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Information technology — Telecommunications and information exchange between systems — Local and metropolitan area networks — Specific requirements —

Part 6:

Distributed Queue Dual Bus (DQDB) access method and physical layer specifications

*Technologies de l'information — Communication de données et échange
d'informations entre systèmes — Réseaux locaux et métropolitains — Exigences
spécifiques —*

*Partie 6: Bus distribué à double queue (DQDB) et spécifications pour la couche
physique*



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Abstract: This standard is part of a family of standards for local area networks (LANs) and metropolitan area networks (MANs) that deals with the Physical and Data Link Layers as defined by the ISO Open Systems Interconnection Reference Model. It defines a high-speed shared medium access protocol for use over a dual, counterflowing, unidirectional bus subnetwork. The Physical Layer and Distributed Queue Dual Bus (DQDB) Layer are required to support a Logical Link Control (LLC) Sublayer by means of a connectionless Medium Access Control (MAC) Sublayer service in a manner consistent with other IEEE 802 networks. Additional DQDB Layer functions are specified as a framework for other services. These additional functions will support Isochronous Service Users and Connection-Oriented Data Service Users, but their implementation is not required for conformance.

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**Information technology—
Telecommunications and information
exchange between systems—
Local and metropolitan area networks—
Specific requirements—**

**Part 6: Distributed Queue Dual Bus
(DQDB) access method and physical
layer specifications**

Sponsor

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International Standard ISO/IEC 8802-6 : 1994

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International Standard ISO/IEC 8802-6 was prepared by Joint Technical Committee ISO/IEC JTC 1, *Information technology*.

ISO/IEC 8802 consists of the following parts, under the general title *Information technology—Local and metropolitan area networks*:

- *Part 1: Overview and Architecture*
- *Part 2: Logical link control*
- *Part 3: Carrier sense multiple access with collision detection (CSMA/CD) access method and physical layer specifications*
- *Part 4: Token-passing bus access method and physical layer specifications*
- *Part 5: Token ring access method and physical layer specifications*
- *Part 6: Distributed Queue Dual Bus (DQDB) access method and physical layer specifications*
- *Part 7: Slotted ring access method and physical layer specification*

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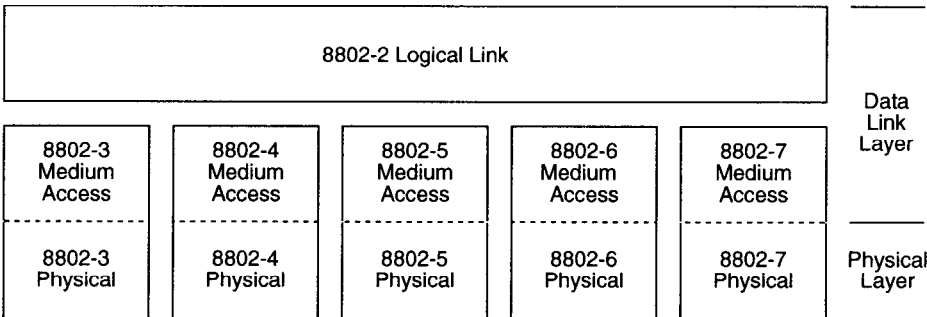
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Foreword to International Standard ISO/IEC 8802-6 : 1994

This standard is part of a family of standards for Local and Metropolitan Area Networks. The relationship between this standard and the other members of the family is shown below. (The numbers in the figure refer to ISO standard numbers.)



This family of standards deals with the Physical and Data Link layers as defined by the ISO Open Systems Interconnection Basic Reference Model (ISO 7498 : 1984). The access standards define five types of medium access technologies and associated physical media, each appropriate for particular applications or system objectives. Other types are under investigation.

The standards defining the access technologies are as follows:

- a) ISO/IEC 8802-3 [ANSI/IEEE Std 802.3, 1993 Edition], a bus utilizing CSMA/CD as the access method.
- b) ISO/IEC 8802-4 [ANSI/IEEE Std 802.4-1990], a bus utilizing token passing as the access method.
- c) ISO/IEC 8802-5 [ANSI/IEEE Std 802.5-1992], a ring utilizing token passing as the access method.
- d) ISO/IEC 8802-6 [ANSI/IEEE Std 802.6, 1994 Edition], a dual bus utilizing distributed queuing as the access method. DQDB subnetworks provide a range of telecommunications services within a metropolitan area.
- e) ISO 8802-7, a ring utilizing slotted ring as the access method.

ISO 8802-2 [ANSI/IEEE Std 802.2-1989], *Logical Link Control protocol*, provides for data transfer between medium access standards and network layer protocol.

ISO/IEC 10038 [ANSI/IEEE Std 802.1D, 1993 Edition], *Media access control (MAC) bridges*, specifies an architecture and protocol for the interconnection of IEEE 802 LANs below the level of the logical link control protocol.

The reader of this document is urged to become familiar with the complete family of standards.

ANSI/IEEE Std 802.6, 1994 Edition

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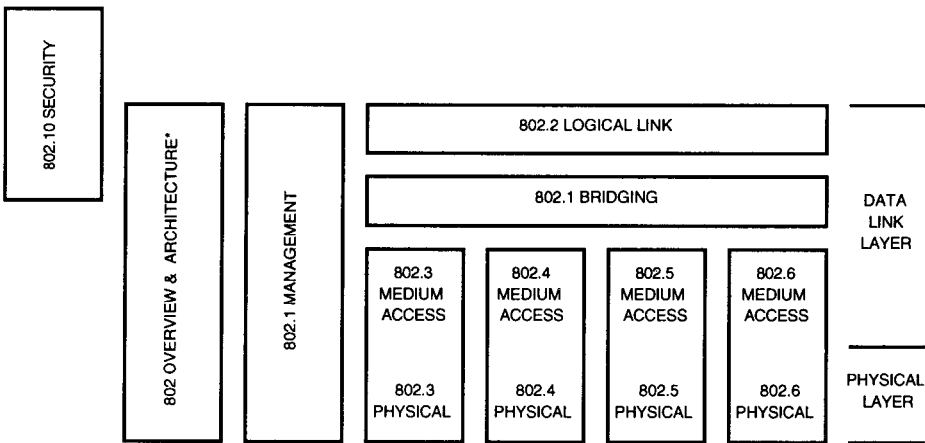
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Introduction

(This introduction is not a part of ANSI/IEEE Std 802.6, 1994 Edition or of ISO/IEC 8802-6 : 1994.)

This standard is part of a family of standards for local and metropolitan area networks. The relationship between the standard and other members of the family is shown below. (The numbers in the figure refer to IEEE standard numbers.)



* Formerly IEEE Std 802.1A.

This family of standards deals with the Physical and Data Link layers as defined by the International Organization for Standardization (ISO) Open Systems Interconnection Basic Reference Model (ISO 7498 : 1984). The access standards define several types of medium access technologies and associated physical media, each appropriate for particular applications or system objectives. Other types are under investigation.

The standards defining these technologies are as follows:

- IEEE Std 802¹: Overview and Architecture. This standard provides an overview to the family of IEEE 802 Standards. This standard forms part of the 802.1 scope of work.
- IEEE Std 802.1B [ISO DIS 15802-2]: LAN/MAN Management. Defines an Open Systems Interconnection (OSI) management-compatible architecture, and service and protocol elements for use in a LAN/MAN environment for performing remote management.
- ISO/IEC 10038 [ANSI/IEEE Std 802.1D]: MAC Bridging. Specifies an architecture and protocol for the interconnection of IEEE 802 LANs below the MAC service boundary.
- IEEE Std 802.1E [ISO DIS 15802-4]: System Load Protocol. Specifies a set of services and protocol for those aspects of management concerned with the loading of systems on IEEE 802 LANs.

¹The 802 Architecture and Overview Specification, originally known as IEEE Std 802.1A, has been renumbered as IEEE Std 802. This has been done to accommodate recognition of the base standard in a family of standards. References to IEEE Std 802.1A should be considered as references to IEEE Std 802.

- ISO 8802-2 [ANSI/IEEE Std 802.2]: Logical Link Control
- ISO/IEC 8802-3 [ANSI/IEEE Std 802.3]: CSMA/CD Access Method and Physical Layer Specifications
- ISO/IEC 8802-4 [ANSI/IEEE Std 802.4]: Token Bus Access Method and Physical Layer Specifications
- ISO/IEC 8802-5 [ANSI/IEEE Std 802.5]: Token Ring Access Method and Physical Layer Specifications
- ISO/IEC 8802-6 [ANSI/IEEE Std 802.6]: Distributed Queue Dual Bus (DQDB) Access Method and Physical Layer Specifications
- IEEE Std 802.10: Interoperable Local Area Network (LAN) Security, *Currently Contains Secure Data Exchange (SDE)*

In addition to the family of standards, the following is a recommended practice for a common technology:

- IEEE Std 802.7: IEEE Recommended Practice for Broadband Local Area Networks

The reader of this standard is urged to become familiar with the complete family of standards.

Conformance test methodology

An additional standards series, identified by the number 1802, has been established to identify the conformance test methodology documents for the 802 family of standards. This makes the correspondence between the various 802 standards and their applicable conformance test requirements readily apparent. Thus the conformance test documents for 802.3 are numbered 1802.3, the conformance test documents for 802.5 will be 1802.5, and so on. Similarly, ISO will use 18802 to number conformance test standards for 8802 standards.

IEEE Std 802.6, 1994 Edition

The purpose of this standard is to lay the foundation for a set of standards to allow DQDB subnetworks to provide a range of telecommunications services within a metropolitan area. Based on this set of standards, equipment vendors will be able to build the components that will allow telecommunications services to be offered to end users within a metropolitan area.

The interconnection of DQDB subnetworks to form a metropolitan area network will be possible via a Multiport Bridge² or via dual-port bridges, routers, and gateways, as shown in Fig A. DQDB subnetworks can be used to provide switching, routing, and concentration of high-speed data, voice, and certain video services, as well as to interconnect LANs, hosts, workstations, and PBXs.

The user's attention is called to the possibility that compliance with this standard may require the use of inventions covered by patent rights. All enquiries should be sent to the IEEE Standards Board at the address given at the end of the introduction.

This edition of the standard defines general principles for the operation of a number of physical layers. A physical layer capable of using a DS3 transmission system (ANSI T1.102 and T1.107) is defined in this edition of the standard. It is anticipated that future editions of the standard will provide additional implementations of the physical layer to support different needs (for example, media, data rates, network operator requirements).

²The services of a Multiport Bridge are the subject of ongoing work under IEEE Project Authorization P802.6a, Multiple Port Bridging for Metropolitan Area Networks.

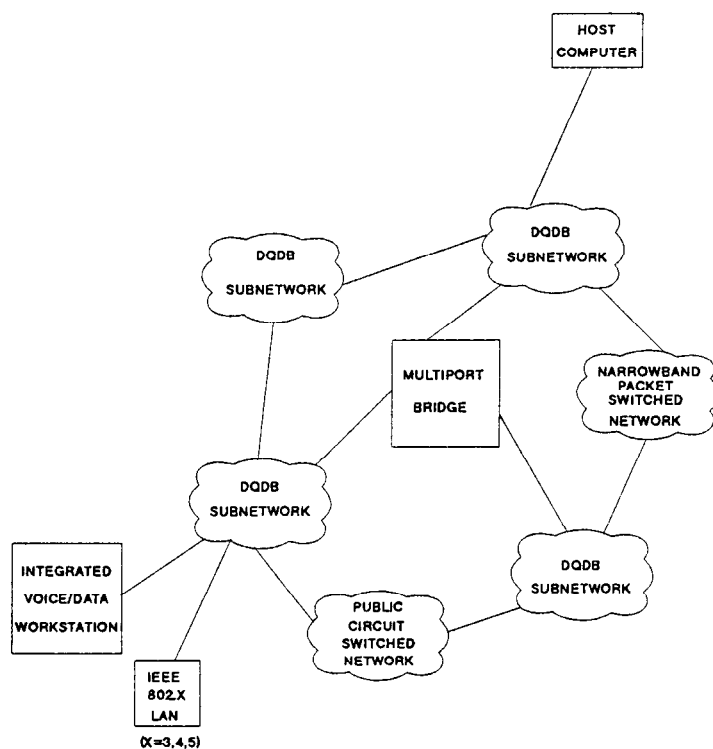


Figure A—DQDB metropolitan area network

This edition of the standard incorporates IEEE Std 802.6d-1993, *Physical Layer Convergence Procedure (PLCP) for CCITT Recommendations G.707, G.708, and G.709 SDH-based systems (155.520 Mbit/s)*, as clause 17. It also incorporates IEEE Std 802.6f-1993, *Protocol Implementation Conformance Statement (PICS) proforma*, as annex A (normative). The two standards are not available as separate publications.

Three additional IEEE 802.6 standards are currently available. These include the following:

- IEEE Std 802.6k-1992, Supplement to Media Access Control (MAC) Bridges (IEEE Std 802.1D-1990): IEEE 802.6 Distributed Queue Dual Bus (DQDB) Subnetwork of a Metropolitan Area Network (MAN)
- IEEE Std 802.6c-1993, Supplement to ISO/IEC 8802-6: Physical Layer Convergence Procedure (PLCP) for DS1-based systems (clause 12) *packaged with IEEE Std 802.6h-1993 (see below)*
- IEEE Std 802.6h-1993, Supplement to ISO/IEC 8802-6: Isochronous Service on a DQDB Subnetwork of a MAN

This standard contains state-of-the-art material. The area covered by this standard is undergoing evolution. Revisions are anticipated within the next few years to clarify existing material, to correct possible errors, and to incorporate new related material. Information on the current revision state of this and other IEEE 802 standards may be obtained from

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The working group is especially grateful for the dedicated work of Peter Evans and Sayeed Ghani, editors of IEEE Std 802.6-1990, who translated the complexities of a protocol involving many thousands of silicon gates into an accurate and usable document. Particular thanks are due also to Dan Sze, who authored several proposals in earlier stages of the MAN standard development and also served as vice chair of the committee; to Mary Kelly, whose work behind the scenes helped to make it all come together.

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The final conditions for approval of IEEE Std 802.6d-1993 and IEEE Std 802.6f-1993 were met on July 12, 1993. These standards were conditionally approved by the IEEE Standards Board on June 17, 1993, with the following membership:

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Peter Evans was the first technical editor of the International Standard (ISO/IEC 8802-6 : 1994). Ingrid Fromm succeeded him in 1993 and saw the document through publication. Kristin Dittmann was the IEEE Standards Project Editor.

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**Information technology—
Telecommunications and information exchange between
systems—
Local and metropolitan area networks—
Specific requirements—**

Part 6: Distributed Queue Dual Bus (DQDB) access method and physical layer specifications

1. Overview

1.1 Scope

This part of ISO/IEC 8802 for a Distributed Queue Dual Bus (DQDB) subnetwork of a metropolitan area network (MAN) defines a high-speed shared medium access protocol for use over a dual, counter-flowing, unidirectional bus subnetwork. This International Standard specifies the Physical Layer and DQDB Layer required to support

- *Logical Link Control (LLC) Sublayer*, by a connectionless Medium Access Control (MAC) Sublayer service provided to support an LLC Sublayer in a manner consistent with other parts of this ISO/IEC 8802 series of International Standards.

This part of ISO/IEC 8802 also specifies additional DQDB Layer functions as a framework for other services. In this edition of the standard, they are provided for completeness as part of the overall architecture; their implementation is not required for conformance.

These additional functions will support

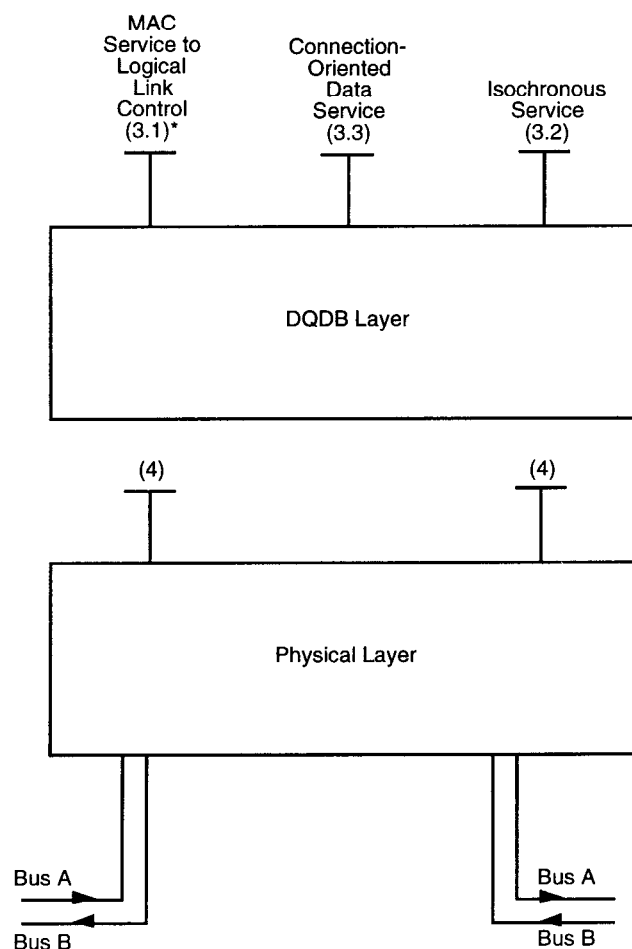
- *Isochronous Service Users (ISUs)*, by a connection-oriented service, which may be used to transport isochronous data; for example, conventional digitized voice.
- *Connection-Oriented Data Service Users*, by an asynchronous, connection-oriented service, which may be used to transport bursty data; for example, signaling.

The scope of this part of ISO/IEC 8802 is shown diagrammatically in figure 1-1.

The DQDB Layer supports these services by employing the following access methods:

- A *Queued Arbitrated access method*, with three priority queues for medium access arbitration and fixed-length slots for data transfer. Each priority level provides distributed queue access for the support of connectionless MAC service and connection-oriented data service.
- A *Pre-Arbitrated access method*, which uses assigned octet positions in particular slots for the transfer of individual octets of data. This access method supports isochronous connection-oriented services.

The Queued Arbitrated access method uses fixed-length slots for data transfer. This transfer mechanism is enhanced by a MAC convergence function to the service expected by the LLC Sublayer. The MAC convergence function is defined in this part of ISO/IEC 8802.



*Numbers in parentheses refer to clauses and subclauses of this part of ISO/IEC 8802 where the services are defined.

Figure 1-1—Scope of this part of ISO/IEC 8802

This part of ISO/IEC 8802 defines the DQDB Layer functions to allow access to the medium to read and write data to support the isochronous service and the connection-oriented data service. The signaling procedures for establishing, maintaining, and clearing a connection are outside the scope of this part of ISO/IEC 8802.¹

The Physical Layer is defined to allow the use of different underlying transmission systems. This International Standard supports a DS3 interface operating at 44.736 Mbit/s, as specified in ANSI T1.102² and T1.107. It also supports transmission systems based on CCITT Recommendation G.703 operating at 2.048 Mbit/s, 34.368 Mbit/s, and 139.264 Mbit/s, and Synchronous Digital Hierarchy (SDH) transmission systems operating at 155.520 Mbit/s, as specified in CCITT Recommendations G.707, G.708, and G.709. Other appropriate media and transmission rates may be considered for standardization in the future.

¹ Annex B gives a tutorial description of the sequence of events that are required to support complete isochronous communication capability.

² Information on references can be found in 1.3.

Each transmission system may need to be enhanced to provide the Physical Layer service that is required by the DQDB Layer. This enhancement function, if required, is provided by a suitable Physical Layer Convergence Procedure (PLCP). This part of ISO/IEC 8802 will define a PLCP for each transmission system that is supported.

1.2 Applicability

DQDB subnetworks are intended to support a range of service types including statistical data traffic. Subnetworks should therefore be dimensioned by the network provider so as to ensure that periods of expected peak demands do not cause excessive delay to the particular services that are intended to be supported.

During the periods of overload, which may still occur, a DQDB subnetwork will continue to operate at close to maximum throughput, but some service degradation must necessarily occur. To ensure manageable degradation of the services supported by a DQDB subnetwork, three mechanisms have been provided:

- Pre-Arbitrated access
- Queued Arbitrated access with priority
- Bandwidth balancing

A description of these mechanisms is provided in 2.1.2, and their areas of application are outlined in 2.1.4.

1.3 Normative references

The following standards contain provisions which, through reference in this text, constitute provisions of this part of ISO/IEC 8802. At the time of publication, the editions indicated were valid. All standards are subject to revision, and parties to agreements based on this part of ISO/IEC 8802 are encouraged to investigate the possibility of applying the most recent editions of the standards listed below. Members of IEC and ISO maintain registers of currently valid International Standards.

ANSI T1.102-1987, American National Standard for Telecommunication—Digital Hierarchy—Electrical Interfaces.³

ANSI T1.107-1988, American National Standard for Telecommunications—Digital Hierarchy—Formats Specifications.

ANSI T1.107a-1990, American National Standard for Telecommunications—Digital Hierarchy—Supplement to Formats Specifications (DS3 Format Applications).

CCITT Recommendation E.164 (1988), Numbering plan for the ISDN era. In vol. II, fascicle II.2 of the *CCITT Blue Books*—Telephone network and ISDN—Operation, numbering, routing and mobile service.⁴

CCITT Recommendation G.703 (1988), Physical/Electrical characteristics of hierarchical digital interfaces. In vol. III, fascicle III.4 of the *CCITT Blue Books*—General aspects of digital transmission systems; terminal equipments.

CCITT Recommendation G.704 (1988), Synchronous frame structures used at primary and secondary hierarchical levels. In vol. III, fascicle III.4 of the *CCITT Blue Books*—General aspects of digital transmission systems; terminal equipments.

³ ANSI publications are available from the Sales Department, American National Standards Institute, 11 West 42nd Street, 13th Floor, New York, NY 10036, USA.

⁴ CCITT Recommendations are available from the Telecommunication Standardization Bureau of the International Telecommunication Union, Place des Nations, CH-1211, Geneva 20, Switzerland.

CCITT Recommendation G.707 (1988), Synchronous digital hierarchy bit rates. In vol. III, fascicle III.4 of the *CCITT Blue Books*—General aspects of digital transmission systems; terminal equipments.

CCITT Recommendation G.708 (1988), Network node interface for the synchronous digital hierarchy. In vol. III, fascicle III.4 of the *CCITT Blue Books*—General aspects of digital transmission systems; terminal equipments.

CCITT Recommendation G.709 (1988), Synchronous multiplexing structure. In vol. III, fascicle III.4 of the *CCITT Blue Books*—General aspects of digital transmission systems; terminal equipments.

CCITT Recommendation G.751 (1988), Digital multiplexing equipments operating at the third order bit rate of 34 368 kbit/s and the fourth order bit rate of 139 264 kbit/s and the fourth order bit rate of 139 264 kbit/s and using positive justification. In vol. III, fascicle III.4 of the *CCITT Blue Books*—General aspects of digital transmission systems; terminal equipments.

CCITT Recommendation G.783 (1990), Characteristics of Synchronous Digital Hierarchy (SDH) multiplexing equipment functional blocks.

CCITT Recommendation I.432 (1990), B-ISDN User-Network Interface—physical layer specification.

ISO 2382-25 : 1992, Information technology—Vocabulary—Part 25: Local area networks.⁵

ISO DIS 2382-26...,⁶ Information technology—Vocabulary—Part 26: OSI architecture.

ISO 7498 : 1984, Information processing systems—Open Systems Interconnection—Basic Reference Model.

ISO/TR 8509 : 1987, Information processing systems—Open Systems Interconnection—Service conventions.

ISO 8802-2 : 1989 [ANSI/IEEE Std 802.2-1989], Information processing systems—Local area networks—Part 2: Logical link control.

ISO/IEC 9595 : 1990, Information processing systems—Open Systems Interconnection—Common management information service definition.

ISO/IEC 9596 : 1990, Information processing systems—Open Systems Interconnection—Common management information protocol specification.

ISO/IEC 10039 : 1991, Information technology—Open Systems Interconnection—Local area networks—Medium Access Control (MAC) service definition.

⁵ISO and ISO/IEC publications are available from the ISO Central Secretariat, Case Postale 56, 1 rue de Varembe, CH-1211, Genève 20, Switzerland/Suisse. ISO publications are also available in the United States from the Sales Department, American National Standards Institute, 11 West 42nd Street, 13th Floor, New York, NY 10036, USA.

⁶Current at state of Draft International Standard.

1.4 Definitions

This clause defines the terms used in this part of ISO/IEC 8802. Words in *italics* indicate terms that are defined elsewhere in the list of definitions.

1.4.1 Access Control Field (ACF): The *protocol control information* in a *slot*, which is used to support the *access control function*.

1.4.2 access control function: The generic name for the *Queued Arbitrated (QA) Access* and *Pre-Arbitrated (PA) Access functions* in the *DQDB Layer* that control access to the medium in this part of ISO/IEC 8802.

1.4.3 access unit (AU): The functional unit in a *node* that performs the *DQDB Layer* functions to control access to both *buses*. Access units attach to each bus via a *write* connection and a *read* tap placed *upstream* of the write connection.

1.4.4 address: An identifier that tells where a *service access point (SAP)* may be found (ISO 7498).

1.4.5 address field: The part of a *protocol data unit (PDU)* that contains an *address* that identifies one or more addressable entities. (The address may be a single-source address, single-destination address, or multiple-destination address [*multicast*].)

1.4.6 bandwidth balancing mechanism: A procedure to facilitate effective sharing of the bandwidth, whereby a *node* occasionally skips the use of *empty Queued Arbitrated (QA) slots*.

1.4.7 bridge: A functional unit that interconnects two subnetworks that use a single *Logical Link Control (LLC) procedure* but may use different *Medium Access Control (MAC) procedures*. *Local area networks (LANs)* and *metropolitan area networks (MANs)* are examples of the subnetworks that a bridge may interconnect.

1.4.8 broadcast address: A predefined destination *address* that denotes the set of all *service access points (SAPs)* within a given *layer*.

1.4.9 bus: The concatenation of the *transmission links* between *nodes* and the data path within nodes that provides unidirectional transport of the digital bit stream from the *Head of Bus function* past the *access unit (AU)* of each node to the end of bus.

NOTE—This differs from the bidirectional bus as used in ISO/IEC 8802-3 (Carrier Sense Multiple Access with Collision Detection [CSMA/CD]) and ISO/IEC 8802-4 [Token Bus].

1.4.10 busy slot: A *slot* that contains information and is not available for *Queued Arbitrated (QA) access*.

1.4.11 Configuration Control function: The function that ensures that the resources of all *nodes* of a *DQDB subnetwork* are configured into a correct Dual Bus topology. The resources that are managed are the *Head of Bus function*, the *External Timing Source function*, and the *Default Slot Generator function*.

1.4.12 connection: An association established by a *layer* between two or more users of the layer service for the transfer of information (ISO 7498).

1.4.13 convergence function: A function or procedure that provides sufficient additional services to enable a *layer* or *sublayer* to provide the services expected by a particular higher layer user. (For example, the MAC Convergence Function enables the capabilities of the *Queued Arbitrated access function* to be enhanced to provide the *Medium Access Control (MAC) Sublayer* service to the *Logical Link Control Sublayer*.)

1.4.14 Data Link Layer: In Open Systems Interconnection (OSI) architecture, the *layer* that provides services to transfer data over a *transmission link* between open systems.

In the ISO/IEC series of *local area network (LAN)* standards, the Data Link Layer is formed by the *Logical Link Control (LLC) Sublayer* and the *Medium Access Control (MAC) Sublayer*.

Using this part of ISO/IEC 8802, the Data Link Layer is formed by the operation of the LLC Sublayer (ISO 8802-2) over the *Medium Access Control (MAC) Sublayer* service offered by the *DQDB Layer*.

1.4.15 Default Slot Generator function: The function that defines the identity (i.e., Bus A or Bus B) for each *bus* of a *Dual Bus subnetwork*. Additionally, the function that provides *Head of Bus functions* for both Bus A and Bus B in a *looped Dual Bus subnetwork*. (If the looped Dual Bus subnetwork is reconfigured to an *open Dual Bus subnetwork* due to *transmission link* faults, then the Head of Bus functions for either one or both buses are assigned to other nodes for the duration of the fault.⁷)

1.4.16 Derived MAC Protocol Data Unit (DMPDU): The *Protocol Data Units (PDUs)* of a length of 48 octets formed by the addition of *protocol control information* (including *message identifier* and error protection information) to each of the 44-octet *segmentation units* created from the *segmentation* of an *Initial MAC Protocol Data Unit (IMPDU)*. Each DMPDU is carried as the payload of a *Queued Arbitrated (QA) segment*.

1.4.17 Distributed Queue: The *Medium Access Control (MAC) procedure* in this part of ISO/IEC 8802 for *Queued Arbitrated (QA) access*.

1.4.18 downstream: The direction of data flow along a *bus*, i.e., away from the *Head of Bus function*.

1.4.19 DQDB Layer: The *sublayer* in this part of ISO/IEC 8802 that uses the services of the *Physical Layer* to provide the following:

- *Medium Access Control (MAC) Sublayer* service to the *Logical Link Control (LLC) Sublayer*, and
- *isochronous* service, and
- connection-oriented data service.

1.4.20 Dual Bus: A pair of *buses* carrying digital bit streams flowing in opposite directions. One bus is referred to as Bus A and the other bus as Bus B.

1.4.21 empty Queued Arbitrated (QA) slot: A *Queued Arbitrated (QA) slot* that was designated by the *Head of Bus function* as being available for transfer of a *QA segment*, and that does not contain a *QA segment*.

1.4.22 External Timing Source function: The function of providing the primary point of synchronization of the *DQDB subnetwork* to some external timing reference, for example, that provided by a public network operator.

1.4.23 gateway: A functional unit that interconnects a *local area network (LAN)* with another network having different higher layer protocols.

1.4.24 group address: A predefined destination *address* that denotes a set of selected *service access points (SAPs)* from the *Medium Access Control (MAC) Sublayer* service offered by the *DQDB Layer* to the *Logical Link Control (LLC) Sublayer*.

⁷ Note that all *DQDB subnetworks*, whether in an *open* or *looped Dual Bus* configuration, contain exactly one *node* with an active Default Slot Generator function at network initialization. This node is nominated prior to network initialization. In an open Dual Bus subnetwork, the Default Slot Generator function may be activated at any single node having this capability.

1.4.25 Head of Bus function: The function that generates *empty Queued Arbitrated (QA) slots*, *Pre-Arbitrated (PA) slots*, and *management information octets* at the point on each bus where data flow starts. The Head of Bus function also inserts the *virtual channel identifier* in the *PA segment header* of PA slots.

1.4.26 individual address: An *address* that identifies a single source or destination *service access point*.

1.4.27 Initial MAC Protocol Data Unit (IMPDU): A *protocol data unit (PDU)* formed in the *DQDB Layer* by the addition of *protocol control information* (including address information) to a *MAC Service Data Unit* received from the *Logical Link Control (LLC) Sublayer*. The IMPDU is segmented into 44-octet *segmentation units* for transfer in *Derived MAC Protocol Data Units (DMPDUs)*.

1.4.28 island: An operating part of a *DQDB subnetwork* that is isolated from the node containing the *default slot generator function*.

1.4.29 isochronous: The time characteristic of an event or signal recurring at known, periodic time intervals.

1.4.30 isochronous service octet: A single *octet* of data passed *isochronously* between the *DQDB Layer* and the *Isochronous Service User (ISU)*.

1.4.31 Isochronous Service User (ISU): The entity that uses the *isochronous* service provided by the *DQDB Layer* to transfer *isochronous service octets* over an established *isochronous connection*.

1.4.32 layer: A subdivision of the Open Systems Interconnection (OSI) architecture, constituted by *subsystems* of the same rank (ISO 7498).

1.4.33 layer management: Functions related to the administration of a given Open Systems Interconnection (OSI) *layer*. These functions are performed in the layer itself according to the protocol of the layer and partly performed as a subset of *network management* or *systems management*.

1.4.34 Layer Management Entity (LME): The entity in a *layer* that performs local management of a layer. The LME provides information about the layer, effects control over it, and indicates the occurrence of certain events within it.

1.4.35 Layer Management Interface (LMI): The service interface provided by the *Layer Management Entity (LME)* to the *Network Management Process (NMP)*.

1.4.36 local area network (LAN): A non-public data network in which serial transmission is used without store and forward techniques for direct data communication among data stations located on the user's premises.

1.4.37 Logical Link Control (LLC) procedure: In a *local area network (LAN)* or a metropolitan area network (MAN), the part of the protocol that governs the assembling of *Data Link Layer* frames and their exchange between data stations independently of how the *transmission medium* is shared.

1.4.38 Logical Link Control (LLC) Sublayer: In a *local area network (LAN)* or metropolitan area network (MAN), that part of the *Data Link Layer* that supports medium-independent data link functions, and uses the *Medium Access Control (MAC) Sublayer* service to provide services to the *Network Layer*.

1.4.39 looped Dual Bus: A *DQDB subnetwork* with the *Head of Bus functions* for both Bus A and Bus B collocated.

1.4.40 MAC address: An address that identifies a particular *Medium Access Control (MAC) Sublayer service access point (SAP)*.

1.4.41 MAC Service Data Unit (MSDU): The user data unit received in an MA-UNITDATA request for transfer by the *Medium Access Control (MAC) Sublayer*.

1.4.42 management information octets: *DQDB Layer Protocol Data Units (PDUs)* used to carry *DQDB Layer Management Protocol* information between peer *DQDB Layer Management Entities (LMEs)*.

1.4.43 Medium Access Control (MAC) procedure: In a *local area network (LAN)* or metropolitan area network (MAN), the part of the protocol that governs access to the *transmission medium* independently of the physical characteristics of the medium, but taking into account the topological aspects of the *subnetwork*, in order to enable the exchange of data between *nodes*.

The MAC procedures include framing, error protection, and acquiring the right to use the underlying transmission medium.

1.4.44 Medium Access Control (MAC) Sublayer: In a *local area network (LAN)*, the part of the *Data Link Layer* that supports topology-dependent functions and uses the services of the *Physical Layer* to provide service to the *Logical Link Control (LLC) Sublayer* defined in ISO/IEC 10039.

In this part of ISO/IEC 8802, the combined set of functions in the *DQDB Layer* that support the MAC Sublayer service to the *Logical Link Control (LLC) Sublayer*.

1.4.45 message identifier: An identifier used to identify *Derived MAC Protocol Data Units (DMPDUs)* derived from the same *Initial MAC Protocol Data Unit (IMPDU)*.

1.4.46 metropolitan area network (MAN): A network for connecting a group of individual stations and networks [for example, *local area networks (LANS)*] located in the same urban area.

NOTE—A MAN generally operates at a higher speed than the networks interconnected, crosses network administrative boundaries, may be subject to some form of regulation, and supports several access methods.

1.4.47 MID (Message Identifier) page: A set of one *message identifier* value.

1.4.48 multicast address: See *group address*.

1.4.49 multiport bridge: A *bridge* that interconnects two or more *DQDB subnetworks*.

1.4.50 Network Layer: In Open Systems Interconnection (OSI) architecture, the *layer* that provides services to establish a path between open systems with a predictable quality of service.

1.4.51 network management: Within this series of standards, the functions related to the management of *Data Link Layer* and *Physical Layer* resources and their status across an IEEE 802 *local area network (LAN)* or metropolitan area network (MAN).

1.4.52 Network Management Process (NMP): The entity that provides access to *network management* functions on behalf of the user of the network management services. In order to perform this function, NMPs may intercommunicate in a peer-to-peer manner and may use the services of NMPs in other *nodes* via a network management protocol. The NMP at a node is the user of the service provided at the *Layer Management Interface (LMI)*.

1.4.53 node: A device that consists of an *access unit (AU)* and a single point of attachment of the access unit to each *bus* of a *DQDB subnetwork* for the purpose of transmitting and receiving data on that subnetwork. Adjacent nodes are connected by a *transmission link*.

1.4.54 octet: A group of eight adjacent bits.

1.4.55 offset: The *octet* position relative to the start of a *Pre-Arbitrated (PA) segment* used to carry an *isochronous service octet* for a particular *Isochronous Service User (ISU)*.

1.4.56 open Dual Bus: A DQDB *subnetwork* with the *Head of Bus function* for Bus A and the Head of Bus function for Bus B at different *nodes*.

1.4.57 Physical Layer: In this part of ISO/IEC 8802, the subdivision that provides the protocol to allow transfer of *slot octets*, *management information octets*, and *DQDB Layer* timing information over the *transmission link* between *DQDB Layer subsystems* at adjacent *nodes*. The Physical Layer provides the service to the *DQDB Layer*.

1.4.58 Physical Layer Convergence Procedure (PLCP): In this part of ISO/IEC 8802, the part of the *Physical Layer* that supports the transfer of *slot octets*, *management information octets*, and *DQDB Layer* timing information in a manner that adapts the capabilities of the *transmission system* to the service expected by the *DQDB Layer*.

1.4.59 pipelining: The function of forwarding in sequence some or all of the Beginning of Message (BOM) and Continuation of Message (COM) *Derived MAC Protocol Data Units (DMPDUs)* before receipt of the End of Message (EOM) DMPDU.

1.4.60 Pre-Arbitrated (PA) Access function: The *access control function* in this part of ISO/IEC 8802 that uses assigned *offsets* in *Pre-Arbitrated (PA) slots* for the transfer of *isochronous service octets*.

1.4.61 Pre-Arbitrated (PA) segment: A multiuser *segment* transferred using *Pre-Arbitrated Access (PA) functions*. The payload of the *PA segment* contains *isochronous service octets* from zero or more *Isochronous Service Users (ISUs)*.

1.4.62 Pre-Arbitrated (PA) slot: A *slot* that is dedicated by the *Head of Bus function* for transfer of *isochronous service octets* in the payload of a *PA segment*.

1.4.63 protocol control information: Information exchanged between entities to coordinate their joint operation.

1.4.64 Protocol Data Unit (PDU): Information that is delivered as a unit between peer entities of a *local area network (LAN)* or a *metropolitan area network (MAN)* and that contains control information, address information, and may contain user data.

1.4.65 Queued Arbitrated (QA) Access function: The *access control function* in this part of ISO/IEC 8802 that uses the *Distributed Queue* to access *empty Queued Arbitrated (QA) slots* for the transfer of *QA segments*.

1.4.66 Queued Arbitrated (QA) segment: A *segment* transferred using *Queued Arbitrated (QA) Access functions*.

1.4.67 Queued Arbitrated (QA) slot: A *slot* that is used for the transfer of a *QA segment*.

1.4.68 read: The process of an *access unit (AU)* copying bits of a data stream as they pass on the *bus*.

1.4.69 reassembly: The function in the *DQDB Layer* that provides for the reconstruction of an *Initial MAC Protocol Data Unit (IMPDU)*. Reassembly is performed by concatenating the *segmentation units* received in *Derived MAC Protocol Data Units (DMPDUs)*. This is the inverse process to segmentation.

1.4.70 reconfiguration: The process by which the *Configuration Control function* activates and deactivates resources of a DQDB *subnetwork* to take account of a change in the operational status of a cluster, *node*, or *transmission link* in the subnetwork.

1.4.71 router: A functional unit that interconnects two computer networks that use a single *Network Layer* procedure but may use different *Data Link Layer* and *Physical Layer* procedures.

1.4.72 segment: The *protocol data unit (PDU)* of 52 octets transferred between peer *DQDB Layer* entities as the information payload of a *slot*. It contains a *segment header* of 4 octets and a *segment payload* of 48 octets. There are two types of segments: *Pre-Arbitrated (PA) segments* and *Queued Arbitrated (QA) segments*.

1.4.73 segment header: The *protocol control information* in a *segment*.

1.4.74 segment payload: The unit of data carried by a *segment*.

1.4.75 segmentation: The function in the *DQDB Layer* that fragments a variable length *Initial MAC Protocol Data Unit (IMPDU)* into fixed-length *segmentation units* for transfer in *Derived MAC Protocol Data Units (DMPDUs)* (cf., *reassembly*).

1.4.76 segmentation unit: The fixed-length data units of 44 octets formed by the fragmentation of an *Initial MAC Protocol Data Unit (IMPDU)*.

1.4.77 service access point (SAP): The point at which services are provided by one *layer* (or *sublayer*) to the layer (or sublayer) immediately above it (ISO 7498).

1.4.78 service data unit (SDU): Information that is delivered as a unit between peer *service access points (SAPs)*.

1.4.79 service primitive; primitive: An abstract, implementation-independent interaction between a service user and the service provider.

1.4.80 slot: The *protocol data unit (PDU)* of 53 octets used to transfer *segments*. It contains a segment of 52 octets and a 1 octet *Access Control Field (ACF)*. There are two type of slots: *Pre-Arbitrated (PA) slots* and *Queued Arbitrated (QA) slots*.

1.4.81 sublayer: A subdivision of a *layer* in the Open System Interconnection (OSI) reference model.

1.4.82 subnetwork: In this part of ISO/IEC 8802, a functional unit comprised of a single *Dual Bus* pair and those *access units (AUs)* attached to it. Subnetworks are physically formed by connecting adjacent *nodes* with *transmission links*.

1.4.83 subnetwork configuration: The topological arrangement of *nodes* to form a *subnetwork*. In normal operation a DQDB subnetwork can have one of two configurations, *open Dual Bus* or *looped Dual Bus*.

1.4.84 subsystem: An element in a hierarchical division of an open system that interacts directly only with elements in the next higher division or the next lower division of that open system (ISO 7498).

1.4.85 systems management: Functions in the application *layer* related to the management of various Open Systems Interconnection (OSI) resources and their status across all layers of the OSI architecture.

1.4.86 transmission link: The physical unit of a DQDB *subnetwork* that provides the transmission connection between adjacent *nodes*. Each transmission link accommodates both *buses* of the *Dual Bus* pair between the adjacent nodes.

1.4.87 transmission medium: The material on which information signals may be carried; e.g., optical fiber, coaxial cable, and twisted-wire pairs.

1.4.88 transmission system: The interface and *transmission medium* through which peer *Physical Layer* entities transfer bits.

1.4.89 upstream: The direction along a *bus* that is towards the *Head of Bus function*. This is opposite to the direction of data flow along a bus.

1.4.90 Virtual Channel Identifier (VCI): A label that is used to distinguish between the different virtual channels. A virtual channel is a logical association between entities that enables unidirectional transfer of *segments* between the entities. In the context of this part of ISO/IEC 8802, the VCI label can be used to allow a transmitter to distinguish between different outgoing *protocol data units (PDUs)*, and is used to allow a receiver to determine whether to receive an incoming segment as well as to distinguish between incoming PDUs.

1.4.91 write: The process of an *access unit (AU)* sending data downstream on a bus by logically ORing its outgoing data with the data pattern (normally all zeros) arriving from *upstream* on that bus.

1.5 Abbreviations and acronyms

ACF	Access Control Field
ANSI	American National Standards Institute
AU	Access Unit
BAsize	Buffer Allocation size
BCD	Binary Coded Decimal
BEtag	Beginning-End tag
BIF	Bus Identification Field
BOM	Beginning Of Message
BWBM	BandWidth Balancing Machine
BWB_CNTR	BandWidth Balancing CouNTer
BWB_MOD	BandWidth Balancing MODulus
CC	Configuration Control (function)
CC_1	Configuration Control type 1 function
CC_2	Configuration Control type 2 function
CC_D	Default Configuration Control function
CCITT	the International Telegraph and Telephone Consultative Committee
CD	CountDown (counter)
CE	Connection Endpoint
CIB	CRC32 Indicator Bit
COCF	Connection-Oriented Convergence Function
COM	Continuation Of Message
CMIP	Common Management Information Protocol
CMIS	Common Management Information Service
CRC	Cyclic Redundancy Check
CRC32	32-bit Cyclic Redundancy Check
DA	Destination Address
DIS	Draft International Standard
DMPDU	Derived MAC Protocol Data Unit
DQDB	Distributed Queue Dual Bus
DQSM	Distributed Queue State Machine
DSG	Default Slot Generator
DSGS	Default Slot Generator Subfield

EOM	End Of Message
ETS	External Timing Source
ETSS	External Timing Source Subfield
FEBE	Far End Block Error
FERF	Far End Receive Failure
GPSM	Get Page State Machine
HCS	(segment) Header Check Sequence
HEL	Header Extension Length
HOB	Head of Bus
HOBS	Head of Bus Subfield
HOB_A	Head of Bus A
HOB_B	Head of Bus B
ICF	Isochronous Convergence Function
IEC	International Electrotechnical Commission
IMPDU	Initial MAC Protocol Data Unit
ISDN	Integrated Services Digital Network
ISDU	Isochronous Service Data Unit
ISO	International Organization for Standardization
ISU	Isochronous Service User
kbit/s	Kibobits per second
KPSM	Keep Page State Machine
LAN	Local Area Network
LLC	Logical Link Control
LME	Layer Management Entity
LMI	Layer Management Interface
LSS	Link Status Signal
MAC	Medium Access Control
MAN	Metropolitan Area Network
Mbit/s	Megabits per second
MCF	MAC Convergence Function
MCP	MAC Convergence Protocol
MID	Message IDentifier
MPA	MID Page Allocation
MPAF	MID Page Allocation Field
MSAP	MAC Service Access Point
MSDU	MAC Service Data Unit
NMP	Network Management Process
OSI	Open Systems Interconnection
PA	Pre-Arbitrated
PCC	Page Counter Control (subfield)
PCI	Protocol Control Information
PCM	Page Counter Modulus (subfield)
PCSM	Page Counter State Machine
PDU	Protocol Data Unit
Ph-SAP	Physical layer Service Access Point
PI	Protocol Identification
PL	PAD Length
PLCP	Physical Layer Convergence Procedure
PLCSM	Physical Layer Connection State Machine
POH (SDH)	Path Overhead
POI	Path Overhead Identifier (octets)
PR	Page Reservation (subfield)
PSR	Previous Segment Received
QA	Queued Arbitrated

QOS	Quality Of Service
REQ	REQuest
RIT	Reassembly IMPDU Timer
RQ	ReQuest (counter)
RQM	REQ Queue Machine
RSM	Reassembly State Machine
SA	Source Address
SAP	Service Access Point
SDH	Synchronous Digital Hierarchy
SDU	Service Data Unit
SG_1	Slot Generator type 1 function
SG_2	Slot Generator type 2 function
SG_D	Default Slot Generator function
SNCF	SubNetwork Configuration Field
SSM	Single Segment Message
VCI	Virtual Channel Identifier

1.6 Conformance

An implementation may be called conformant with this part of ISO/IEC 8802 if all mandatory functions in this part of ISO/IEC 8802 are implemented according to the respective mandatory specification.

Additional optional features need not be implemented. However, if they are implemented, they should conform with the specification in this part of ISO/IEC 8802, if applicable.

By way of example, support for 48-bit universally administered MAC Service Access Point (MSAP) addresses is mandatory. However, if 60-bit MSAP addresses are implemented to carry numbers allocated according to CCITT Recommendation E.164, the coding shall be as specified in this part of ISO/IEC 8802.

1.7 Notation

1.7.1 Service specification method and notation

This clause covers the method of specification for the services required of the DQDB Layer by the LLC Sublayer, by the ISUs, and by the Connection-Oriented Data Service Users. It also covers the method of specification for the services required of the Physical Layer by the DQDB Layer.

In general, the services of an (N)-layer (or sublayer) are the capabilities that it offers to an (N)-user in the next higher layer (or sublayer). In order to provide its service, a layer (or sublayer) builds its functions on the services it requires from the next lower layer (or sublayer). Figure 1-2 illustrates the layer service model.

Services are specified by describing the information flow between the (N)-user and the (N)-layer (or Sublayer). This information flow is modeled by discrete, instantaneous events, which characterize the provision of a service. Each event consists of passing a service primitive from one layer (or Sublayer) to the other through the (N)-layer (or sublayer) service access point (SAP) associated with an (N)-user. Service primitives convey the information required in providing a particular service. These service primitives are an abstraction in that they specify only the service provided rather than the means by which the service is provided. This definition of service is independent of any particular interface implementation.

Services are specified by describing the service primitives and parameters that characterize each service. A service may have one or more related primitives that constitute the activity that is related to the particular service. Each service primitive may have zero or more parameters that convey the information required to provide the service.

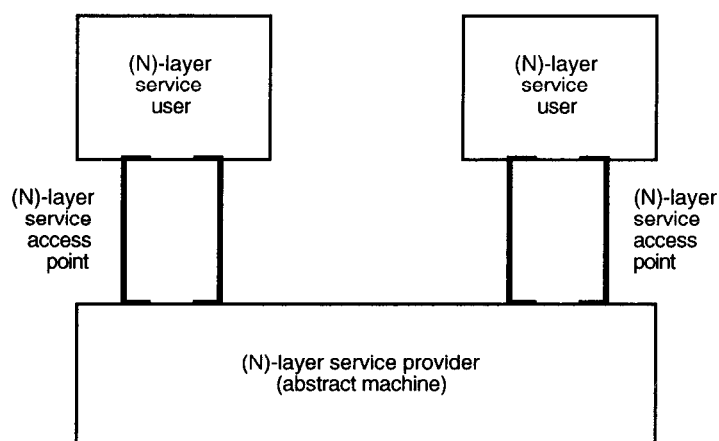


Figure 1-2—Layer service model

Primitives are of four generic types:

- a) **REQUEST.** The request primitive is passed from the (N)-user to the (N)-layer (or sublayer) to request that a service be initiated.
- b) **INDICATION.** The indication primitive is passed from the (N)-layer (or sublayer) to the (N)-user to indicate an internal (N)-layer (or sublayer) event that is significant to the (N)-user. This event may be logically related to a remote service request, or may be caused by an event internal to the (N)-layer (or sublayer).
- c) **RESPONSE.** The response primitive is passed from the (N)-user to the (N)-layer (or sublayer) to complete a procedure previously invoked by an indication primitive.
- d) **CONFIRM.** The confirm primitive is passed from the (N)-layer (or sublayer) to the (N)-user to convey the results of one or more associated previous service request(s).

This part of ISO/IEC 8802 only requires the use of request and indication primitives. Possible relationships among primitive types are illustrated by the time sequence diagrams shown in figure 1-3. The figure also indicates the time relationship of the primitive types. The key of figure 1-3 depicts the time sequence framework used to illustrate the interactions between the two correspondent service users and their associated service provider. Primitive types occurring earlier in time and connected by dotted lines are the logical antecedents of subsequent primitive types. If there is no specific relationship between events, in that it is impossible to predict which will occur first, but both must occur within a finite period of time, then a tilde (~) is used.

This part of ISO/IEC 8802 uses the conventions for naming service primitives contained in annex A of ISO/TR 8509.

1.7.2 State machine notation

The operation of a protocol can be described by subdividing the protocol into a number of interrelated functions. The operation of a function can be described by state machines. Each machine represents the domain of a function and consists of a group of connected, mutually exclusive states. Only one state of a function is active at any given time.

Each state that the function can assume is represented by a vertical line. All permissible transitions between states of a function are represented by arrows from one state to the next state. The condition that causes a change from one state to another is shown above the transition arrow. The affected signals and flags are shown below the transition arrow.

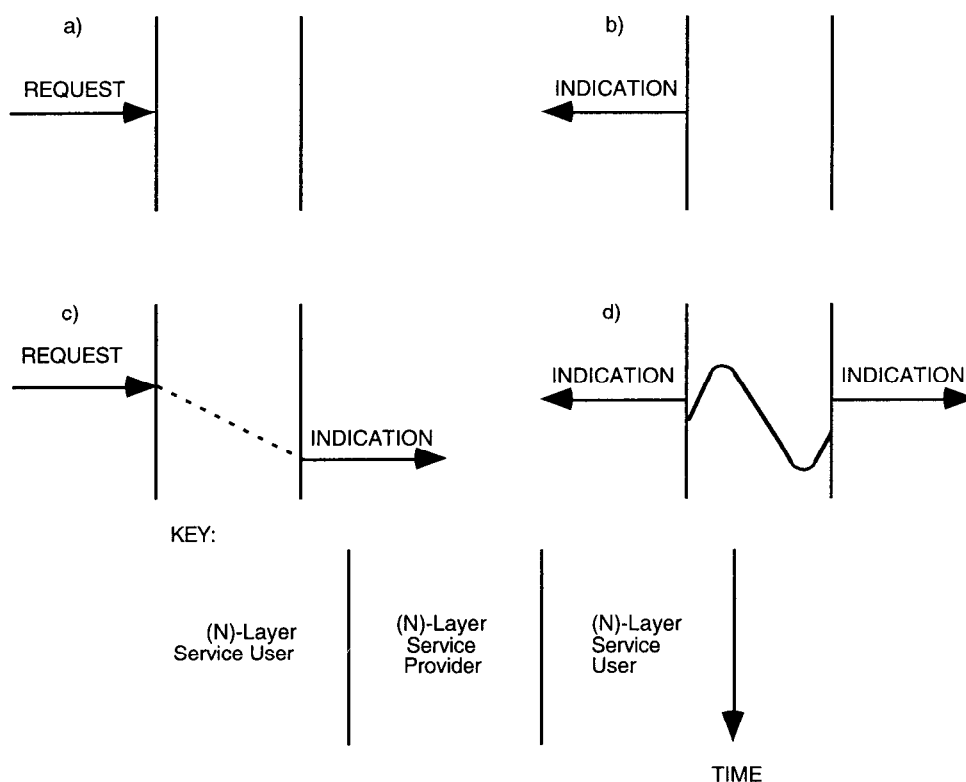


Figure 1-4—Time sequence diagram

Signals and flags refer to the form by which one functional block communicates with another. Both terms arise from the use of the state machines as a descriptive method in this part of ISO/IEC 8802. Signals are transient communications (in theory, they are of infinitesimal time width) that occur during a change of state. Flags are steady, nontransient indications that remain constant until explicitly changed.

1.8 Organization of this part of ISO/IEC 8802

Clause 2 contains an overview of the operation of a DQDB subnetwork and the communication services supported. It also provides an introduction to the functional architecture of a DQDB node.

Clause 3 contains the service definition for the DQDB Layer, which gives a formal definition for the MAC Sublayer service offered to LLC and for the isochronous service, and a statement on the requirements for the connection-oriented data service.

Clause 4 contains the definition of the Physical Layer service to the DQDB Layer.

Clause 5 gives a rigorous description of the internal operation of a DQDB node, with forward references to clauses 6, 7, 8, 9, and 10. Clause 6 describes the DQDB Layer PDU formats (with the exception of the management information octets, which are described in clause 10). Clause 7 describes the facilities required for operation of the DQDB Layer, such as timers, counters, parameters, and status indicators. Clause 8 specifies the state machines that describe peer-to-peer communication by DQDB Layer entities.

DQDB Layer Management is described in clauses 9 and 10. Clause 9 describes the services offered at the DQDB LMI to the NMP. This interface can support remote management of a DQDB node using management protocols such as those specified in ISO 9595 (CMIS) and ISO 9596 (CMIP). Clause 10 describes the

operation of the DQDB Layer Management Protocol, which is the mechanism by which DQDB Layer Management Entities communicate to perform layer-wide management of DQDB Layer resources.

Clause 11 describes the principles of operation of the Physical Layer. Clause 13 describes the operation of the Physical Layer Convergence Procedure (PLCP) required to support a DS3 rate (as specified in ANSI T1.102 and T1.107). Subsequent clauses describe the operation of the PLCs required to support transmission systems based on CCITT Recommendation G.703 and operating at 2.048 Mbit/s (clause 14), 34.368 Mbit/s (clause 15), and 139.264 Mbit/s (clause 16). Clause 17 describes the operation of the PLCP required to support SDH transmission systems operating at 155.520 Mbit/s, as specified in CCITT Recommendations G.707, G.708, and G.709.

2. Overview

This clause is intended as an introduction to the DQDB access mechanism and is not intended as a rigorous definition of the protocol. The first part of this clause, 2.1, provides an overview of the operation of a DQDB subnetwork. The second part, 2.2, provides a focused view describing the functional architecture of a DQDB node. Rigorous descriptions of the node functional operation, and of the DQDB Layer protocols, are given in clause 5, and clauses 8 and 10, respectively.

2.1 DQDB subnetwork

The purpose of a MAN is to provide integrated services, such as data, voice, and video, over a large geographical area. This part of ISO/IEC 8802 describes the DQDB subnetwork that can be used as a component part of a MAN. Typically, a MAN would consist of interconnected DQDB subnetworks. The interconnection of subnetworks could be via bridges, routers, or gateways, but this is not covered in this International Standard. MANs can be interconnected by suitable means to form a wide area network.

The DQDB subnetwork is a distributed multiaccess network that supports integrated communications. In particular, the DQDB subnetwork supports connectionless data transfer, connection-oriented data transfer, and isochronous communications (e.g., voice). The various communication transfer functions share flexibly the total capacity of the subnetwork.

In the first part of this clause, 2.1.1, there is a description of the Dual Bus architecture and the format by which nodes communicate. The second part of the clause, 2.1.2, describes the control of access to the Dual Bus. This includes Distributed Queue access control for non-isochronous communications and Pre-Arbitrated access control for isochronous communications.

The third part of this clause, 2.1.3, describes the functions specified in addition to the access control mechanism to support the DQDB Layer services. For example, it describes how LLC is supported by the DQDB subnetwork.

The fourth part of this clause, 2.1.4, provides guidance to network administrators in the application of the standard.

The final part of this clause, 2.1.5, describes the Configuration Control function, which is important in subnetwork start-up and in subnetwork reconfiguration in the case of bus faults.

2.1.1 Dual Bus architecture

The Dual Bus architecture of a DQDB subnetwork consists of two unidirectional buses and a multiplicity of nodes along the buses, as shown in figure 2-1. The buses, denoted Bus A and Bus B, support communications in opposite directions allowing full duplex communications between any pair of nodes on the subnetwork. Therefore, it is desirable for a node to know which bus to use to communicate to another node. Annex C provides information as to how this can be done for connectionless services.

As both buses are operational at all times, the capacity of the subnetwork is up to twice the capacity of a single bus. The operation of the two buses in the transfer of data is independent.

The data on each bus are formatted into either DQDB Layer Management information octets or fixed-length slots, this format being generated by the node at the head of each bus. Hence, there must be one node at the Head of Bus A and one node at the Head of Bus B. The management information octets are used to maintain the operational integrity of the subnetwork. The slots are used to carry data between the nodes. That is, nodes may write into slots under the control of the access protocol. All data flow terminates at the end of the bus.

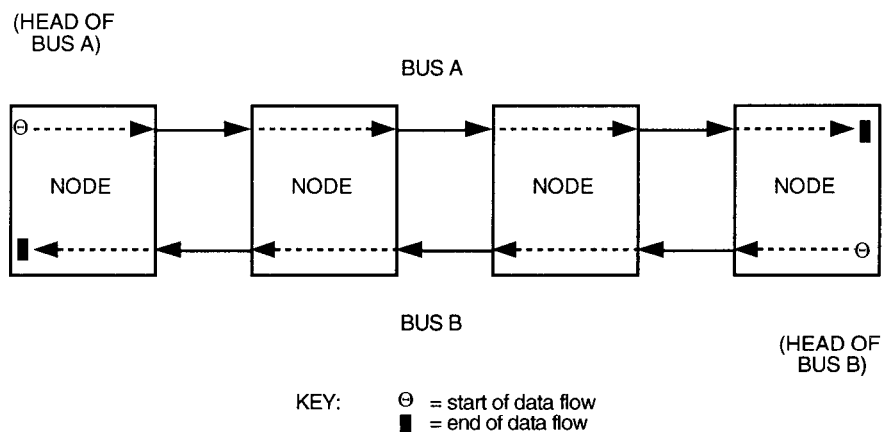


Figure 2-1—Dual Bus architecture

The nodes on the subnetwork consist of an access unit and the attachment of the access unit to the two buses. The access unit performs the node's DQDB Layer functions and attaches to a bus via one read and one write connection. The writing of data to the bus is a logical OR of the data from upstream with the data from the access unit. The read connection is placed logically ahead of the write connection and allows all data to be copied from the bus unaffected by the unit's own writing.

The read and write functions can be implemented using either passive or active techniques.⁸ An example implementation of the AU connection is shown in figure 2-2. With this connection arrangement, an AU can copy data that passes on the bus, but can never remove it, and only alter data when it is permitted by the access protocol.

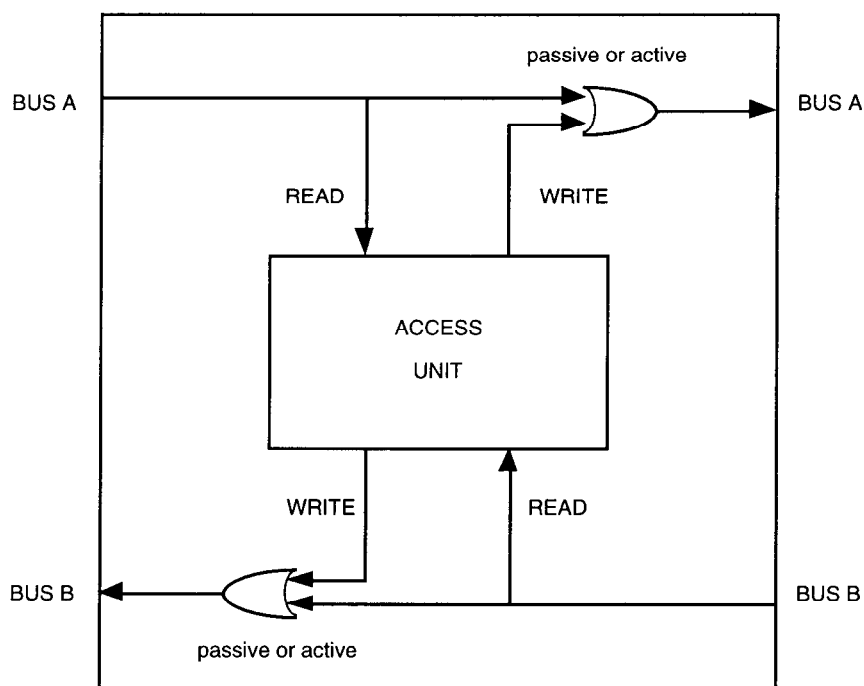


Figure 2-2—Example access unit attachment

⁸ A passive optical implementation would limit the number of nodes on the bus according to the optical power budget.

With the exception of AUs performing the Head of Bus functions, the operation of the bus is only loosely coupled to the operation of each AU; an AU can fail (provided that it does not destructively write to the bus), or even be removed from the subnetwork, with no consequence on the operation of the rest of the subnetwork.

Under normal conditions there is a single source of slot timing for the subnetwork, which is used to ensure that all nodes in the subnetwork transfer data (i.e., slots and management information octets) at the same average rate. This is necessary to ensure stable operation of the Distributed Queue access mechanism (see annex D) and to ensure that isochronous communications can be supported without slips being introduced by different timing sources at the two communicating entities.

The timing source can be traceable to an external timing source, such as that provided by a public telecommunications network. Any node or nodes can be designated by network operations procedures to provide the External Timing Source function. If the subnetwork is supporting certain isochronous services and is connected to a public network, then the external timing source may be required. However, in the absence of an external timing source, one node with the capability specified below will be selected by subnetwork configuration control procedures to be the source of subnetwork timing. In this case, the selected node must have available and operate a local clock that has a nominal period of 125 μ s.

2.1.2 Access control to the Dual Bus subnetwork

The DQDB Layer protocols and protocol data units allow a DQDB subnetwork to be operated at a variety of speed and distance combinations. This permits a wide range of transmission systems to be used in the underlying Physical Layer.

The DQDB Layer provides two modes of access control to the Dual Bus. These are Queued Arbitrated (QA) and Pre-Arbitrated (PA), which use QA and PA slots for access respectively. Each slot contains an ACF and a segment, which forms the payload of the slot.

Queued Arbitrated access is controlled by the Distributed Queueing protocol and would be used typically to provide non-isochronous services. Pre-Arbitrated access would be used typically to provide isochronous services. The Distributed Queue and Pre-Arbitrated access protocols are described in the following subclauses.

2.1.2.1 The Distributed Queue access protocol

Distributed Queueing is a media access protocol that controls the access to the payload of QA slots on the DQDB bus. The fixed-length payload of a QA slot is called a QA segment.

The protocol provides deterministic access control for services that are typically bursty in nature. In particular, the protocol can enable all of the payload bandwidth to be used and the average slot access delay approximates that of a perfect scheduler up to high levels of subnetwork loading.

The operation of the protocol is based on two control fields: the BUSY bit, which indicates whether or not a slot is used, and the REQUEST field, which is used to indicate when a segment has been queued for access. Each node, by counting the number of requests it receives and unused slots that pass, can determine the number of segments queued (i.e., in line) ahead of it. This counting operation establishes a single queue across the subnetwork of segments queued for access to each bus.

With such queued access, levels of priority can be established by operating a number of queues, one for each level. Within each level the performance characteristics described above are maintained. Segments will gain access as soon as capacity becomes available, but priority is always given to segments in higher level queues. This part of ISO/IEC 8802 specifies three levels of priority.

The operation of Distributed Queueing is fundamentally different from other MAC protocols. In Distributed Queueing, information that explicitly indicates the queueing state of the subnetwork is kept in the nodes. Hence, when a node has data to transmit, the node need not first derive information from the subnetwork to tell it when it can gain access.

With Distributed Queueing, a current state record that holds the number of segments awaiting access to the bus is kept in every node. When a node has a segment for transmission, it uses this count to determine its position in the distributed queue. If no segments are waiting, permission for access is immediate, otherwise deference is given only to those segments already queued.

To facilitate effective sharing of the bandwidth, the Distributed Queue protocol includes a bandwidth balancing mechanism to occasionally skip the use of empty QA slots.

2.1.2.1.1 The basic Distributed Queueing algorithm

The operation of the basic Distributed Queueing algorithm for access to Bus A is illustrated in figure 2-3. In this case, Bus A is the forward bus and Bus B is the reverse bus. An identical, but independent, arrangement applies for access to the opposite bus, Bus B.

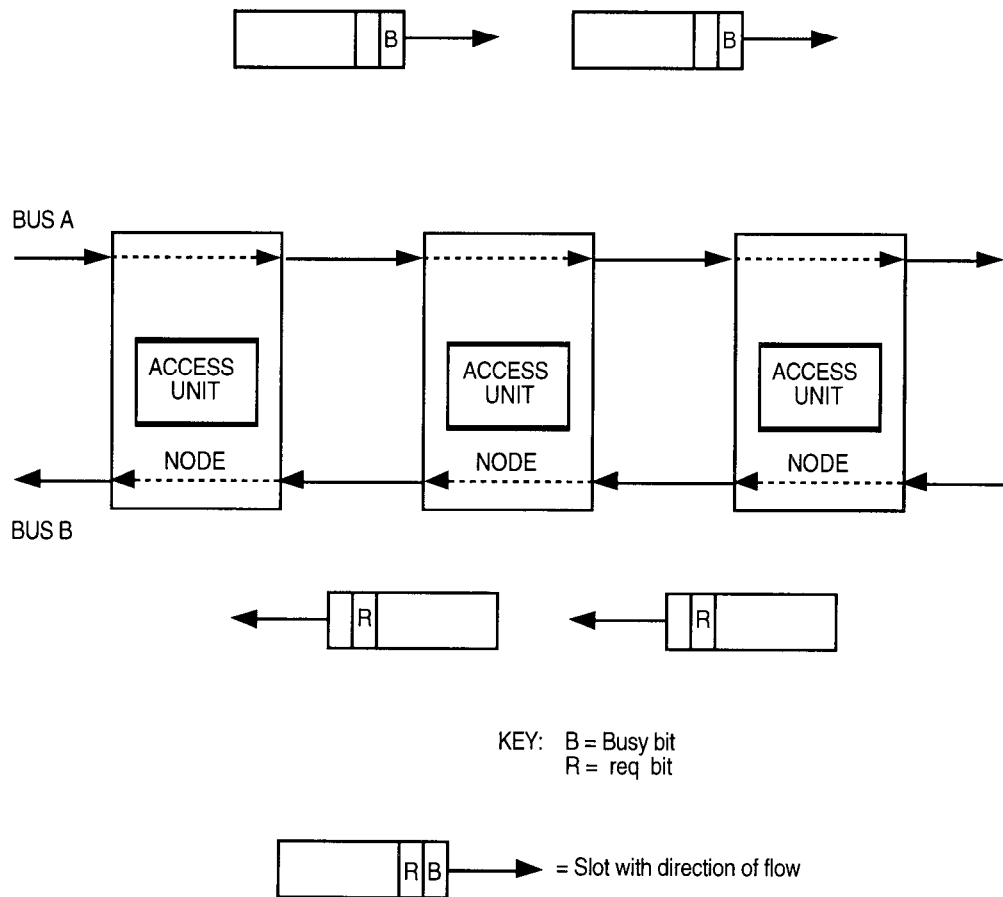


Figure 2-3—Queue formation on Bus A

Each slot on the buses (whether a QA or PA slot) contains an ACF that includes a BUSY bit and a REQUEST field of three Request bits (henceforth referred to as REQ bits), one for each priority level. The

BUSY bit indicates whether or not the slot is used. The REQ bits are to signal when a QA segment has been queued on the reverse bus.

The operation of the Distributed Queueing algorithm within a single priority level is described first. The description of the multiple priority scheme follows, as does a description of the bandwidth balancing mechanism.

When an AU has a QA segment for transmission on the forward bus it will cause a single REQ to be sent on the reverse bus. This REQ will eventually be written into the next free REQ bit of the required priority on the reverse bus. The bit once written will pass to all upstream AUs, where upstream is defined in relation to data flow on the forward bus, in this case Bus A. This REQ bit serves as an indicator to the upstream AUs that an additional QA segment is now queued for access. For each AU, the Distributed Queueing algorithm allows, at most, one QA segment per priority level, to be queued for access to each bus.

Each AU keeps track of the number of QA segments queued downstream from itself for access to the forward bus by counting the REQ bits as they pass on the reverse bus, as shown in figure 2-4a). The request (RQ) counter is incremented for each REQ passing on the reverse bus. For a node that is not queued to send, one REQ in the RQ counter is cancelled each time an empty slot passes on the forward bus. This is done since the empty slot that passes the AU will be used by one of the downstream queued QA segments. Hence, with these two actions, the RQ counter keeps a record of the number of segments queued downstream.

In addition to issuing the REQ for the reverse bus, an AU with a QA segment to send will transfer the current value of the RQ counter to another counter, the countdown (CD) counter, as shown in figure 2-4b), with the RQ counter then being reset to zero. This action loads the CD counter with the number of downstream segments queued ahead of it. This, along with the issuing of the REQ for the AU's segment, effectively places the QA segment in the distributed queue. The distributed queue at a given level of priority approximates a first-in-first-out (FIFO) queue of the QA segments at the heads of the local queues in each node.

To ensure that the segments registered in the CD counter gain access before the newly queued segment in the given AU, the CD counter is decremented for every empty slot that passes on the forward bus. This operation is shown in figure 2-4b). The given AU can transmit its QA segment in an empty slot provided the CD count is zero. For this single priority description, this is equivalent to claiming the first free slot after the CD count reaches zero, which ensures that no downstream segment that queued after the given segment can access out of order.

During the time that the AU is waiting for access for its segment, any new REQs received from the reverse bus are added to the RQ counter, as shown in figure 2-4b). Hence, the RQ counter still tracks the number of segments newly queued downstream and the count will be correct for the next QA segment access.

It should be noted that the control of the access of the QA segment to the forward bus is determined solely by the values of the counters. Access is not inhibited if the value of the CD counter is zero but the REQ bit associated with the QA segment has not yet been written to the reverse bus. That is, the operation of writing REQ bits and sending QA segments is independent.

Thus, with the use of two counters in each AU, one counting unsatisfied access requests and the other counting down before access, a FIFO queue is established for access to the forward bus. The queue formation is also such that, except for bandwidth balancing considerations (described later in this clause), a slot is never wasted on the subnetwork if there is a segment queued for it. This is guaranteed since the CD count in AUs with segments queued represents the number of downstream segments ahead in the queue. Since at any point in time⁹ one segment must have queued first, then at least one AU is guaranteed to have a CD count of zero. It is the most upstream of these that will gain access.

⁹ More correctly, at the same point in time normalized for propagation delay of the slot along the bus.

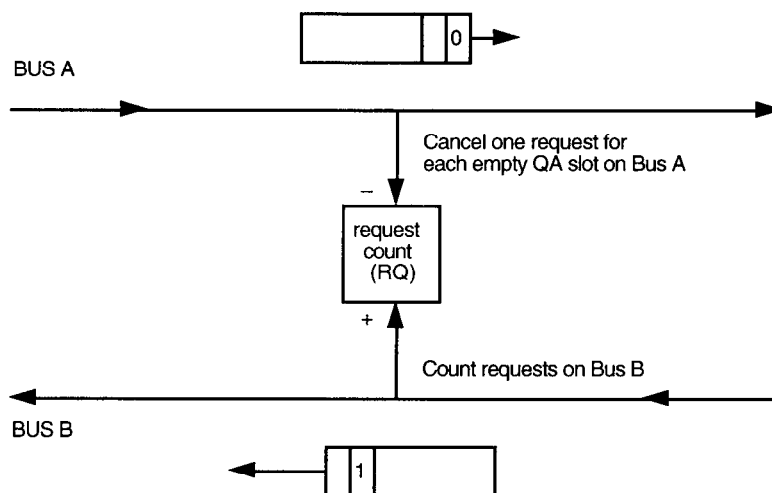


Figure 2-4a)—Node not queued to send (Bus A)

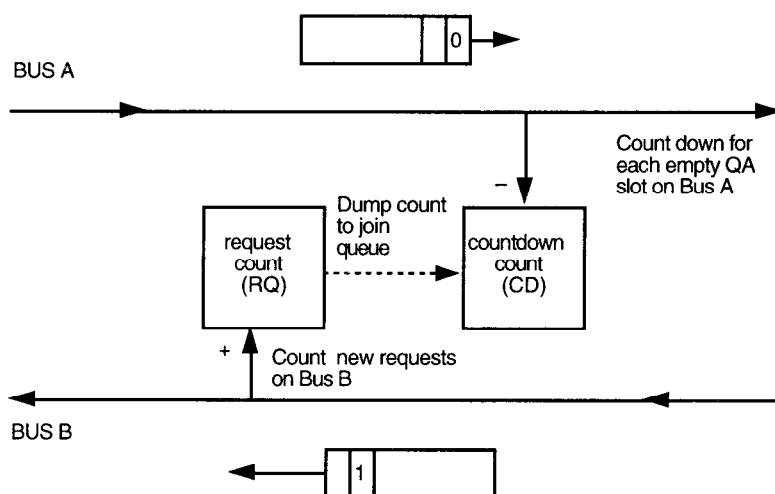


Figure 2-4b)—Node queued to send (Bus A)

As an example of the operation of Distributed Queueing, consider figure 2-5, which shows a subnetwork consisting of five AUs, three of which will queue for access. Assume that each request counter, the value of which is indicated by RQ at the AU, and each CD counter, the value of which is indicated by CD at the AU, are all initially zero and that all slots passing on Bus A are used. AUs 5, 2, and 3 queue in that order, using the operations described above. The counters in the final state can be interpreted as follows:

- AU 1: There are three AUs downstream on Bus A that have QA segments queued for access.
- AU 2: There is one downstream AU with a QA segment queued for access before AU 2, as represented by the one in the CD counter. There is also one downstream AU with a QA segment queued for access after AU 2, as represented by the one in the RQ counter.
- AU 3: There is one downstream AU with a QA segment queued for access before AU 3, as represented by the one in the CD counter, and no downstream AU with a QA segment queued after AU 3, as represented by the zero in the RQ counter. (Note that the REQ from AU 2 is not registered at AU 3.)

AU 4: There is one downstream AU with a QA segment queued for access.

AU 5: There are no downstream AUs from AU 5; hence, both the RQ and CD counts are zero.

Now consider figure 2-6, which shows the AUs gaining access. Assume no further requests are received on the reverse bus and that empty QA slots pass on bus A. The AUs then gain access in the order 5, 2, then 3, which is the order of queueing. Note that after AU 5 gains access both AU 2 and AU 3 have a CD count of zero. AU 2 will gain access first as it is the most upstream.

2.1.2.1.2 Priority Distributed Queueing

The Distributed Queueing protocol supports the assignment of priority to QA segment access. However, all connectionless data segments must be sent at the lowest priority, priority level 0 (see 7.3.3). Future use of priorities for other services is under study. For upward compatibility, Distributed Queues for each level of priority are required in this part of ISO/IEC 8802.

The operation of separate distributed queues for each level of priority is achieved by using a separate REQ bit on the reverse bus for each level of priority and separate RQ and CD counters for each priority level. The counters operate similarly to the single priority case, except that account must be taken of REQs at higher levels.

That is, for an AU that does not have a QA segment queued at a particular priority level, the RQ counter operating at that level will count REQs at the same *and higher* priority levels. Thus, the RQ counter records all queued segments at equal and higher priorities.

If the AU does have a QA segment queued at a particular priority level, then the RQ counter operating at that level will only count REQs at the same priority level. However, the operation of the CD counter at that priority level is slightly altered. The CD counter will, in addition to counting down the received empty slots, increment for REQs received at higher priority levels. This allows the higher priority segments to claim access ahead of already queued segments.

It should be noted that decrementing of the RQ counters and CD counters at all priority levels occurs when an empty QA slot is received at the function within the node operating the Distributed Queue, not when it is sent by this function. This ensures that the correct counter values are maintained if the highest priority segment queued by the AU gains access to the empty slot and the AU marks the slot as busy.

The intent of the priority scheme is for the access performance of the highest priority traffic to be unaffected by lower priority traffic. Such a feature is very important in network signaling. However, it should be noted that priority should not be relied upon for all subnetwork distance, speed, and loading conditions. (See 2.1.4.)

This part of ISO/IEC 8802 specifies three levels of access priority. The definition of how these levels are used is a matter for the operator of the particular DQDB subnetwork, although it is likely that the highest level would be used for network management and signaling. The DQDB Layer assumes that the user of the QA access requests the appropriate priority level for a particular segment transfer.

2.1.2.1.3 Bandwidth balancing

To facilitate effective sharing of the bandwidth, the Distributed Queue access protocol includes a bandwidth balancing mechanism to occasionally skip the use of empty QA slots. To support this function, the Distributed Queueing function in the node operates a bandwidth balancing counter for each bus. There is also a system parameter called the Bandwidth Balancing Modulus, denoted BWB_MOD. If BWB_MOD is set to zero, the bandwidth balancing mechanism is disabled. If BWB_MOD is set to a value greater than zero, it indicates when the value of the bandwidth balancing counter should roll over to zero, as explained below.

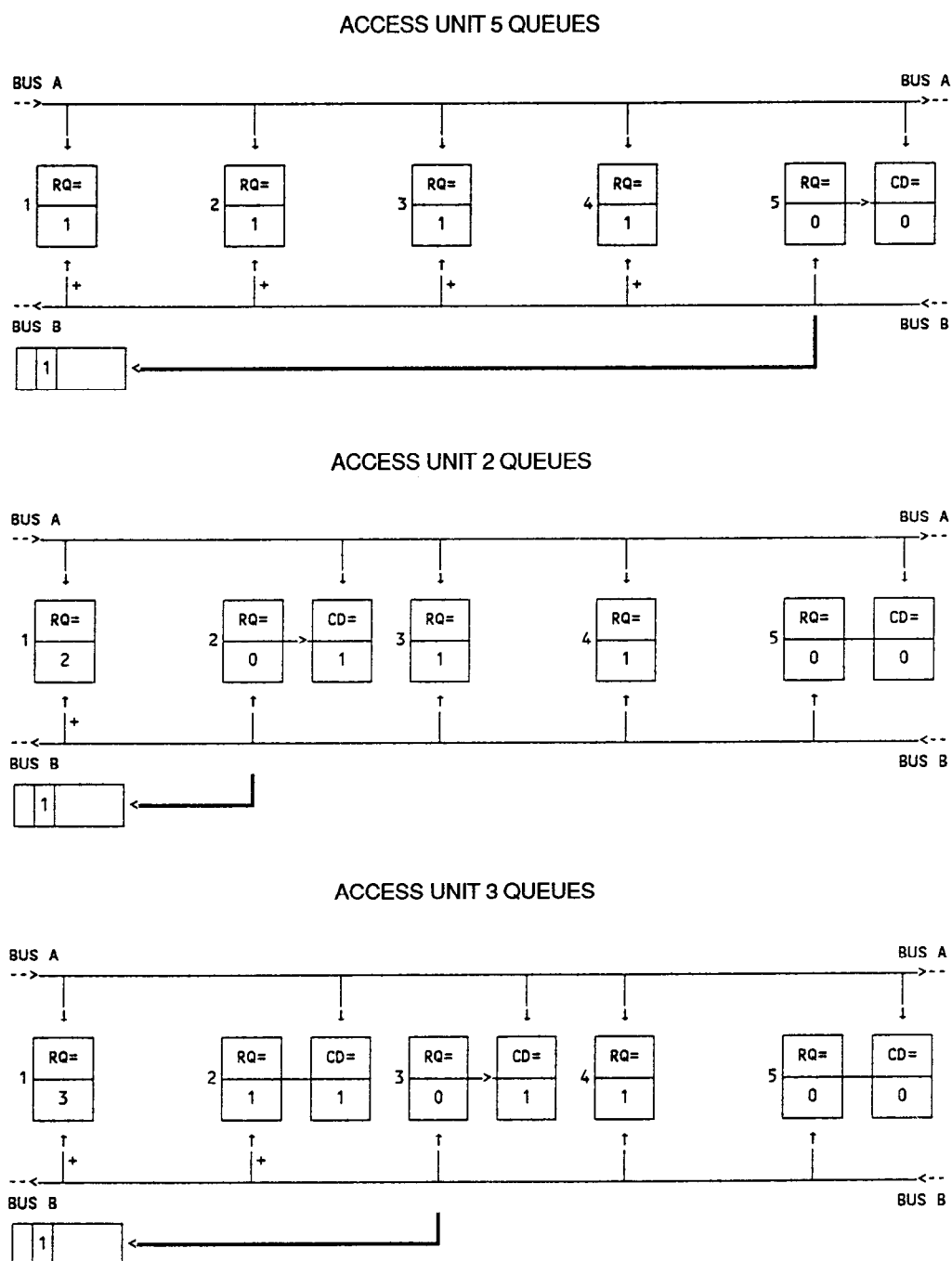
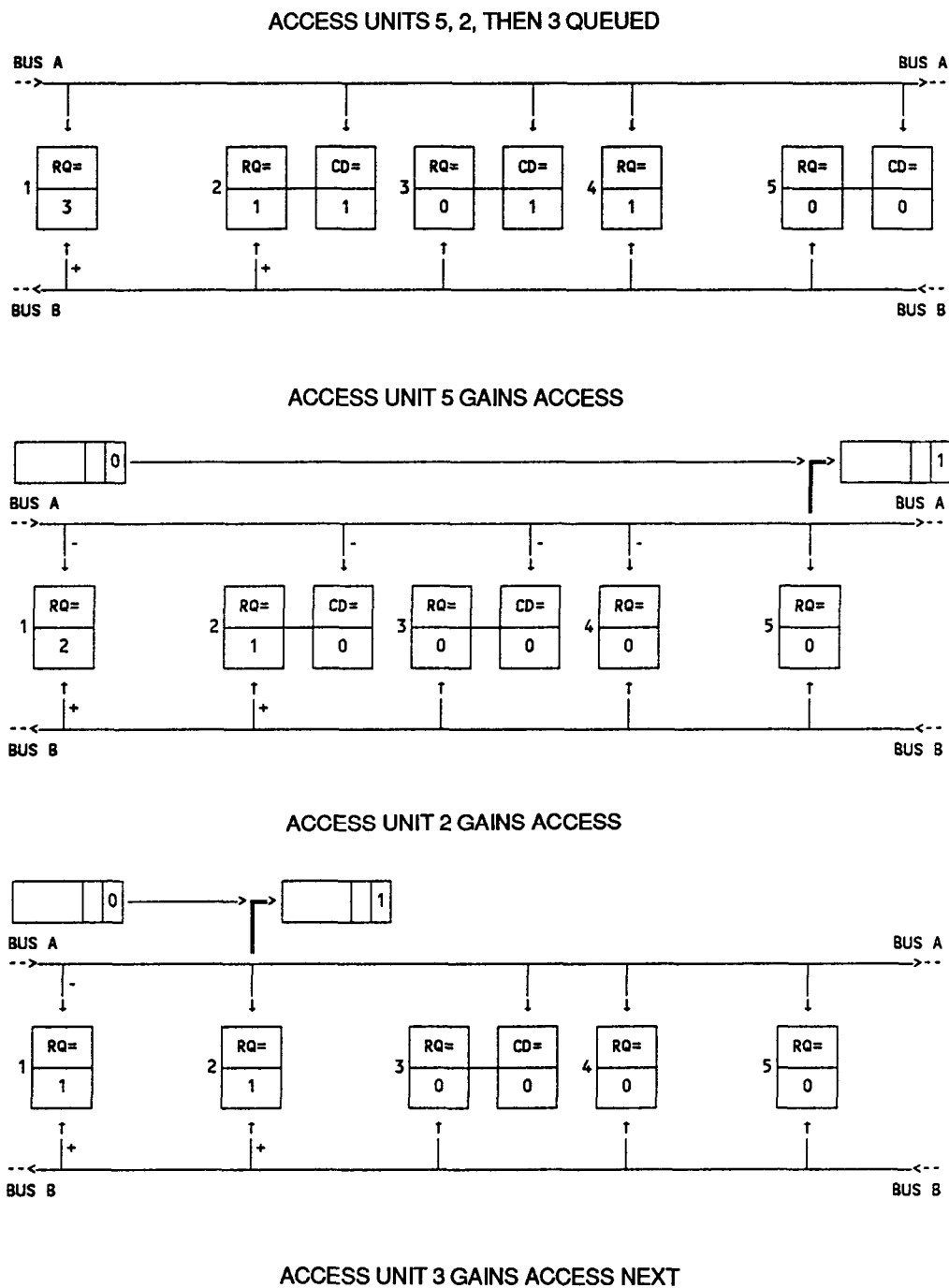


Figure 2-5—Queueing for access (Bus A)



Whenever the node transmits a QA segment on one of the buses, the bandwidth balancing counter for that bus is incremented by one, provided that the counter does not have a value of $(BWB_MOD - 1)$. If the counter does have a value of $(BWB_MOD - 1)$ when the QA segment is transmitted, the bandwidth balancing counter is reset to zero instead.

When the bandwidth balancing counter for a bus is reset to zero, this is an indication that the node should skip the use of one empty QA slot on that bus. To make that happen, the Distributed Queueing function increments the request counters for that bus for all priority levels for which no QA segment is queued, and increments the countdown counters for that bus for all priority levels for which a QA segment is queued.

NOTES

1—Resetting the bandwidth balancing counter for a bus has no effect on the queue of requests waiting to be sent on the opposite bus.

2—If the value of BWB_MOD is one, the bandwidth balancing counter always has the value zero. This means that every time the node transmits a QA segment on the bus, the request or countdown counter corresponding to each priority level for that bus is incremented, as described above.

2.1.2.2 Pre-Arbitrated access control

Pre-Arbitrated (PA) slot access will be used typically to provide for transfer of isochronous service octets. The access to PA slots and the use of PA segment payloads differ from that for QA access. The access differs in that PA slots are designated by the node at the head of bus and that more than one AU may share access to the slot. A PA segment payload consists of a number of octets, each of which can be used by a different AU. Therefore, an AU may write zero, one or more isochronous service octets into designated positions of a PA segment payload. The AU is notified of the offsets of these octet positions relative to the start of the PA segment payload via the DQDB Layer Management procedures.

The designation of slots to PA access and the marking of the PA slot and PA segment header is controlled by functions in the node at the head of the bus. In particular, the Head of Bus function must write the Virtual Channel Identifier (VCI) field into PA slots. The Head of Bus function must also ensure that the PA slots for each VCI value are provided in a periodic manner on the bus to guarantee that sufficient bandwidth is available for the Isochronous Service Users (ISUs).

The access to the PA slots by an AU commences by examining the VCI. For each VCI value that the AU must access, the AU will have a table that indicates which octet offsets within the slot the AU should use for reading and writing. The AU will write isochronous service octets into those write positions and will read from positions the table has marked for reading. The PA slot is ignored if the VCI is not one in use by the AU.

2.1.3 Provision of DQDB Layer services

The Queued Arbitrated and Pre-Arbitrated functions of the DQDB Layer provide access control to the Dual Buses. These access control functions are used by a range of convergence functions to create the DQDB Layer services shown in figure 2-7 and defined formally in clause 3. This part of ISO/IEC 8802 specifies the convergence function for the provision of the ISO/IEC 10039 MAC Sublayer service to the LLC Sublayer. Also specified is a general isochronous service. Future services, such as a connection-oriented data service, are the subject of further study.

At no time do two convergence functions simultaneously use the same VCI at a given node. Hence, if a node is to receive a particular segment, the VCI can be used to direct the segment to the appropriate convergence function at the node. An overview of the convergence functions specified in this part of ISO/IEC 8802 is given in the following subclauses.

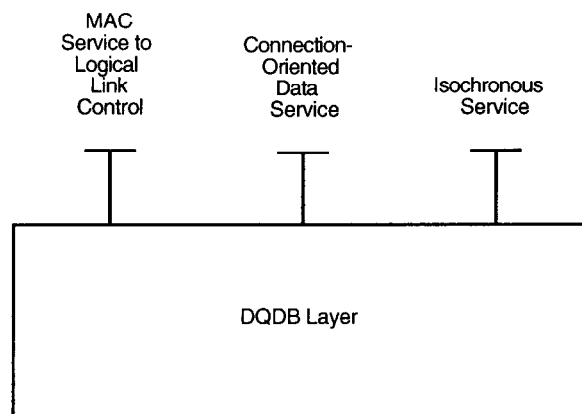


Figure 2-7—DQDB Layer services

2.1.3.1 Provision of MAC service to LLC

The provision of MAC service to LLC consists of the segmentation of the MAC Service Data Unit (MSDU) at the source into fixed-length units and the transfer of these fixed-length units to the destination, which reassembles them into the MSDU.

The segmentation process follows the formation of an Initial MAC Protocol Data Unit (IMPDU) by the addition of an IMPDU header,¹⁰ an optional Header Extension, an optional 32-bit CRC, a Common PDU trailer,¹¹ and a variable-length PAD field to the MSDU. The PAD field ensures that all of the fields added to the MSDU are 32-bit aligned. The IMPDU is fragmented into fixed-length segmentation units, as shown in figure 2-8, for transfer in QA segment payloads. There may be padding of the IMPDU with trailing zero octets to ensure complete filling of the last segmentation unit.

All segment payloads that support the MAC service are called Derived MAC Protocol Data Units (DMPDUs) and consist of a header field and a trailer field¹² along with the segmentation unit, as shown in figure 2-9. The DMPDU header field consists of three subfields. The first is the Segment Type subfield. The second is the Sequence Number subfield. The third subfield is the Message Identifier (MID). The DMPDU trailer field consists of two subfields.

The first is the Payload Length subfield. The second subfield is the Payload CRC.

The MID is used to provide the logical linking between segmentation units derived from the same IMPDU, and should be unique on a subnetwork while the IMPDU is being transferred. Each AU will have at least one unique MID. The allocation of the MID numbers is controlled by the MID page allocation scheme, which is a distributed method for claiming and keeping MID values that are unique across the whole subnetwork.¹³

¹⁰ The IMPDU header is formed by the concatenation of two types of header information. The first four octets of the 24-octet IMPDU header are called the Common PDU header. The remaining 20 octets of the IMPDU header are called the MAC Convergence Protocol (MCP) header. The MCP header is specific to the transfer of a MSDU by the DQDB Layer, and is not used to support the connection-oriented data service.

¹¹ The Common PDU header and the Common PDU trailer are intended to support a common segmentation and reassembly procedure across all of the bursty data services that can be provided by the DQDB Layer, such as the MAC service and the connection-oriented data service. The format of the Common PDU header and the Common PDU trailer is the same for both of these services, although the exact use of the fields differs.

¹² As with the Common PDU header and Common PDU trailer for the IMPDU, the DMPDU header and DMPDU trailer are also intended to support the common segmentation and reassembly procedure across all of the bursty data services that can be provided by the DQDB Layer. Hence, the QA segment payloads used to support all of these services will contain a header and trailer with identical format to those carried in the DMPDU.

¹³ The MID page allocation scheme is part of the DQDB Layer Management protocol.

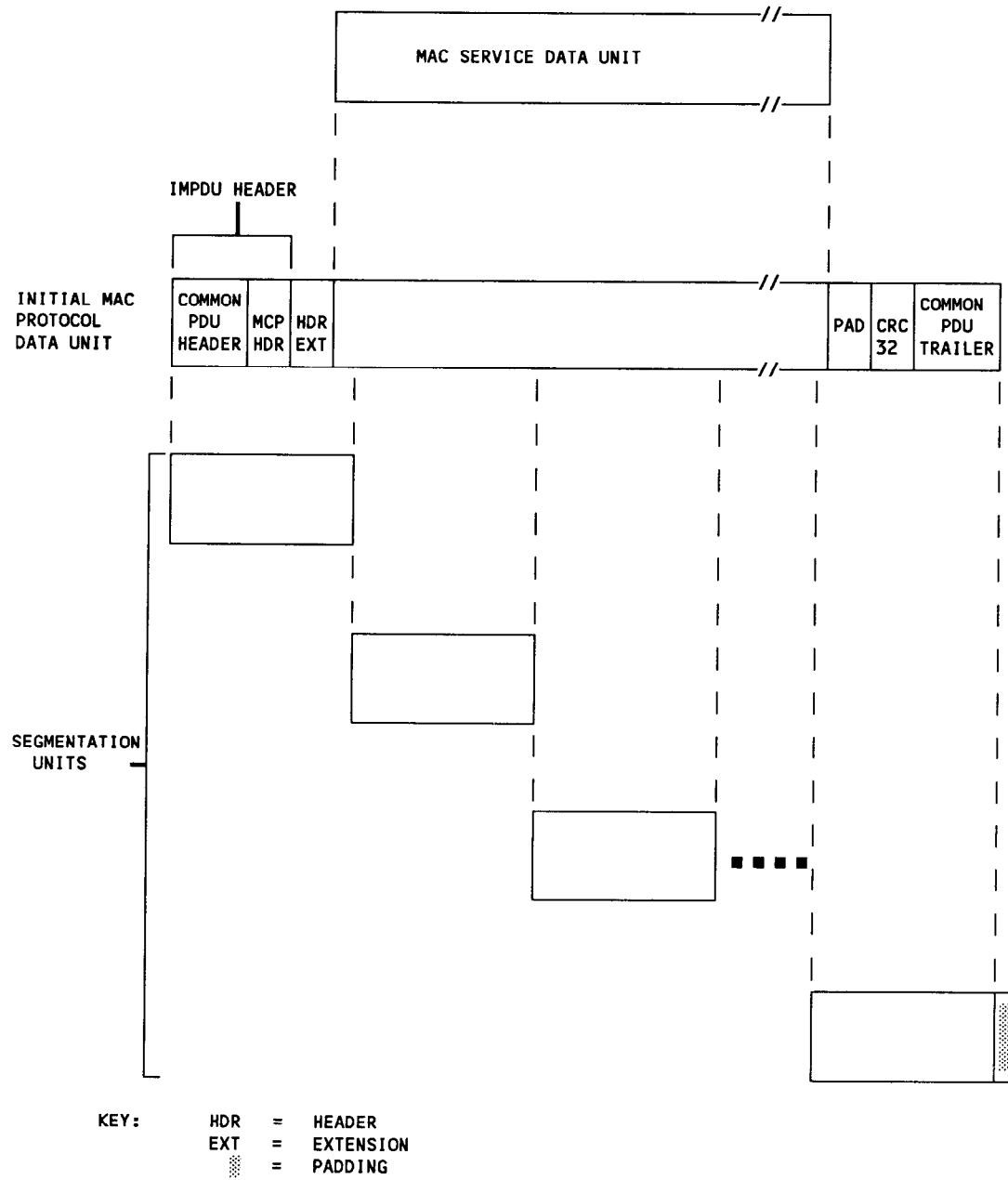


Figure 2-8—Segmentation of an IMPDU

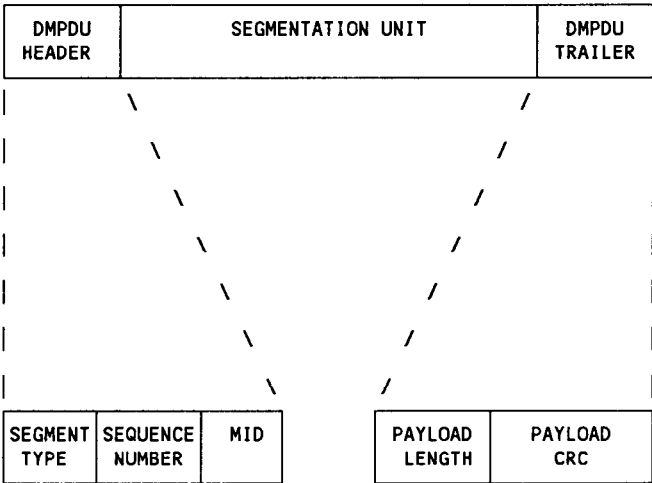


Figure 2-9—Format of a DMPDU

The MID identifies all of the DMPDUs derived from a single IMPDU, and is used in the reassembly of the segmented IMPDU at the destination. To describe the operation of this scheme, the segmentation of the IMPDU at the source is described first and the reassembly at the destination is considered later.

2.1.3.1.1 Segmentation at the source

The train of DMPDUs sent as QA segment payloads by the source is shown in figure 2-10. The first segmentation unit of a multi-segmentation unit IMPDU is carried in a Beginning of Message (BOM) DMPDU. This DMPDU is identified by the BOM code in the Segment Type subfield and signifies the start of a new IMPDU transfer. The MID subfield of this DMPDU carries one of the MIDs obtained by the source and not currently being used for the sending of another IMPDU by the source. The Sequence Number subfield of this DMPDU carries the initial value of the sequence of numbers to be associated with sequential DMPDUs. The segmentation unit of the BOM DMPDU will include the IMPDU header and any header extension plus the first octets of the MSDU, sufficient to fill the segmentation unit.

All subsequent segmentation units of the IMPDU until the last are placed in the payload field of the train of segments following the first segment. These DMPDUs are identified by the COM (Continuation of Message) code in the Segment Type subfield. The transfer of the multisegmentation unit IMPDU is completed by sending the last segmentation unit in a DMPDU that contains the EOM (End of Message) code in the Segment Type subfield. The COM and EOM DMPDUs carry the value of Sequence Number in the BOM DMPDU incremented by one for each successive DMPDU. All COM DMPDUs and the EOM DMPDU derived from an IMPDU carry the same MID value as the BOM DMPDU.

For the transfer of an IMPDU that only requires a single segmentation unit, the SSM (Single Segment Unit Message) code is used in the Segment Type subfield of the DMPDU. The MID is not used in this case, thus it is set to the reserved value of zero.

Each DMPDU trailer is constructed by writing into the Payload Length subfield the number of IMPDU octets used in the segmentation unit. For the connectionless MAC service to LLC, this number is always 44 for BOM and COM DMPDUs. The number written into EOM DMPDUs indicates the remaining number of octets in the IMPDU that need to be transferred. This number can be any multiple of 4 in the range 4 to 44, inclusive. The number written into SSM DMPDUs indicates the length of the entire IMPDU. This number

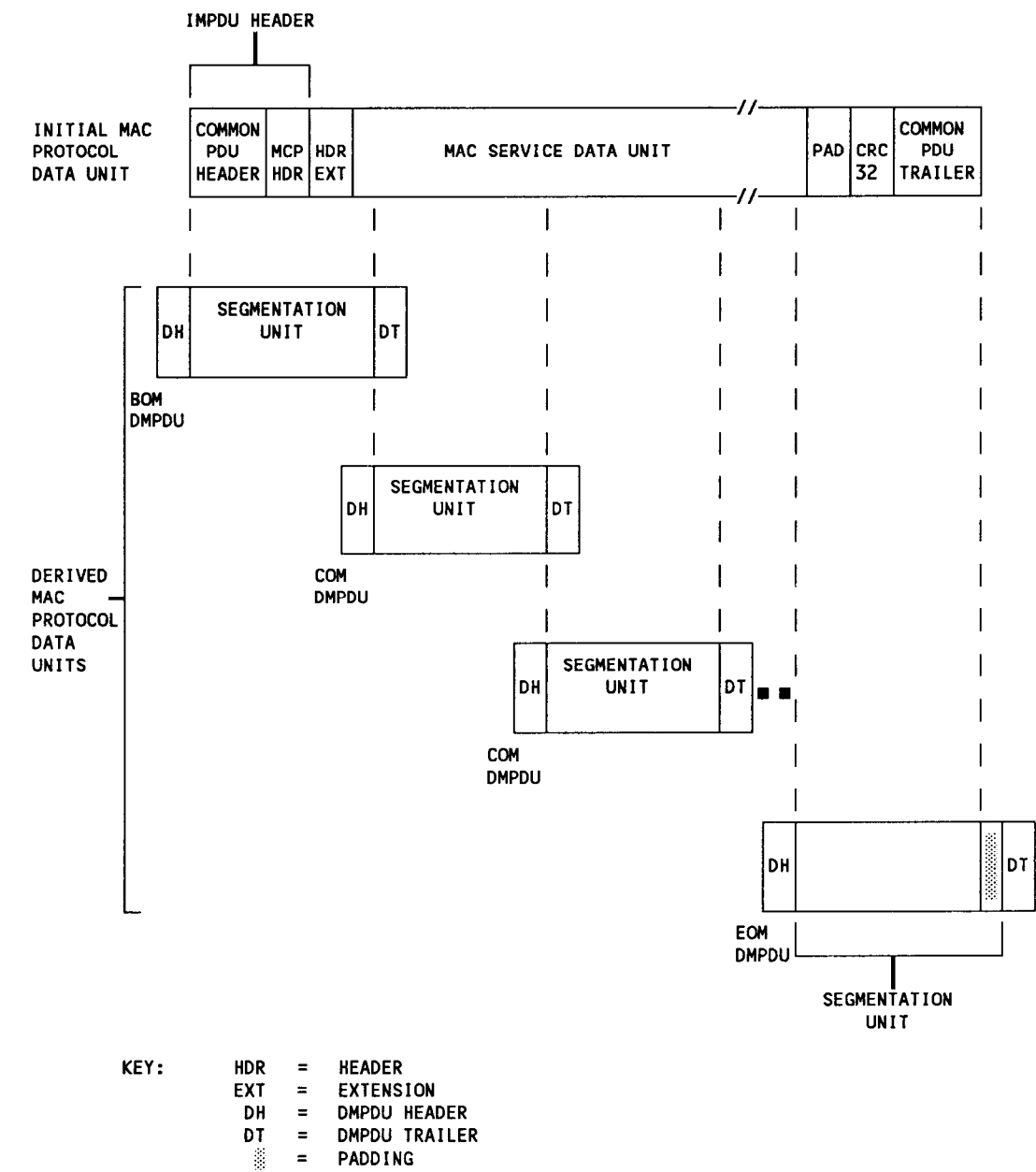


Figure 2-10—Transfer of an IMPDU

can be any multiple of 4 in the range 28 to 44, inclusive. The Payload CRC subfield is a CRC computed over all octets of the segment payload, including the DMPDU header, segmentation unit, and DMPDU trailer.

2.1.3.1.2 Reassembly at the destination

To receive IMPDUs segmented as described above, each AU will monitor all segments passing on the bus. All DMPDUs contain one of a particular set of VCI values in the segment header.¹⁴ If the VCI value is one that the AU is programmed to receive, the AU will verify the DMPDU by means of the Payload CRC subfield. If the CRC verification fails then the DMPDU is discarded.

For each valid DMPDU with a BOM code in the Segment Type subfield, the AU will inspect the MCP header, which will be within the BOM segmentation unit. If the MCP header indicates that the IMPDU is addressed to the AU then it will copy the BOM segmentation unit and continue as described below. If the IMPDU is not addressed to the AU, the segmentation unit is not copied.

To receive the remainder of the IMPDU associated with the BOM segmentation unit, the AU will also record the sequence number and MID value from the BOM DMPDU. The following DMPDUs derived from the same IMPDU should all be received with the same VCI value in the segment header, and an incremented sequence number for each successive DMPDU, and the same MID value in the DMPDU header. The DMPDUs are recognized by the AU using the Payload CRC subfield to verify the DMPDU header of all DMPDUs received on the same VCI, and then comparing the verified MID value with the one recorded. When a match is made, the AU copies the verified segmentation unit of the DMPDU, provided the sequence number is the expected value. The complete IMPDU is received when a verified DMPDU with matching MID value, expected value of sequence number, and EOM code in the Segment Type subfield is received and the segmentation unit copied. Since segments are guaranteed to be delivered in order across the DQDB subnetwork, the AU can reassemble the received segmentation units into the original IMPDU by concatenating them in the order they were received.

The concatenation of all received segmentation units of an IMPDU is finally verified by using two pieces of information contained in the Common PDU header and Common PDU trailer. The length of the IMPDU, minus the length of the Common PDU header and Common PDU trailer, is sent in both the Common PDU header and Common PDU trailer. The length value received in the Common PDU trailer is compared against the number of octets received for the IMPDU. A mismatch causes the receiver to discard the IMPDU. This check is used to ensure that the correct number of DMPDUs have been received, and thus protects against loss or insertion of COM DMPDUs.

The second piece of information is a Beginning-End tag (BETag). The same value of BETag is sent in both the Common PDU header and Common PDU trailer of a given IMPDU. The BETag value is incremented for the next IMPDU sent by the node. The two BETag values for an IMPDU are compared at the receiver, and a mismatch causes the receiver to discard the IMPDU. The BETag is used to ensure that the BOM DMPDU and the EOM DMPDU of a reassembled IMPDU were both actually derived from the same source IMPDU. This protects against loss of the EOM DMPDU from one IMPDU, loss of the BOM DMPDU from a subsequent IMPDU, and loss of the appropriate number of COM DMPDUs such that the IMPDU reassembled from the received DMPDUs is still of the length specified in the Common PDU trailer.

If the node supports checking of the IMPDU using the 32-bit CRC, and if the CRC32 field is present in the received IMPDU, as indicated by a bit in the MCP header, the AU will verify the IMPDU by means of the CRC32 field. If the CRC verification fails, then the IMPDU is discarded. If the CRC verification passes, the IMPDU is accepted as valid.

If all of these comparisons pass, then the IMPDU is accepted as valid.

¹⁴ Note that the VCI value of all bits being set to one is reserved as a default value for MAC service to LLC. All conforming stations must recognize this VCI value and process the DMPDUs.

On the receipt of all DMPDUs of the IMPDU, the recorded MID value must be cleared because the source may reuse the same MID for a different IMPDU transfer.

To protect against the loss of the EOM DMPDU, which normally terminates the IMPDU transfer, a timer must be used to determine how long the MID value remains valid within the AU. If no EOM DMPDU with matching MID value is received within the timer period, the recorded MID is cleared and the received segmentation units associated with the MID are discarded.

With the above scheme, an AU can receive and reassemble multiple IMPDUs concurrently. This is achieved by recording multiple MID values for each VCI which the AU is programmed to receive and receiving all DMPDUs which match the recorded MID values on the appropriate VCIs.

Single segmentation unit IMPDUs are received in a similar manner to the BOM DMPDUs. The AU will verify the DMPDU using the Payload CRC subfield, inspect the MCP Header field in the segmentation unit, and copy the segmentation unit if the Payload CRC passes and the IMPDU is addressed to the AU. However, since with SSMs there are no following DMPDUs, there is no need to record the MID value. The IMPDU is completely received in this first DMPDU, and is then validated using the Length value in the Common PDU trailer, the BTag values in the Common PDU header and Common PDU trailer, and possibly the CRC32 field as described above.

2.1.3.2 Provision of isochronous service

The Pre-Arbitrated access control mechanism does not necessarily accept or deliver isochronous service octets in an isochronous fashion. The variation from isochronous delivery occurs when the node at the head of bus does not generate Pre-Arbitrated slots in an isochronous fashion. If the user of the isochronous service requires the DQDB Layer to accept and/or deliver the octets isochronously, there is a requirement for some buffering. The nature of this buffering depends on the isochronous service user and is performed as part of an isochronous convergence function.

The isochronous service described in this part of ISO/IEC 8802 provides only the DQDB Layer functions required to access the medium to read and write isochronous service octets. This includes the Pre-Arbitrated Access functions within each AU and the periodic generation of the PA slots at the head of bus. The signaling procedures for establishing, maintaining, and clearing a connection are outside the scope of this part of ISO/IEC 8802. Similarly, the signaling procedures for indicating to the head of bus which PA slots are to be generated are also outside the scope of this part of ISO/IEC 8802.

However, annex B gives a tutorial description of the sequence of events that would be required to support complete isochronous communication capability. In particular, it explains the interactions that occur via the DQDB LMI and are defined in this part of ISO/IEC 8802.

2.1.3.3 Provision of Connection-Oriented Data Service

The Queued Arbitrated access control mechanism provides support for the transfer of QA segments. Therefore, support for the MAC service to LLC (described in 2.1.3.1) requires the definition of a convergence function.

Another convergence function is under study that will allow for the DQDB Layer to support a connection-oriented data service that uses Queued Arbitrated access. This will require some functions that differ from the convergence functions to support the MAC to LLC service. However, the connection-oriented data service will use a segmentation and reassembly procedure that is common with that used to support MAC service to LLC. In particular, it will use the same format for the Common PDU header and Common PDU trailer. It will also carry the same fields in QA segment payloads as are carried in the DMPDU header and DMPDU trailer.