3.1.1 Address Lines (A<19..0>). These unidirectional lines specify the address of the referenced memory or I/O location or, during a *vector fetch* response to an attention request, the level of the request being acknowledged. The most significant bit is A19 and A0 is the least significant.

The following table details the usage of the address lines during various types of operations.

Operation	Valid Lines	Total Addressed Range
Memory read or write	A<19..0>	1 048 576 bytes
I/O read or write	A<11..0>	4096 locations
Vector-fetch	A<2..0>	8 levels

3.1.2 Data Lines (D<7..0>). These eight bidirectional lines carry information between masters and slaves. The most significant bit is D7 and D0 is the least significant.

3.1.3 Command Lines (CM<2..0>). These signals are used by the current master to convey coded data to the slave describing the type of the current data transfer according to Table 1.

Command codes marked *reserved* shall not be used, and IEEE Std 1000 boards shall not respond to or utilize these codes for any purpose so as to be considered in compliance with this standard. This is to guarantee compatibility with boards that may be designed to conform to future revisions of this standard.

Table 1
Command Codes

CM2	CM1	CM0	Transfer
0	0	0	Reserved
0	0	1	Reserved
0	1	0	Reserved
0	1	1	Vector-fetch
1	0	0	I/O Write
1	0	1	I/O Read
1	1	0	Memory write
1	1	1	Memory read

3.2 Synchronization Lines. The following signals are classified as synchronization lines:

Signal	Function
ADRSTB*	Address strobe
DATSTB*	Data strobe
DATAACK*	Data transfer acknowledge
TFRERR*	Transfer error

3.2.1 Address Strobe (ADRSTB*). This signal indicates the presence of valid data on the address lines.

3.2.2 Data Strobe (DATSTB*). This signal indicates the presence of valid data on the command lines CM<2..0>. During a read, or vector-fetch transfer, this signal indicates that the addressed slave may place data on the data lines. During a write transfer this signal indicates the presence of valid data on the data lines.

3.2.3 Data Transfer Acknowledge (DATAACK*). This signal is used to indicate to the master that the command has been performed: that data have been placed on, or accepted from, the data lines.

3.2.4 Transfer Error (TFRERR*). This signal may be asserted by any board to indicate an error during the current transfer. Specific timing requirements for this signal are detailed in Section 5 of this standard.

3.3 Attention Request Lines (ATNRQ<7..0>*). These signals are configured for indicating user-specific events when a IEEE Std 1000 Bus system is commissioned. Such events may include, but are not limited to, interrupt requests, DMA requests, or notification of conditions, which exist either at the board or system level (for example, failure). Eight attention request lines are available. Three optional response protocols are described in Section 6 of this standard for the use of attention request lines as traditional interrupts.

These signals may be used by any board to request the attention of other boards within IEEE Std 1000 Bus systems. Any board within the system may be connected to any of the eight attention request lines. Multiple boards may be connected to the same attention request line allowing for the broadcast of events to one or more boards within a system. There is an implied priority with ATNRQ0* having highest priority and ATNRQ7* having the lowest.

3.4 Bus Allocation Lines. The bus allocation lines are:

Signal	Function
BUSRQ<1..0>*	Bus request lines
BUSAK<1..0>*	Bus acknowledge lines

3.4.1 Bus Request Lines (BUSRQ<1..0>*). These signals may be asserted by any potential master that desires allocation of the IEEE Std 1000 Bus. In systems utilizing prioritized arbitration, BUSRQ0* shall have priority over BUSRQ1*.

3.4.2 Bus Acknowledge Lines (BUSAK<1..0>*). These signals are used by the arbiter to indicate to a master requesting bus allocation that it may take control of the bus. BUSAK1* indicates a grant to the master requesting by way of BUSRQ1*, and BUSAK0* indicates a grant to the master requesting by way of BUSRQ0*.

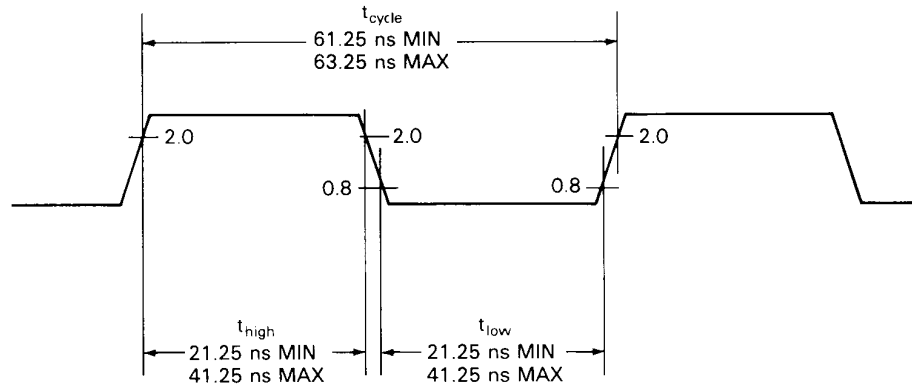
3.5 Utility Lines. The following signals are classified as utility lines:

Signal	Function
SYSCLK	System clock
SYSRST*	System reset

3.5.1 System Clock (SYSCLK). The system clock is a periodic signal of constant frequency that may be used as a generalized facility by masters or slaves. The system clock is independent of the protocol of any other bus signals, and is provided by the system controller. Figure 2 contains timing waveforms for SYSCLK.

3.5.2 System Reset (SYSRST*). The system reset signal is used to place the system in a known initial state. While SYSRST* is active all boards shall inhibit any access to the system bus. SYSRST* may be driven by any board. It is recommended that during a power-up sequence, any board capable of performing on-board diagnostic self-tests hold SYSRST* active until the successful completion of such tests.

The system controller shall provide an initial power-on system reset signal that is not <200 ms and not >500 ms in duration, measured from the point at which the +5 V dc supply reaches its designated minimum specification (see Section 7). The system controller also shall assert SYSRST* at any time that the



$$\begin{aligned} t_{cycle} &= 62.5 \text{ ns} \pm 1 \text{ ns} \text{ (16.00 MHz)} \\ t_{high} &= 31.25 \text{ ns} \pm 10 \text{ ns} \\ t_{low} &= 31.25 \text{ ns} \pm 10 \text{ ns} \end{aligned}$$

Fig 2
SYSCLK Timing

system supply falls below its minimum specified tolerance, and shall continue to assert it for the entire period during which the supply is out of tolerance. The rise time of this signal shall not exceed 100 ns (10%-90%).

A recommended power-fail protocol, for implementation of systems where an early indication of primary power supply failure is available, is provided in Section 7.

3.6 IEEE Std 1000 Bus Connector Pin Allocations

3.6.1 Connector Type. The IEEE Std 1000 Bus connector pair shall be C064 from IEC 603-2 (1980). The male connector shall be on the board and the female connector shall be on the backplane.

3.6.2 Connector Pin Allocation. The allocation of the signals described in this standard to pins on the IEEE Std 1000 Bus connector are in Fig 3.

The backplane tracking arrangement shown is recommended so as to allow the signal 0 V return traces to provide guard track protection between strobes and signal sets.

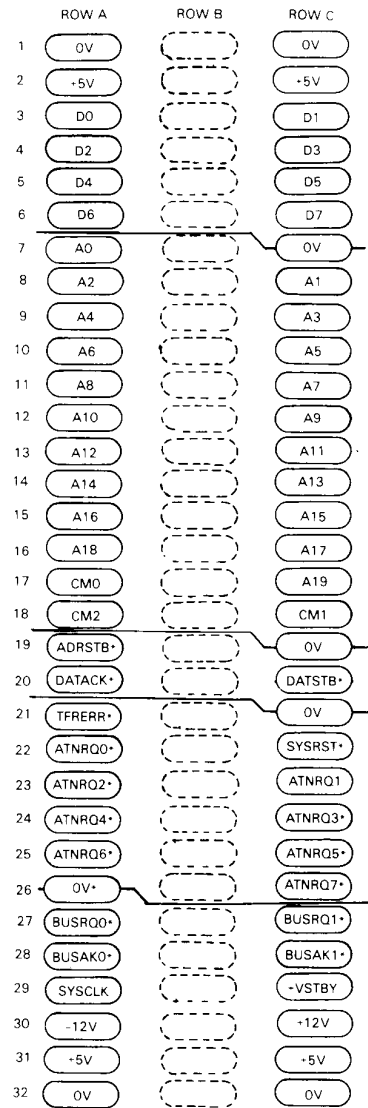


Fig 3
IEEE Std 1000 Bus Connector

4. Arbitration

This section defines the protocol that is used by IEEE Std 1000 Bus masters to gain control of the IEEE Std 1000 Bus.

The following signals are used to effect allocation of bus control within IEEE Std 1000 Bus systems:

BUSRQ<1..0>*

BUSAK<1..0>*

4.1 Arbitration Algorithm. On power-up, or following a system reset, the arbiter shall have control of the bus. All masters (except default masters) must request and receive a control allocation grant from the arbiter prior to controlling the bus. Only the arbiter may effect a change in allocation of bus control within an IEEE Std 1000 Bus system.

Any algorithm may be used by the arbiter to determine which of the bus request levels will be granted when the current master releases the bus, although, by convention, the BUSRQ0*/BUSAK0*, request/grant pair have priority over BUSRQ1*/BUSAK1*. The arbiter may award preference to one requesting level on a priority basis, a *round-robin* basis, or by means of any other algorithm implementable within the control transfer protocol.

4.2 Bus Requests. A Bus request may be initiated by any master within an IEEE Std 1000 Bus system at any time by asserting one of the BUSRQ* lines. Bus requests are level-triggered, meaning that a requesting condition is indicated by the relative voltage level present on a BUSRQ n * line rather than by the active transition from one state to another.

A master, having asserted one of the BUSRQ n * lines, shall continue to assert it for the entire period during which it desires to control the bus.

NOTE: Multiple masters may be connected to a single BUSRQ n * line only when it can be assured that more than one master cannot request the bus at the same time, on the same requesting level.

4.3 Bus Grants. BUSAK0* or BUSAK1* are asserted by the arbiter in response to the corresponding requests, BUSRQ0* and BUSRQ1*. The arbiter shall not assert either of these grant lines in the absence of a corresponding bus request. In addition, the arbiter shall not assert both bus acknowledge lines simultaneously under any circumstances.

A master asserting a bus request line and detecting the corresponding bus acknowledge line asserted may assume control of the bus. The arbiter, having asserted a bus acknowledge line, shall continue to assert it for the entire period

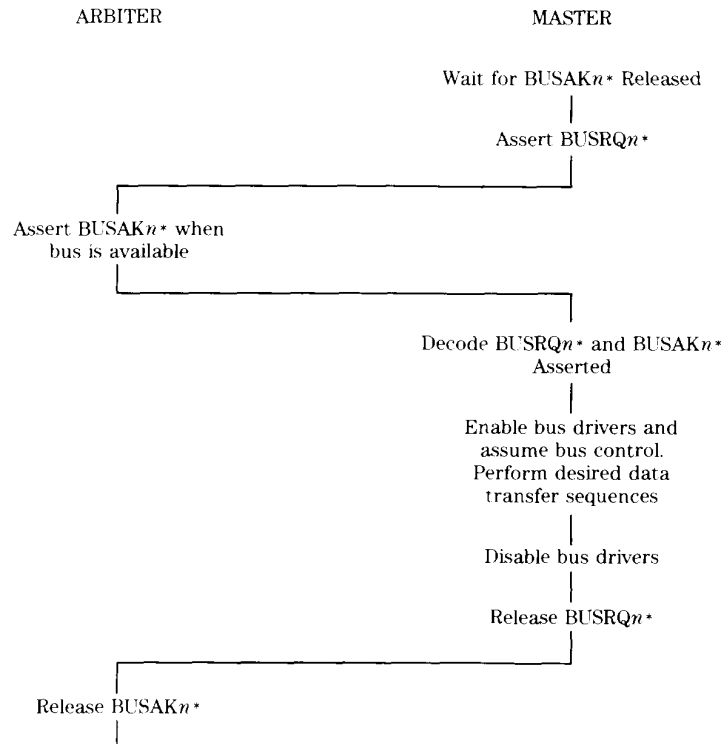
during which the current master continues to assert the corresponding bus request line.

4.4 Control Allocation Sequence. Figure 4 shows the handshake-signal flow during an allocation of bus control. The sequence of events for all allocations of bus control shall conform to the following description:

A master may request allocation of the IEEE Std 1000 Bus at any time by asserting one of the BUSRQ n^* lines except that a master may not assert BUSRQ n^* until the corresponding BUSAK n^* has been released from a previous cycle.

The arbiter, upon detecting BUSRQ n^* active, and having first assured that the bus is available, shall assert BUSAK n^* indicating that the bus is available for use

Fig 4
Control Transfer Flow Diagram



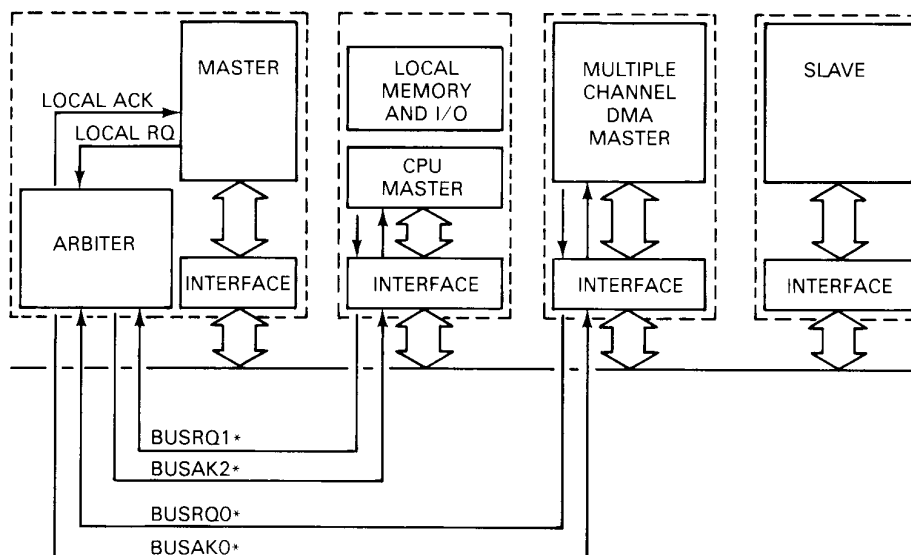
by the requesting master. The arbiter shall not assert $BUSAK_n^*$ in the absence of a corresponding $BUSRQ_n^*$. The bus is available for arbitration if the arbiter is not issuing a bus grant to any master, either by asserting $BUSAK_n^*$ or by granting to the bus to its own on-board master (if one is present).

The requesting master, upon detecting that $BUSAK_n^*$ has been made active, shall assume control of the bus.

Upon completing the desired data transfer sequences, the current master shall cease driving all bus lines and shall then release $BUSRQ_n^*$.

The arbiter, upon detecting that the $BUSRQ_n^*$ is no longer asserted, shall release $BUSAK_n^*$.

Fig 5
Example of Multiple Master Implementation



5. Data Transfer Protocol

This section defines the protocol and timing requirements necessary to effect data transfer sequences across the IEEE Std 1000 Bus. All data are transferred as single bytes using an asynchronous interlocked handshake. The asynchronous nature of the protocol allows communication between masters and slaves of widely differing speeds, and is essentially self-adaptive to changing system configurations.

Indivisible read-modify-write sequences are accommodated within the protocol, allowing synchronization of tasks in a multiple processor system by use of lock variables.

A high-speed *burst* mode of data transfer is also specified, allowing for rapid movement of contiguous data blocks between boards.

Figure 6 illustrates the various types of sequences accommodated by the IEEE Std 1000 Bus data transfer protocol.

The signals used for data transfer sequences are

Signal	Source
A<19..0>	Master
D<7..0>	Master during write, slave during read, and vector-fetch
CM<2..0>	Master
ADRSTB*	Master
DATSTB*	Master
DATAACK*	Slave
TFRERR*	Any board

Throughout this section the word *master* implies current master, that is, it is assumed that the master discussed has been allocated control of the IEEE Std 1000 Bus by means of the protocol described in Section 4 of this standard.

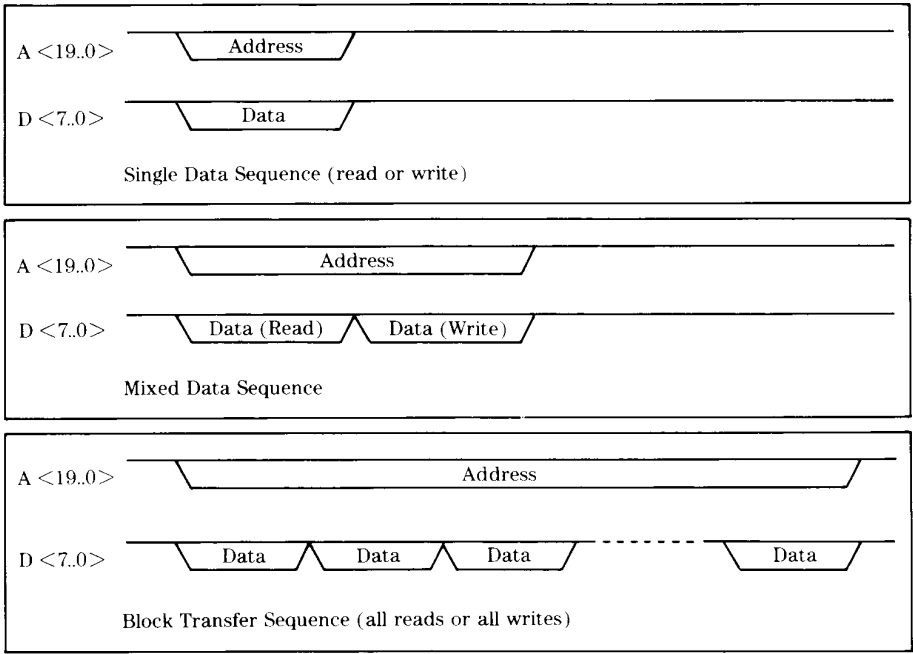


Fig 6
IEEE Std 1000 Bus Data Transfer Sequences

5.1 Read Sequence. Data transfers from slave to master are designated read sequences. Figure 7 illustrates the handshake-signal flow, and Fig 10 specifies the signal timing parameters for a read sequence. Read sequences shall conform to the following description:

(1) The master places the address of the targeted memory or I/O location on the address lines.

NOTE: The master may assert the appropriate command code at this time.

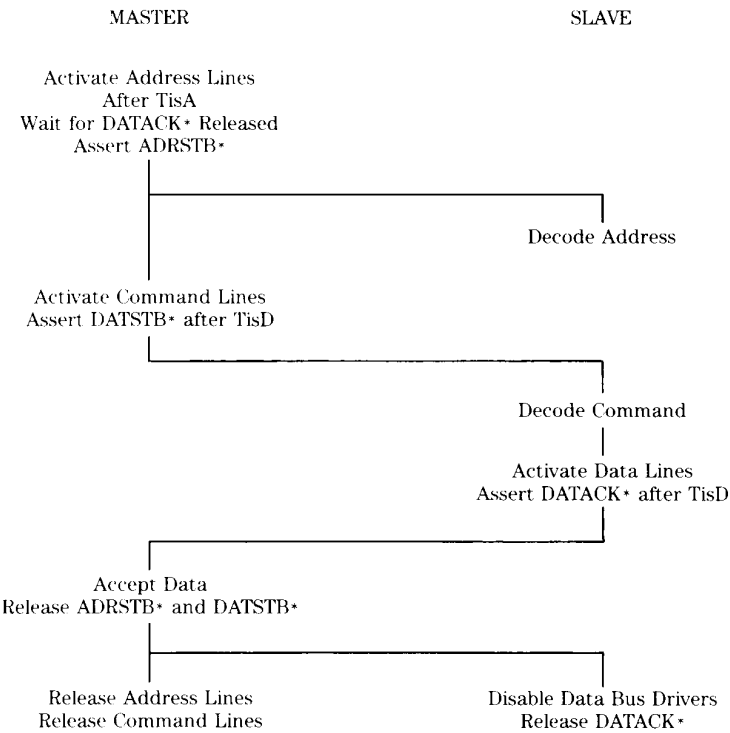


Fig 7
Read Sequence Flow Diagram

- (2) After a setup time, during which the address lines become valid, and having ensured that DATAACK* is released from the previous cycle, the master shall assert ADRSTB*.
- (3) The master shall activate the command lines, CM<2..0>, to indicate the type of transfer.

CM2	CM1	CM0	Transfer
1	0	1	I/O read
1	1	1	Memory read

(4) After a setup time during which the command lines become valid, the master shall assert DATSTB* indicating the presence of valid data on the command lines, and that it is ready to accept data.

(5) The addressed slave shall enable its data bus drivers, placing requested data on the data lines.

(6) After a setup time, during which the data lines become valid, the addressed slave shall assert DATAACK* indicating that the data is available.

(7) In response to DATAACK* the master shall accept the data and shall release DATSTB* indicating to the slave that it must remove the data from the data lines.

(8) Upon detecting either DATSTB* or ADRSTB* released, the slave disables its data bus drivers and releases DATAACK* indicating that the sequence is complete.

5.2 Write Sequence. Data transfers from masters to slaves are designated write sequences. Figure 8 illustrates the handshake signal flow, and Fig 11 specifies the signal timing parameters for a write sequence. Write sequences shall conform to the following description:

(1) The master places the address of the targeted memory or I/O location on the address lines.

NOTE: The master may assert the appropriate command code at this time.

(2) After a setup time, during which the address lines become valid, and having ensured DATAACK* is released from the previous cycle, the master shall assert ADRSTB*.

(3) The master shall activate the command lines, CM<2..0>, to indicate the type of transfer, and shall place the data to be transferred on the data lines.

CM2	CM1	CM0	Transfer
1	0	0	I/O write
1	1	0	Memory write

(4) After a setup time, during which the data lines and the command lines become valid, the master shall assert DATSTB* indicating the presence of valid data on the data and command lines.

(5) The addressed slave shall accept the data and shall assert DATAK* to indicate that the transfer may be terminated.

(6) In response to DATAK* the master shall release ADRSTB* and DATSTB*.

(7) Upon detecting either DATSTB* or ADRSTB* released, the slave shall release DATAK* indicating that the sequence is complete.

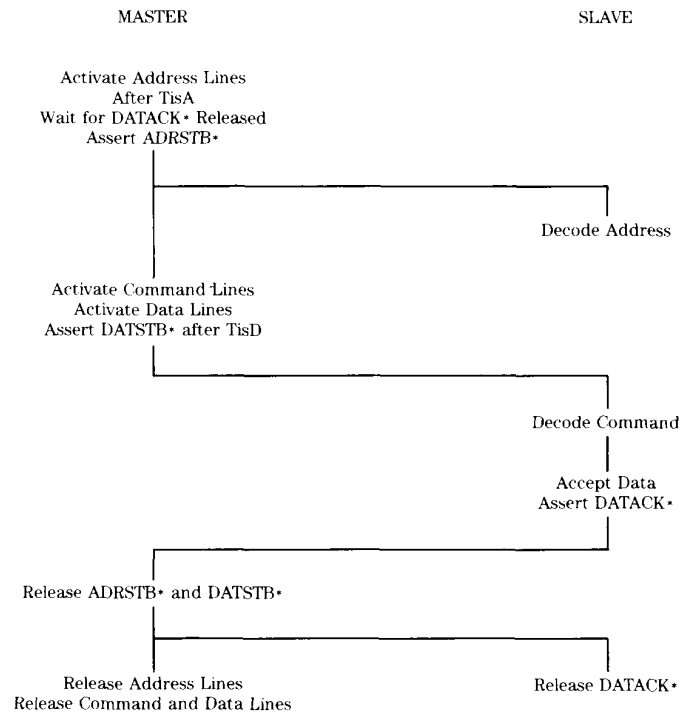


Fig 8
Write Sequence Flow Diagram

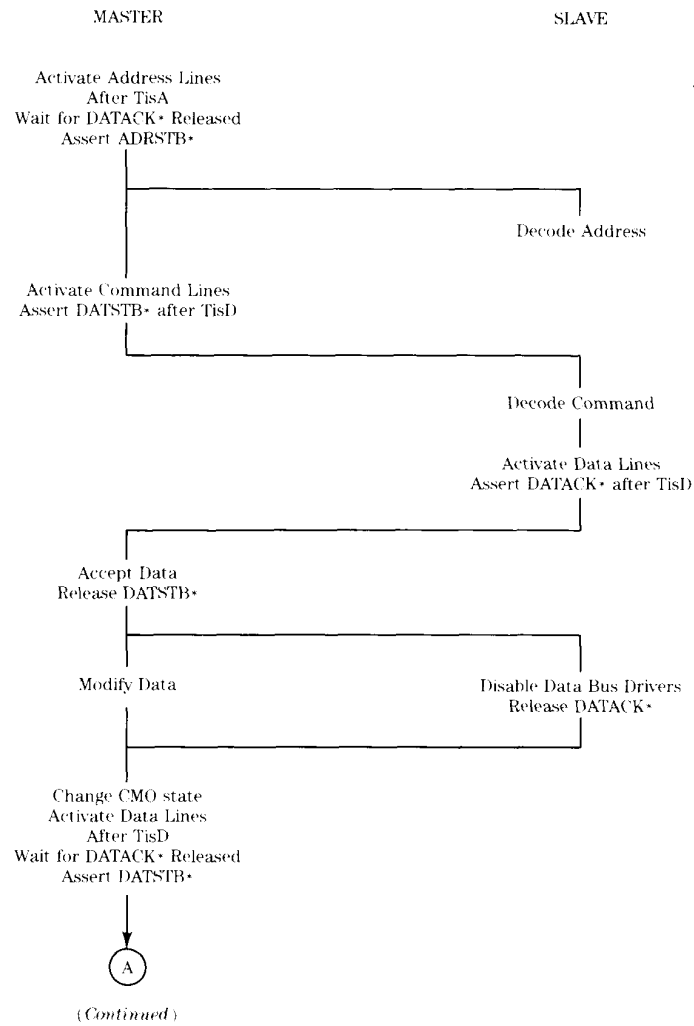


Fig 9
Read-Modify-Write Sequence Flow Diagram

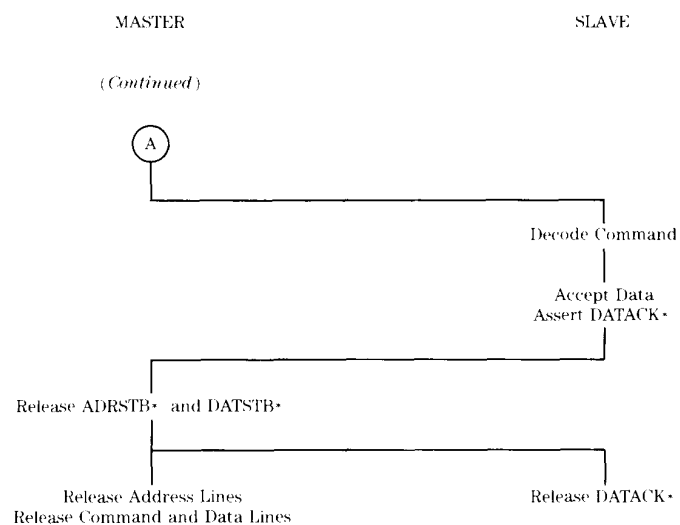


Fig 9 (Continued)
Read-Modify-Write Sequence Flow Diagram

5.3 Read-Modify-Write Sequence. Sequences during which the data is transferred to the master, operated on by the master, and subsequently transferred from the master to the same address, in a single indivisible sequence are designated read-modify-write sequences. Figure 9 illustrates the handshake signal flow, and Fig 12 specifies the signal timing parameters for a read-modify-write sequence. Read-modify-write sequences shall conform to the following description:

(1) The master places the address of the targeted memory or I/O location on the address lines.

NOTE: The master may assert the appropriate command code at this time.

(2) After a setup time, during which the address lines become valid, and having ensured DATAACK* is released from the previous cycle, the master shall assert ADRSTB*.

(3) The master shall activate the command lines, CM<2..0>, to indicate the type of transfer.

CM2	CM1	CM0	Transfer
1	0	1	I/O read
1	1	1	Memory read

(4) After a setup time during which the command lines become valid the master shall assert DATSTB* indicating the presence of valid data on the command lines, and that it is ready to accept data.

(5) The addressed slave shall enable its data bus drivers, placing the requested data on the data lines. A slave shall not respond to DATSTB* unless ADRSTB* is asserted.

(6) After setup time, during which the data lines become valid, the addressed slave shall assert DATAACK* indicating that the data is available.